
Fondamenti della Programmazione: Metodi Evoluti

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Lezione 5: Logica

Reminder: contracts

Associated with an individual feature:

- Pre-conditions (must be true BEFORE feature execution)
- Post-conditions (must be true AFTER feature execution)

Associated with a class:

- Class invariants (expresses consistency requirements between queries of a class)

Logic

How to express conditions in contracts?

We need a mathematical notation since conditions have to be automatically checked

Logic is the answer!

Reasoning and programming

Logic is the basis of

- Mathematics: proofs are only valid if they follow the rules of logic.
- Software development:
 - Conditions in program actions: “If m is positive, then execute this instruction”
 - Conditions in contracts:
“ x must not be zero, so that we can calculate $\frac{x+7}{x}$ ”

Boolean expressions

A condition is expressed as a **boolean expression**.

It consists of

- **Boolean variables** (identifiers denoting boolean values)
- **Boolean operators** (**not**, **or**, **and**, **=**, **implies**)

and represents possible

- **Boolean values** (truth values, either **True** or **False**)

Examples

Examples of boolean expressions
(with *rain_today* and *cuckoo_sang_last_night* as boolean variables):

- *rain_today*
(a boolean variable is a boolean expression)
- **not** *rain_today*
- **(not** *cuckoo_sang_last_night*) **implies** *rain_today*

(Parentheses group sub-expressions)

Negation (not)

a	$\text{not } a$
True	False
False	True

For any boolean expression e and any values of variables:

- Exactly one of e and **not** e has value **True**
- Exactly one of e and **not** e has value **False**
- One of e and **not** e has value **True** (Principle of the Excluded Middle)
- Not both of e and **not** e have value **True** (Principle of Non-Contradiction)

Disjunction (or)

a	b	a or b
True	True	True
True	False	True
False	True	True
False	False	False

or operator is **commutative**

or operator is **associative**:

- a or (b or c) = (a or b) or c

Disjunction principle:

- An **or** disjunction has value **True** except if both operands have value **False**

NB: differently from 'or' in common language **or** is non-exclusive

Conjunction (and)

<i>a</i>	<i>b</i>	<i>a and b</i>
True	True	True
True	False	False
False	True	False
False	False	False

and operator is commutative

and operator is associative

- $a \text{ and } (b \text{ and } c) = (a \text{ and } b) \text{ and } c$

Conjunction principle:

- An **and** conjunction has value **False** except if both operands have value **True**

Truth assignment and truth table

Truth assignment for a set of variables: particular choice of values (**True** or **False**), for every variable

A truth assignment satisfies an expression if the value for the expression is **True**

A truth table for an expression with n variables has

- $n + 1$ columns
- 2^n rows

Combined truth table for basic operators

<i>a</i>	<i>b</i>	not <i>a</i>	<i>a</i> or <i>b</i>	<i>a</i> and <i>b</i>
True	True	False	True	True
True	False		True	False
False	True	True	True	False
False	False		False	False

Tautologies

Tautology: a boolean expression that has value **True** for every possible truth assignment

Examples:

- $a \text{ or } (\text{not } a)$
- $\text{not } (a \text{ and } (\text{not } a))$
- $(a \text{ and } b) \text{ or } ((\text{not } a) \text{ or } (\text{not } b))$

Contradictions

Contradiction: a boolean expression that has value **False** for every possible truth assignment

Examples:

- a and (not a)
- not (a or (not a))

Satisfiable: for at least one truth assignment the expression yields **True**

- Any tautology is satisfiable
- No contradiction is satisfiable.

