Preliminary considerations, $A \in \mathbb{C}^{n \times n}$

Given a $n \times n$ matrix A with eigenvalues $\lambda_1, \lambda_2, \ldots, \lambda_n$, the following result holds.

Theorem (inverse power). If λ_i^* is an approximation of the eigenvalue λ_i of a $n \times n$ matrix A, i.e. $|\lambda_i - \lambda_i^*|$ is smaller than $|\lambda_j - \lambda_i^*|, \forall \lambda_j \neq \lambda_i$, if $m_a(\lambda_i) = m_g(\lambda_i)$, and if $\mathbf{v}_0 \in \mathbb{C}^n$ is not orthogonal to the space spanned by the eigenvectors of Acorresponding to λ_i , then the sequence $\{\mathbf{v}_k\}$ generated by the algorithm

$$(A - \lambda_i^* I)\mathbf{a}_k = \mathbf{v}_{k-1}, \ \mathbf{v}_k = \frac{\mathbf{a}_k}{\|\mathbf{a}_k\|}, \quad k = 1, 2, \dots$$

converges to an eigenvector of A corresponding to the eigenvalue λ_i . If A is diagonalizable, then the rate of convergence is $O((\max_{j: \lambda_j \neq \lambda_i} | \frac{\lambda_i - \lambda_i^*}{\lambda_i - \lambda_i^*} |)^k)$.

Remark. In the general case the rate of convergence is

$$\max_{j: \lambda_j \neq \lambda_i} \max_{s_{\lambda_j}} O(|p_{s_{\lambda_j}-1}(k)|| \frac{\lambda_i - \lambda_i^*}{\lambda_j - \lambda_i^*} |^k)$$

where, given the block diagonal Jordan form of $A, J = X^{-1}AX$, for each $\lambda_j \neq \lambda_i$, the number s_{λ_j} indicates the dimension of the generic Jordan block associated with λ_j , and $p_{s_{\lambda_j}-1}(k)$ is a polynomial of degree $s_{\lambda_j} - 1$, whose coefficients depend on $\frac{1}{(\lambda_j - \lambda_i^*)^r}$, $r = 1, \ldots, s_{\lambda_j} - 1$, and on the coefficients of \mathbf{v}_0 with respect to the column vectors of X corresponding to the Jordan block under consideration.

proof: See the Appendix. \Box

Let λ_1 be such that $|\lambda_1| = \rho(A)$ and assume that all λ_i such that $|\lambda_i| = \rho(A)$ are equal to λ_1 (in such case we say that λ_1 dominates the eigenvalues of A). Assume, moreover, that the algebraic and geometric multiplicity of λ_1 are equal. Then the *power method* (see the Theorem below) can be used to compute λ_1 and an eigenvector corresponding to λ_1 .

Theorem (power). If λ_1 dominates the eigenvalues of a $n \times n$ matrix A, if $m_a(\lambda_1) = m_g(\lambda_1)$, and if $\mathbf{v}_0 \in \mathbb{C}^n$ is not orthogonal to the space spanned by the eigenvectors of A corresponding to λ_1 , then the sequence $\{\mathbf{v}_k\}$ generated by the algorithm

$$\mathbf{a}_k = A\mathbf{v}_{k-1}, \ \mathbf{v}_k = \frac{\mathbf{a}_k}{\|\mathbf{a}_k\|}, \quad k = 1, 2, \dots$$

converges to an eigenvector of A corresponding to the eigenvalue λ_1 . Moreover,

$$rac{\mathbf{u}^H \mathbf{v}_{k+1}}{\mathbf{u}^H \mathbf{v}_k} \to \lambda_1, \quad k o +\infty$$

for any **u** for which $\mathbf{u}^H \mathbf{v}_k \neq 0$. If A is diagonalizable, then the rate of convergence is $O((\max_{j: \lambda_j \neq \lambda_1} | \frac{\lambda_j}{\lambda_1} |)^k)$.

Remark. In the general case the rate of convergence is

$$\max_{j: \lambda_j \neq \lambda_1} \max_{s_{\lambda_j}} O(|p_{s_{\lambda_j}-1}(k)|| \frac{\lambda_j}{\lambda_1}|^k),$$

where, given the block diagonal Jordan form of A, $J = X^{-1}AX$, for each $\lambda_j \neq \lambda_1$, the number s_{λ_j} indicates the dimension of the generic Jordan block associated with λ_j , and $p_{s_{\lambda_j}-1}(k)$ is

a polynomial of degree $s_{\lambda_j} - 1$, whose coefficients depend on $\frac{1}{\lambda_i^r}$, $r = 1, \ldots, s_{\lambda_j} - 1$, and on the coefficients of \mathbf{v}_0 with respect to the column vectors of X corresponding to the Jordan block under consideration.

proof: See the Appendix. \Box

For our purposes it is useful to recall also the following classic result on matrix deflation: how to introduce a matrix whose eigenvalues are all equal to the eigenvalues of A except one, which, instead, is zero.