Equivalence (=)

a	b	$a = b$
True	True	True
True	False	False
False	True	False
False	False	True

= operator is commutative ($a = b$ has same value as $b = a$)

= operator is reflexive ($a = a$ is a tautology for any a)

Substitution:

- For any expressions u , v and e , if $u = v$ is a tautology and e' is the expression obtained from e by replacing every occurrence of u by v , then $e = e'$ is a tautology

Types of propositions

Tautology

- **True** for all truth assignments
 - P **or** (**not** P)
 - **not** (P **and** (**not** P))
 - (P **and** Q) **or** ((**not** P) **or** (**not** Q))

Contradiction

- **False** for all truth assignments
 - P **and** (**not** P)

Satisfiable

- **True** for at least one truth assignment

Equivalent

- φ and χ are equivalent if they are satisfied under exactly the same truth assignments, or if $\varphi = \chi$ is a tautology

De Morgan's laws

They show the duality between **and** and **or**: negating an expression is equivalent to negating variables and swapping **and** and **or**

Tautologies

- $\text{not } (a \text{ or } b) = (\text{not } a) \text{ and } (\text{not } b)$
- $\text{not } (a \text{ and } b) = (\text{not } a) \text{ or } (\text{not } b)$

- $a \text{ or } b = \text{not } (\text{not } a) \text{ and } (\text{not } b)$
- $a \text{ and } b = \text{not } (\text{not } a) \text{ or } (\text{not } b)$

More tautologies (distributivity):

- $(a \text{ and } (b \text{ or } c)) = ((a \text{ and } b) \text{ or } (a \text{ and } c))$
- $(a \text{ or } (b \text{ and } c)) = ((a \text{ or } b) \text{ and } (a \text{ or } c))$

Syntax convention: binding of operators

Order of binding (starting with tightest binding): **not**, **and**, **or**, **implies** (to be introduced), = .

Style rules:

When writing a boolean expression, drop the parentheses:

- Around the expressions of each side of "=" if whole expression is an equivalence.
- Around successive elementary terms if they are separated by the same associative operators.

Implication (implies)

a	b	a implies b
True	True	True
True	False	False
False	True	True
False	False	True

a **implies** b , for any a and b , is the value of **(not a) or b**

In a **implies** b : a is **antecedent**, b **consequent**

Implication principle:

- An implication has value **True** except if its antecedent has value **True** and its consequent has value **False**
- In particular, always **True** if antecedent is **False**

Implication in ordinary language

implies in ordinary language often means causation, as in “if ... then ...”

- *“If the weather stays like this, skiing will be great this weekend”*
- *“If you put this stuff in your hand luggage, they won’t let you through.”*

Misunderstanding implication

Whenever a is **False**,
 a **implies** b is **True**, regardless of b :

- “Today is Wednesday implies $2+2=5$.”
- “ $2+2=5$ implies today is Wednesday.”

Both of the above implications are **True**

Cases in which a is **False** tell us nothing about the truth of the consequent

Reversing implications (1)

It is not generally true that

~~$a \text{ implies } b = (\text{not } a) \text{ implies } (\text{not } b)$~~

Example (wrong!):

- *“All the people in Rome who live near Spanish Steps are rich. I do not live near Spanish Steps, so I am not rich.”*

live_near_spanish_steps **implies** *rich* [1]

~~$(\text{not } \textit{live_near_spanish_steps}) \text{ implies } (\text{not } \textit{rich})$~~ [2]

Reversing implications (2)

Correct:

$$a \text{ implies } b = (\text{not } b) \text{ implies } (\text{not } a)$$

Example:

- “*All the people who live near Spanish Steps are rich. She is not rich, so she can't be living near Spanish Steps*”

$$\text{live_near_spanish_steps} \text{ implies } \text{rich} = \\ (\text{not } \text{rich}) \text{ implies } (\text{not } \text{live_near_spanish_steps})$$

Semistrict boolean operators (1)