Theorem. Let A be a $n \times n$ matrix. Let λ_1 be a nonzero eigenvalue of A and \mathbf{y}_1 a corresponding eigenvector, i.e. $A\mathbf{y}_1 = \lambda_1 \mathbf{y}_1$. Call $\lambda_2, \lambda_3, \dots, \lambda_n$ the remaining eigenvalues of A. Then the matrix $W = A - \frac{\lambda_1}{\mathbf{w}^* \mathbf{y}_1} \mathbf{y}_1 \mathbf{w}^*, \ \mathbf{w}^* \mathbf{y}_1 \neq 0$, has eigenvalues $0, \lambda_2, \lambda_3, \ldots, \lambda_n$.

proof. Introduce $S = [\mathbf{y}_1 \mathbf{z}_2 \cdots \mathbf{z}_n]$ non singular, and observe that $p_A(\lambda) =$ $p_{S^{-1}AS}(\lambda) = (\lambda - \lambda_1)q(\lambda), \ p_W(\lambda) = p_{S^{-1}WS}(\lambda) = \lambda q(\lambda).$

Let G be the following matrix

$$G = \begin{bmatrix} \frac{1}{4} & \frac{1}{4} & \frac{1}{2} \\ \frac{3}{4} & \frac{1}{8} & \frac{1}{8} \\ \frac{11}{16} & \frac{1}{4} & \frac{1}{16} \end{bmatrix}.$$

Since $\rho(G) \leq ||G||_{\infty} = 1$, we can say that the spectrum of G lies in the circle $\{z \in \mathbb{C} : |z| \le 1\}.$

Note that $G\mathbf{e} = 1 \cdot \mathbf{e}$, i.e. one eigenvalue of G, $\lambda_1 = 1$, and its corresponding eigenvector, $\mathbf{y}_1 = \mathbf{e} = \begin{bmatrix} 1 & \cdots & 1 \end{bmatrix}^T$, are known. If λ_2, λ_3 denote the remaining eigenvalues of G, then we can define a matrix W, in terms of G, λ_1 , \mathbf{y}_1 , whose eigenvalues are $0, \lambda_2, \lambda_3$:

$$W = G - \frac{\lambda_1}{\mathbf{w}^* \mathbf{y}_1} \mathbf{y}_1 \mathbf{w}^* = G - \frac{1}{\mathbf{w}^* \mathbf{e}} \begin{bmatrix} \mathbf{w}^* \\ \mathbf{w}^* \\ \mathbf{w}^* \end{bmatrix}, \quad \forall \mathbf{w}, \ \mathbf{w}^* \mathbf{e} \neq 0.$$

Since $(\mathbf{e}_i^T G)\mathbf{e} = \mathbf{e}_i^T (1 \cdot \mathbf{e}) = 1 \neq 0$, we choose $\mathbf{w}^* = \mathbf{e}_i^T G$:

$$W = G - \begin{bmatrix} \mathbf{e}_i^T G \\ \mathbf{e}_i^T G \\ \mathbf{e}_i^T G \end{bmatrix}.$$

For i = 1 the matrix W becomes

$$W = \begin{bmatrix} \frac{1}{4} & \frac{1}{4} & \frac{1}{2} \\ \frac{3}{4} & \frac{1}{8} & \frac{1}{8} \\ \frac{11}{16} & \frac{1}{4} & \frac{1}{16} \end{bmatrix} - \begin{bmatrix} \frac{1}{4} & \frac{1}{4} & \frac{1}{2} \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{2} \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{2} \end{bmatrix} = \begin{bmatrix} 0 & 0 & 0 \\ \frac{1}{2} & -\frac{1}{8} & -\frac{3}{8} \\ \frac{7}{16} & 0 & -\frac{7}{16} \end{bmatrix},$$

so, the eigenvalues of A different from 1 are the eigenvalues of

$$\tilde{W} = \begin{bmatrix} -\frac{1}{8} & -\frac{3}{8} \\ 0 & -\frac{7}{16} \end{bmatrix},$$

i.e. $-\frac{1}{8}$ and $-\frac{7}{16}$. Thus, 1, $-\frac{1}{8}$ and $-\frac{7}{16}$ are the eigenvalues of G, and also, of course, of G^T . In particular, 1 is eigenvalue of G^T , but note that the eigenvector \mathbf{p} of G^T

corresponding to 1 is not obvious; it must be computed. For example, it can be computed as the limit of the inverse power sequence $\{\mathbf{v}_k\}$ defined as follows (ε small positive number):

$$\mathbf{v}_0 \in \mathbb{R}^3, \ (G^T - (1 + \varepsilon)I)\mathbf{a}_k = \mathbf{v}_{k-1}, \ \mathbf{v}_k = \frac{\mathbf{a}_k}{\|\mathbf{a}_k\|}, \ k = 0, 1, 2, \dots$$

(rate of convergence: $O(\left|\frac{1-(1+\varepsilon)}{-\frac{1}{8}-(1+\varepsilon)}\right|^k)$). [Here a reference for the inverse power iterations]. We shall see that the vector **p** can be also obtained as the limit of the sequence

$$\mathbf{p}_0 \in \mathbb{R}^3$$
, \mathbf{p}_0 positive, $\|\mathbf{p}_0\|_1 = 1$, $\mathbf{p}_{k+1} = G^T \mathbf{p}_k$, $k = 0, 1, 2, ...$

(rate of convergence: $O(\left|\frac{-\frac{7}{16}}{1}\right|^k)$) [this result is in fact a particular case of the Theorem at the beginning of this section (set $\|\cdot\| = \|\cdot\|_1$, $\mathbf{u} = \mathbf{e}$)]. Even if the rate of convergence of the \mathbf{p}_k is not as good as the rate of convergence of the \mathbf{v}_k , the computation of \mathbf{p}_{k+1} from \mathbf{p}_k is much cheaper than the computation of \mathbf{v}_{k+1} from \mathbf{v}_k . In fact, for analogous problems, but of high dimension, even one step of the inverse power iterations is prohibitive.

The Perron-Frobenius theory: $A \in \mathbb{R}^{n \times n}$, $A \ge 0$ irreducible [Varga]

Lemma. Let A be a $n \times n$ non negative matrix, i.e. $a_{ij} \ge 0, \forall i, j$. Assume that A is not reducible. Then $(I + A)^{n-1}$ is a positive matrix, i.e. its entries are all positive.

proof. We shall prove that the vector $(I + A)^{n-1}\mathbf{x}$ is positive whenever \mathbf{x} is a non negative non null vector (prove that this is equivalent to the thesis!).

Let \mathbf{x} be a non negative non null vector. Set $\mathbf{x}_0 = \mathbf{x}$, $\mathbf{x}_1 = (I + A)\mathbf{x}_0 = \mathbf{x}_0 + A\mathbf{x}_0, \ldots, \mathbf{x}_{k+1} = (I + A)\mathbf{x}_k = \mathbf{x}_k + A\mathbf{x}_k, \ k = 1, \ldots, n-2$. Note that $\mathbf{x}_k = (I + A)^k \mathbf{x}$, in particular $\mathbf{x}_{n-1} = (I + A)^{n-1} \mathbf{x}$. So, our aim is to prove that \mathbf{x}_{n-1} is a positive vector. First observe by induction that all \mathbf{x}_k are non negative vectors (\mathbf{x}_k non negative and A non negative imply $A\mathbf{x}_k$ non negative and $\mathbf{x}_{k+1} = \mathbf{x}_k + A\mathbf{x}_k$ non negative). Then the thesis (\mathbf{x}_{n-1} positive) is now proved by showing that \mathbf{x}_{k+1} must have less zeros than \mathbf{x}_k for each $k \in \{0, \ldots, n-2\}$. Note that \mathbf{x}_{k+1} cannot have more zeros than \mathbf{x}_k , since $A\mathbf{x}_k$, in the definition $\mathbf{x}_k + A\mathbf{x}_k$ of \mathbf{x}_{k+1} , is non negative. Assume \mathbf{x}_{k+1} has the same number of zero entries as \mathbf{x}_k . Of course such zeros must be in the same places, i.e. there exists a permutation matrix P such that $P\mathbf{x}_k = \begin{bmatrix} \alpha \\ 0 \end{bmatrix}$, $P\mathbf{x}_{k+1} = \begin{bmatrix} \beta \\ 0 \end{bmatrix}$, with α, β positive vectors of the same dimension $m, 1 \le m \le n-1$ (why such bounds for m?). Thus, $P\mathbf{x}_{k+1} = P\mathbf{x}_k + PA\mathbf{x}_k = P\mathbf{x}_k + PAP^T P\mathbf{x}_k$. Consider the following partition of the matrix PAP^T

$$PAP^T = \begin{bmatrix} M_{11} & M_{12} \\ M_{21} & M_{22} \end{bmatrix}$$

where M_{11} is $m \times m$. Then

$$\left[\begin{array}{c}\beta\\0\end{array}\right] = \left[\begin{array}{c}\alpha\\0\end{array}\right] + \left[\begin{array}{c}M_{11} & M_{12}\\M_{21} & M_{22}\end{array}\right] \left[\begin{array}{c}\alpha\\0\end{array}\right],$$

and, in particular, $M_{21}\alpha = 0$. The latter condition implies $M_{21} = 0$, being α a positive vector and M_{21} a non negative matrix. But this is equivalent to say that A is reducible, against the hypothesis! \Box

Example. Set

$$I + A = \begin{bmatrix} 1 & 0 & 0 & 1 \\ a & 1 & 0 & 0 \\ 0 & b & 1 & 0 \\ 0 & 0 & c & 1 \end{bmatrix}, \ a, b, c \text{ positive.}$$

Note that A is a non-negative irreducible matrix, and in fact $(I + A)^3$ is a positive matrix, i.e. its entries are positive. Moreover, 3 is the minimum j for which $(I + A)^j$ is positive.

Let A be a $n \times n$ non negative irreducible matrix. Let **x** be a non negative non null vector, and associate to **x** the number

$$r_{\mathbf{x}} := \min_{i:x_i>0} \frac{\sum_j a_{ij} x_j}{x_i} = \min_{i:x_i>0} \frac{(A\mathbf{x})_i}{x_i}.$$

Proposition. $r_{\mathbf{x}}$ is a non negative real number; $r_{\mathbf{x}} = r_{\alpha \mathbf{x}}$ if $\alpha > 0$; $A\mathbf{x} \ge r_{\mathbf{x}}\mathbf{x}$; $r_{\mathbf{x}} = \sup\{\rho \in \mathbb{R} : A\mathbf{x} \ge \rho \mathbf{x}\}.$

proof: easy, left to the reader.