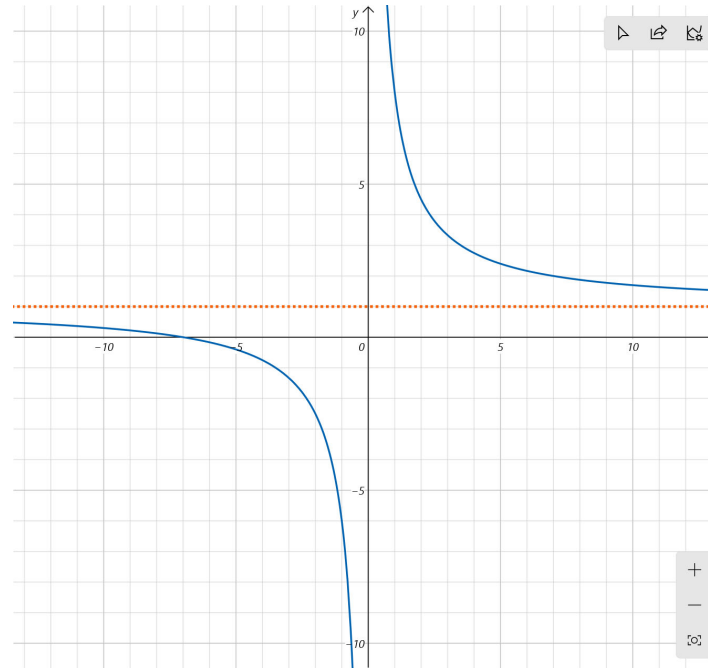
Example boolean-valued expression (x is an integer):

$$\frac{x + 7}{x} > 0$$

True for $x < -7$ or $x > 0$

False for $x \geq -7$ and $x < 0$

Undefined for $x = 0$



Semistrict boolean operators (2)

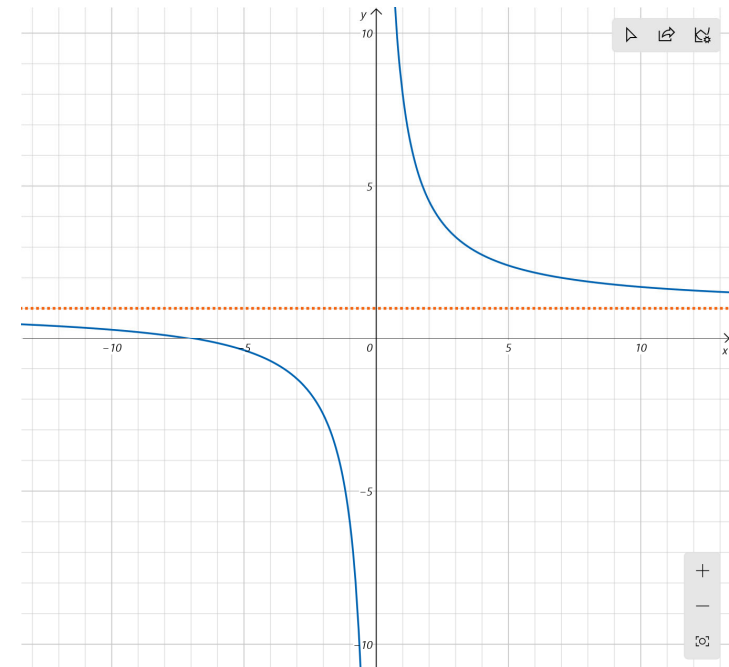
- Avoid division by zero:

$$(x \neq 0) \text{ and } ((x + 7) / x) > 0$$

True for $x < -7$ or $x > 0$

False for $x \geq -7$ and $x \leq 0$

Always defined



Semistrict boolean operators (3)

BUT:

- What happens when evaluating it?
- Program would crash during evaluation of division

We need a **non-commutative** version of **and** (and **or**):

Semistrict boolean operators

Semistrict operators (and then, or else)

a **and then** b : has same value as a **and** b if a and b are both defined, and has **False** whenever a has value **False** even if b is undefined

a **or else** b : has same value as a **or** b if a and b are both defined, and has **True** whenever a has value **True** even if b is undefined

$(x \neq 0)$ **and then** $((x + 7) / x) > 0$

Semistrict operators allow us to define an order of expression evaluation (left to right).