Now associate to A the following number:

$$r = \sup_{\mathbf{x} \ge \mathbf{0}, \, \mathbf{x} \neq \mathbf{0}} r_{\mathbf{x}} = \sup_{\mathbf{x} \ge \mathbf{0}, \, \mathbf{x} \neq \mathbf{0}} \min_{i:x_i > 0} \frac{\sum_j a_{ij} x_j}{x_i}.$$

Proposition. r is a positive real number; if $\mathbf{w} \ge \mathbf{0}$, $\mathbf{w} \ne \mathbf{0}$, is such that $A\mathbf{w} \ge r\mathbf{w}$, then $A\mathbf{w} = r\mathbf{w}$ and $\mathbf{w} > \mathbf{0}$.

proof. If $\mathbf{e} = [1 \ 1 \ \cdots \ 1]^T$, then $r_{\mathbf{e}} = \min_{i: (\mathbf{e})_i > 0} \frac{\sum_j a_{ij}(\mathbf{e})_j}{(\mathbf{e})_i} = \min_i \sum_j a_{ij} \ge 0$. Assume $r_{\mathbf{e}} = 0$. Then for some k we would have $\sum_j a_{kj} = 0$, so the kth row of A would be null, and thus A would be reducible (exchange the k and n rows), against the hypothesis. It follows that $r \ge r_{\mathbf{e}} > 0$.

proof. Set $\eta = A\mathbf{w} - r\mathbf{w}$. We know that $\eta \ge \mathbf{0}$. Assume $\eta \ne \mathbf{0}$. Then, by the Lemma,

$$\begin{aligned} \mathbf{0} &< (I+A)^{n-1}\eta = (I+A)^{n-1}A\mathbf{w} - (I+A)^{n-1}r\mathbf{w} \\ &= A(I+A)^{n-1}\mathbf{w} - r(I+A)^{n-1}\mathbf{w} = A\mathbf{y} - r\mathbf{y}, \ \mathbf{y} > \mathbf{0}. \end{aligned}$$

i.e. $r < (A\mathbf{y})_i/y_i \forall i$. Thus $r < r_{\mathbf{y}}$, which is absurd. It follows that $\eta = \mathbf{0}$, that is, $A\mathbf{w} = r\mathbf{w}$. Then, we also have $\mathbf{w} > \mathbf{0}$ since $\mathbf{0} < (I+A)^{n-1}\mathbf{w} = (1+r)^{n-1}\mathbf{w}$. \Box

In the following, given $\mathbf{v} \in \mathbb{C}^n$ and $M \in \mathbb{C}^{n \times n}$ we denote by $|\mathbf{v}|$ and |M|, respectively, the column vector $(|v_k|)_{k=1}^n$ and the matrix $(|m_{ij}|)_{i,j=1}^n$.

Theorem. There exists a positive vector \mathbf{z} for which $A\mathbf{z} = r\mathbf{z}$; $r = \rho(A)$; if $B \in \mathbb{C}^{n \times n}$, $|B| \leq A$, then $\rho(B) \leq \rho(A)$; if $B \in \mathbb{C}^{n \times n}$, $|B| \leq A$, $|B| \neq A$ then $\rho(B) < \rho(A)$.

proof: We now show that there exists $\mathbf{z} \geq \mathbf{0}$ such that $r = r_{\mathbf{z}}$. Once this is proved we will have the inequality $A\mathbf{z} \geq r_{\mathbf{z}}\mathbf{z} = r\mathbf{z}$ which implies, by the Proposition, $A\mathbf{z} = r\mathbf{z}$ and $\mathbf{z} > \mathbf{0}$.

$$r = \sup_{\mathbf{x} \ge \mathbf{0}, \, \mathbf{x} \neq \mathbf{0}} r_{\mathbf{x}} = \sup_{\mathbf{x} \ge \mathbf{0}, \, \mathbf{x} \neq \mathbf{0}} r_{\frac{\mathbf{x}}{\|\mathbf{x}\|_2}} = \sup_{\mathbf{x} \ge \mathbf{0}, \, \|\mathbf{x}\|_2 = 1} r_{\mathbf{x}} = (*)$$

We have $r_{\mathbf{x}} \leq r_{(I+A)^{n-1}\mathbf{x}}$ since $r_{(I+A)^{n-1}\mathbf{x}} = \sup\{\rho \in \mathbb{R} : A(I+A)^{n-1}\mathbf{x} \geq \rho(I+A)^{n-1}\mathbf{x}\}$ and $A\mathbf{x} \geq r_{\mathbf{x}}\mathbf{x} \Rightarrow A(I+A)^{n-1}\mathbf{x} = (I+A)^{n-1}A\mathbf{x} \geq r_{\mathbf{x}}(I+A)^{n-1}\mathbf{x}$. Thus,

$$(*) \leq \sup_{\mathbf{x} \geq \mathbf{0}, \, \|\mathbf{x}\|_2 = 1} r_{(I+A)^{n-1}\mathbf{x}} = \sup_{\mathbf{y} \in Q} r_{\mathbf{y}} = \max_{\mathbf{y} \in Q} r_{\mathbf{y}} = r_{\mathbf{z}} \leq r_{\mathbf{y}}$$

for some $\mathbf{z} \in Q = \{(I + A)^{n-1}\mathbf{x} : \mathbf{x} \ge \mathbf{0}, \|\mathbf{x}\|_2 = 1\}$ ($r_{\mathbf{y}}$ is continuous in \mathbf{y} (why?) and Q is a compact). Note that $\mathbf{z} > \mathbf{0}$.

proof: Let λ be an eigenvalue of A, i.e. $\exists \mathbf{y} \neq \mathbf{0} | A\mathbf{y} = \lambda \mathbf{y}$. Then $|\lambda| |\mathbf{y}| = |\lambda \mathbf{y}| = |A\mathbf{y}| \leq A|\mathbf{y}|, |\mathbf{y}| \geq \mathbf{0}, |\mathbf{y}| \neq \mathbf{0}$. Thus, by definition of $r_{|\mathbf{y}|}, |\lambda| \leq r_{|\mathbf{y}|} \leq r$, and we have the inequality $\rho(A) \leq r$. But r is an eigenvalue of A, thus $r = \rho(A)$.

proof: Let λ be an eigenvalue of B, i.e. $\exists \mathbf{y} \neq \mathbf{0} \mid B\mathbf{y} = \lambda \mathbf{y}$. Then $|\lambda||\mathbf{y}| = |\lambda \mathbf{y}| = |B\mathbf{y}| \leq |B||\mathbf{y}| \leq A|\mathbf{y}|, |\mathbf{y}| \geq \mathbf{0}, |\mathbf{y}| \neq \mathbf{0}$. Thus, by definition of $r_{|\mathbf{y}|}, |\lambda| \leq r_{|\mathbf{y}|} \leq r$, and we have the inequality $\rho(B) \leq r = \rho(A)$.

proof: Assume $\rho(B) = \rho(A)$, i.e. there exists λ eigenvalue of B ($B\mathbf{y} = \lambda \mathbf{y}$, $\mathbf{y} \neq \mathbf{0}$) such that $|\lambda| = r$. Then we can add an equality in the above arguments,

$$r|\mathbf{y}| = |\lambda||\mathbf{y}| = |\lambda\mathbf{y}| = |B\mathbf{y}| \le |B||\mathbf{y}| \le A|\mathbf{y}|, \quad |\mathbf{y}| \ge \mathbf{0}, \ |\mathbf{y}| \ne \mathbf{0},$$

obtaining the inequality $r|\mathbf{y}| \leq A|\mathbf{y}|, |\mathbf{y}| \geq \mathbf{0}, |\mathbf{y}| \neq \mathbf{0}$. But by the above Proposition, such inequality implies $A|\mathbf{y}| = r|\mathbf{y}|$ with $|\mathbf{y}| > \mathbf{0}$. So we have

$$r|\mathbf{y}| = |\lambda||\mathbf{y}| = |\lambda\mathbf{y}| = |B\mathbf{y}| = |B||\mathbf{y}| = A|\mathbf{y}| = r|\mathbf{y}|, \quad |\mathbf{y}| > \mathbf{0},$$

from which it follows |B| = A, against the hypothesis! Thus $\rho(B) < \rho(A)$. \Box

We conclude this section with the following result of the Perron-Frobenius theory, stated without proof (actually we shall give a proof of such result in the case A is positive):

Result. If A is a non-negative irreducible $n \times n$ matrix, then $r = \rho(A)$ is a simple eigenvalue of A, i.e. $\lambda - r$ divides $p_A(\lambda)$ but $(\lambda - r)^2$ does not divide $p_A(\lambda)$.

Computing the Perron-pair of A irreducible non-negative by the power method

Let A be an irreducible non-negative $n \times n$ matrix. Then we know that $\rho(A)$ is positive and is a simple eigenvalue of A, and the corresponding eigenvector can be chosen positive. Of course, such eigenvector is uniquely defined if we require that its measure in the 1-norm is one. We also know that $\rho(A) = r := \sup_{\mathbf{x} \ge \mathbf{0}, \mathbf{x} \neq \mathbf{0}} \min_{i: x_i > 0} \frac{(A\mathbf{x})_i}{x_i}$.

So, it is uniquely defined the Perron-pair $(\rho(A), \mathbf{z})$. such that $A\mathbf{z} = \rho(A)\mathbf{z}$, $\rho(A)$ positive, \mathbf{z} positive, $\|\mathbf{z}\|_1 = 1$. The power method is a way to compute such Perron-pair.