Important for programming when undefined objects may cause program crashes

Ordinary vs. Semistrict boolean operators

Use

- Ordinary boolean operators (**and** and **or**) if you can guarantee that both operands are defined
- **and then** if a condition only makes sense when another is true
- **or else** if a condition only makes sense when another is false

Example:

- “If you are not single, then your spouse must sign the contract”

is_single **or else** *spouse_must_sign*

not *is_single* **and then** *spouse_must_sign*

Semistrict implication

Example with implies:

- “If you are not single, then your spouse must sign the contract.”

(**not** *is_single*) **implies** *spouse_must_sign*

Definition of **implies**: in our case, **always semistrict!**

- $a \text{ implies } b = (\text{not } a) \text{ or else } b$

Strict or semi-strict?

Hands-On

- $a = 0$ **or** $b = 0$
- $a \neq 0$ **and** $b // a \neq 0$
- $a \neq \mathbf{Void}$ **and** $b \neq \mathbf{Void}$
- $a < 0$ **or** $\text{sqrt}(a) > 2$
- $(a = b$ **and** $b \neq \mathbf{Void})$ **and** $a.\text{age} = 0$

Programming language notation for boolean operators

Eiffel keyword	Common mathematical symbol
not	\sim or \neg
or	\vee
and	\wedge
=	\Leftrightarrow
implies	\Rightarrow

Propositional and predicate calculus

Propositional calculus:

property p holds for a single object

Predicate calculus:

property p holds for several objects

Generalizing or

G : group of objects, p : property

Generalization of **or**:

Does *at least one* of the objects in G satisfy p ?

Is at least one station of Line 8 an exchange?

Station_Balard.is_exchange **or** *Station_Lourmel.is_exchange* **or**
Station_Boucicaut.is_exchange **or**
... (all stations of Line 8)

Existential quantifier: *exists*, or \exists

$\exists s : Stations_8 \mid s.is_exchange$

“There exists an s in *Stations_8*
such that $s.is_exchange$ is true”

Generalizing and

Generalization of **and**:

Does *every* object in G satisfy p ?

Are all stations of Tram 8 exchanges?

Station_Balard.is_exchange **and** *Station_Lourmel.is_exchange*
and *Station_Boucicaut.is_exchange* **and** ...

(all stations of Line 8)

Universal quantifier: *for_all*, or \forall

$\forall s: Stations_8 \mid s.is_exchange$

“For all s in *Stations8* | *s.is_exchange* is true”

Existentially quantified expression

Boolean expression:

$$\exists s : SOME_SET \mid s.some_property$$

- *True* if and only if at least one member of *SOME_SET* satisfies property *some_property*

Proving

- **True:** Find one element of *SOME_SET* that satisfies the property
- **False:** Prove that no element of *SOME_SET* satisfies the property (test all elements)

Universally quantified expression

Boolean expression:

$$\forall s: SOME_SET \mid s.some_property$$

- *True* if and only if every member of *SOME_SET* satisfies property *some_property*

Proving

- **True:** Prove that every element of *SOME_SET* satisfies the property (test all elements)
- **False:** Find one element of *SOME_SET* that does not satisfies the property

Duality

Generalization of DeMorgan's laws:

$$\mathbf{not} (\exists s : SOME_SET \mid P) = \forall s : SOME_SET \mid \mathbf{not} P$$

$$\mathbf{not} (\forall s : SOME_SET \mid P) = \exists s : SOME_SET \mid \mathbf{not} P$$

Empty sets

$\exists s: \text{SOME_SET} \mid \text{some_property}$

If *SOME_SET* is empty: always **False**

$\forall s: \text{SOME_SET} \mid \text{some_property}$

If *SOME_SET* is empty: always **True**

Tautology / contradiction / satisfiable?

Let the range of variables be INTEGER

$x < 0$ **or** $x \geq 0$

tautology

$x > 0$ **implies** $x > 1$

satisfiable

$\forall x \mid x > 0$ **implies** $x > 1$

contradiction

$\forall x \mid x * y = y$

satisfiable

$\exists y \mid \forall x \mid x * y = y$

tautology

Hands-On