Theorem(power). Let A be an irreducible non-negative $n \times n$ matrix. Let \mathbf{a}_0 be any positive vector, and set $\mathbf{v}_0 = \mathbf{a}_0 / \|\mathbf{a}_0\|_1$. Then set

$$\mathbf{a}_{k+1} = A\mathbf{v}_k, \quad \mathbf{v}_{k+1} = \frac{\mathbf{a}_{k+1}}{\|\mathbf{a}_{k+1}\|_1}, \quad k = 0, 1, 2, \dots$$

Note that the sequences $\{\mathbf{a}_k\}, \{\mathbf{v}_k\}$ are well defined sequences of positive vectors, and $\|\mathbf{v}_k\|_1 = 1, \forall k$. Let X be the non singular matrix defining, by similarity, the Jordan block-diagonal form of A, i.e.

$$X^{-1}AX = J = \begin{bmatrix} r & \mathbf{0}^T \\ \mathbf{0} & B \end{bmatrix}, \quad X = \begin{bmatrix} \mathbf{z} & \mathbf{x}_2 & \cdots & \mathbf{x}_n \end{bmatrix},$$

and let $r = \rho(A)$, λ_j , j = 2, ..., n, be the eigenvalues of A ($\lambda_j \neq r, \forall j$). If α in the expression $\mathbf{a}_0 = \alpha \mathbf{z} + \sum_{j=2}^n \alpha_j \mathbf{x}_j$ is nonzero, then, for $k \to +\infty$, we have

$$\mathbf{v}_k - \mathbf{z} \rightarrow \mathbf{0}, \quad \|\mathbf{a}_k\|_1 - \rho(A) \rightarrow 0,$$

provided that $|\lambda_j|$ is smaller than r for all j = 2, ..., n. In the particular case where A is diagonalizable, the rate of convergence is

$$\Big(\max_{j=2\dots n}\frac{|\lambda_j|}{r}\Big)^k$$

[for the general case, use the Remark in Theorem(power) of the section on preliminary considerations].

proof: We prove the Theorem only in the case A diagonalizable, where $A\mathbf{x}_j = \lambda_j \mathbf{x}_j$, j = 2, ..., n. It is easy to observe that $A^k \mathbf{a}_0 = \alpha r^k \mathbf{z} + \sum_{j=2}^n \alpha_j \lambda_j^k \mathbf{x}_j$ is a positive vector, and that

$$\mathbf{v}_{k} = \frac{\alpha r^{k} \mathbf{z} + \sum_{j=2}^{n} \alpha_{j} \lambda_{j}^{k} \mathbf{x}_{j}}{\|\alpha r^{k} \mathbf{z} + \sum_{j=2}^{n} \alpha_{j} \lambda_{j}^{k} \mathbf{x}_{j}\|_{1}} = \frac{\alpha r^{k} \mathbf{z} + \sum_{j=2}^{n} \alpha_{j} \lambda_{j}^{k} \mathbf{x}_{j}}{\mathbf{e}^{T} (\alpha r^{k} \mathbf{z} + \sum_{j=2}^{n} \alpha_{j} \lambda_{j}^{k} \mathbf{x}_{j})}$$
$$= \frac{\alpha r^{k} \mathbf{z} + \sum_{j=2}^{n} \alpha_{j} \lambda_{j}^{k} \mathbf{x}_{j}}{\alpha r^{k} \mathbf{e}^{T} \mathbf{z} + \sum_{j=2}^{n} \alpha_{j} \lambda_{j}^{k} \mathbf{e}^{T} \mathbf{x}_{j}} = \frac{\mathbf{z} + \sum_{j=2}^{n} \frac{\alpha_{j}}{\alpha} \left(\frac{\lambda_{j}}{r}\right)^{k} \mathbf{x}_{j}}{1 + \sum_{j=2}^{n} \frac{\alpha_{j}}{\alpha} \left(\frac{\lambda_{j}}{r}\right)^{k} \mathbf{e}^{T} \mathbf{x}_{j}}$$

Moreover,

$$\mathbf{a}_{k+1} = A\mathbf{v}_k = \frac{r\mathbf{z} + \sum_{j=2}^n \frac{\alpha_j}{\alpha} \left(\frac{\lambda_j}{r}\right)^k \lambda_j \mathbf{x}_j}{1 + \sum_{j=2}^n \frac{\alpha_j}{\alpha} \left(\frac{\lambda_j}{r}\right)^k \mathbf{e}^T \mathbf{x}_j}$$

and, since \mathbf{a}_{k+1} is positive,

$$\|\mathbf{a}_{k+1}\|_{1} = \mathbf{e}^{T}\mathbf{a}_{k+1} = \frac{r + \sum_{j=2}^{n} \frac{\alpha_{j}}{\alpha} \left(\frac{\lambda_{j}}{r}\right)^{k} \lambda_{j} \mathbf{e}^{T} \mathbf{x}_{j}}{1 + \sum_{j=2}^{n} \frac{\alpha_{j}}{\alpha} \left(\frac{\lambda_{j}}{r}\right)^{k} \mathbf{e}^{T} \mathbf{x}_{j}}$$
$$= \frac{r \left(1 + \sum_{j=2}^{n} \frac{\alpha_{j}}{\alpha} \left(\frac{\lambda_{j}}{r}\right)^{k+1} \mathbf{e}^{T} \mathbf{x}_{j}\right)}{1 + \sum_{j=2}^{n} \frac{\alpha_{j}}{\alpha} \left(\frac{\lambda_{j}}{r}\right)^{k} \mathbf{e}^{T} \mathbf{x}_{j}}$$

Exercise. Discuss the convergence of the power method when applied to compute the Perron-pair $\left(1, \begin{bmatrix} \frac{1}{2} \\ \frac{1}{2} \end{bmatrix}\right)$ of the following two 2 × 2 matrices:

$$A = \begin{bmatrix} 0 & 1\\ 1 & 0 \end{bmatrix}, \quad A = \begin{bmatrix} 1/4 & 3/4\\ 3/4 & 1/4 \end{bmatrix}$$

In the second case, find a constant c such that $\|\mathbf{v}_k - \mathbf{z}\| \leq c(\frac{1}{2})^k$.

Exercise. Prove the Theorem in case A (non negative, irreducible) is 4×4 , and there exists X non singular for which

$$X^{-1}AX = \begin{bmatrix} r & & \\ & \lambda_2 & 1 & \\ & & \lambda_2 & 1 \\ & & & \lambda_2 \end{bmatrix},$$

with $|\lambda_2|$ smaller than r. Prove that in such case the rate of convergence is

$$|p_2(k)| \Big(\left| \frac{\lambda_2}{r} \right| \Big)^k,$$

where p_2 is a degree-two polynomial.

Exercise. Set

$$A = \begin{bmatrix} 0 & 0 & 1\\ \frac{1}{4} & 0 & 0\\ \frac{3}{4} & 1 & 0 \end{bmatrix}.$$

Prove that the eigenvalues of A are $\{1, -\frac{1}{2}, -\frac{1}{2}\}$ and that A is not diagonalizable. Find $X = [\mathbf{z} \ \mathbf{x}_2 \ \mathbf{x}_3]$ such that

$$X^{-1}AX = \begin{bmatrix} 1 & 0 & 0\\ 0 & -\frac{1}{2} & 1\\ 0 & 0 & -\frac{1}{2} \end{bmatrix}.$$

The only hypothesis A irreducible, non-negative, does not assure that $r = \rho(A)$ is the unique eigenvalue of A whose absolute value is equal to r. We only know that if $|\lambda_j| = r, j \in \{2, ..., n\}$, then $\lambda_j \neq r$. Let us see examples:

$$A = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}; \text{ Eigenvalues: } -1,1; \text{ Perron-pair: } (1, \begin{bmatrix} \frac{1}{2} \\ \frac{1}{2} \end{bmatrix}).$$
$$A = \begin{bmatrix} 0 & a & 0 \\ 1 & 0 & 1 \\ 0 & 1-a & 0 \end{bmatrix}, a \in (0,1); \text{ Eig: } -1,0,1; \text{ Perron-pair: } (1, \begin{bmatrix} \frac{a}{2} \\ \frac{1}{2} \\ \frac{1-a}{2} \end{bmatrix}).$$
$$A = \begin{bmatrix} 2 & 1 & 1 \\ 1 & 2 & 1 \\ 1 & 1 & 2 \end{bmatrix}; \text{ Eigenvalues: } 1,1,4; \text{ Perron-pair: } (4, \begin{bmatrix} 1/3 \\ 1/3 \\ 1/3 \end{bmatrix}).$$

Note that if in the second example A is replaced by A^T , then there is no need of computation in obtaining the perron-pair, since it is clear that $A^T \mathbf{e} = \mathbf{e}$, $\mathbf{e} = [1 \ 1 \ 1]^T$. Moreover, in the third example $r = \rho(A)$ (which is 4) dominates the remaining eigevalues of A (which are 1, 1). We note that this fact is true for any *positive* matrix A (see Theorem(positive) below), even if, as the following example shows, it is not a peculiarity of positive matrices:

$$A = \begin{bmatrix} 2 & 1 & 0 \\ 1 & 2 & 1 \\ 0 & 1 & 2 \end{bmatrix}; \text{ Eig: } 2 \mp \sqrt{2}, 2; \text{ Perron-pair: } (2 + \sqrt{2}, \frac{1}{2 + \sqrt{2}} \begin{bmatrix} 1 \\ \sqrt{2} \\ 1 \end{bmatrix})$$

(here computation is required to obtain the Perron-pair).

Theorem (positive). If A is a $n \times n$ positive matrix and $r = \rho(A), \lambda_2, \ldots, \lambda_n$ are its eigenvalues, then $|\lambda_j|$ is smaller than r for all $j = 2, \ldots, n$.

proof: Set $W = A - s\mathbf{z}\mathbf{e}^T$. The eigenvalues of W are $r - s, \lambda_2, \ldots, \lambda_n$ (a sketch of the proof:

$$Y := \begin{bmatrix} \mathbf{z} & \mathbf{y}_2 & \cdots & \mathbf{y}_n \end{bmatrix} \text{ non singular,} \\ Y^{-1}AY = \begin{bmatrix} r & \cdots \\ \mathbf{0} & M \end{bmatrix}, \ Y^{-1}WY = \begin{bmatrix} r-s & \cdots \\ \mathbf{0} & M \end{bmatrix}).$$

If there is a value of s for which $|W| \leq A$, $|W| \neq A$, then $\rho(W)$ is smaller than $\rho(A)$, and thus $|\lambda_2|, \ldots, |\lambda_n|$ are smaller than $r = \rho(A)$. If A is positive, then such s exists, $s = \min_{i,j} a_{ij}$. \Box

Irreducible non-negative stochastic-by-columns A and power method

Let A be an irreducible non-negative stochastic by columns $n \times n$ matrix. Then we know that $1 = \rho(A)$ $(A^T \mathbf{e} = \mathbf{e}, \rho(A) \leq ||A||_1 = 1)$, that $r = 1 = \rho(A)$ is a simple eigenvalue of A, and that the corresponding eigenvector can be chosen positive. Of course, such eigenvector is uniquely defined if we require that its measure in the 1-norm is one.

So, it is uniquely defined the Perron-pair $(1 = \rho(A), \mathbf{z})$. such that $A\mathbf{z} = \mathbf{z}, \mathbf{z}$ positive, $\|\mathbf{z}\|_1 = 1$. The computation of such Perron-pair, i.e. the computation of the vector \mathbf{z} , can be performed via the power method. Actually, we now see that the power method in the particular case where A is stochastic-by-columns (besides non-negative and irreducible) can be rewritten in a simpler form and converges independently from the choice of \mathbf{a}_0 (provided that 1 dominates the other eigenvalues). These results follow from some remarks, reported in the following Proposition.

Proposition. Let A be an irreducible non-negative stochastic-by-columns $n\times n$ matrix. Then

i)

$$\mathbf{v} \in \mathbb{C}^n \quad \Rightarrow \quad \sum_i (A\mathbf{v})_i = \sum_i v_i$$
$$(\sum_i (A\mathbf{v})_i = \sum_i \sum_j a_{ij}v_j = \sum_j v_j \sum_i a_{ij} = \sum_j v_j),$$
ii)

v positive, $\|\mathbf{v}\|_1 = 1 \Rightarrow A\mathbf{v}$ positive, $\|A\mathbf{v}\|_1 = 1$

(use the irreducibility of A and assertion i)), iii)

$$A\mathbf{v} = \lambda \mathbf{v}, \ \lambda \neq 1, \ \Rightarrow \ \sum_i v_i = 0$$

$$(\sum_{i} v_i = \sum_{i} (A\mathbf{v})_i = \lambda \sum_{i} v_i, \text{ thus } (\lambda - 1) \sum_{i} v_i = 0).$$

Exercise. Assume that $X^{-1}AX = J$ where J is the Jordan block diagonal form of A, where A is an irreducible non-negative stochastic-by-columns $n \times n$ matrix. Assume that $[J]_{11} = 1$. Prove that $\sum_i [X]_{ij} = 0, j = 2, \ldots, n$.

Corollary(power). Let A be an irreducible non-negative stochastic-by-columns $n \times n$ matrix. Let \mathbf{a}_0 be any positive vector such that $\|\mathbf{a}_0\|_1 = 1$. Then set

$$\mathbf{a}_{k+1} = A\mathbf{a}_k, \quad k = 0, 1, 2, \dots$$

The sequence $\{\mathbf{a}_k\}$ is a well defined sequence of positive vectors, and $\|\mathbf{a}_k\|_1 = 1 = \rho(A), \forall k$. If $k \to +\infty$, then

$$\mathbf{a}_k - \mathbf{z} \rightarrow \mathbf{0},$$

provided that $|\lambda_j|$ is smaller than 1 for all j = 2, ..., n. The latter event is assured by Theorem(positive) when A is positive. In the particular case where A is diagonalizable, the rate of convergence is

$$(\max_{j=2\dots n} |\lambda_j|)^k$$

[for the general case, use the Remark in Theorem(power) of the section on preliminary considerations].

proof: The sequences \mathbf{a}_k and \mathbf{v}_k defined in Theorem(power) coincide, because, by Proposition ii), we have $||A\mathbf{v}_k||_1 = ||\mathbf{v}_k||_1$. It remains to show that α in the expression $\mathbf{a}_0 = \alpha \mathbf{z} + \sum_j \alpha_j \mathbf{x}_j$ is nonzero. Note that $\mathbf{e}^T \mathbf{a}_0 = \alpha \mathbf{e}^T \mathbf{z} + \sum_j \alpha_j \mathbf{e}^T \mathbf{x}_j = \alpha + \sum_j \alpha_j \mathbf{e}^T \mathbf{x}_j$. In the case that A is diagonalizable the thesis follows by Proposition iii) which asserts that $\mathbf{e}^T \mathbf{x}_j$, $j = 2, \ldots, n$, must be zero. In the generic case the proof is left to the reader \ldots

Exercise. Discuss existence, unicity, and computation of the Perron-pair for

$$A = \begin{bmatrix} 0 & 1-a & a \\ a & 0 & 1-a \\ 1-a & a & 0 \end{bmatrix}, \ a \in [0,1], \ A = \begin{bmatrix} 0 & 0 & 1 \\ a & 0 & 0 \\ 1-a & 1 & 0 \end{bmatrix}, \ a \in [0,1]$$
$$A = \begin{bmatrix} 0 & b_1 & & & \\ 1 & 0 & b_2 & & \\ & 1-b_1 & 0 & \ddots & \\ & & 1-b_2 & \ddots & b_{n-2} \\ & & & \ddots & 0 & 1 \\ & & & 1-b_{n-2} & 0 \end{bmatrix}, \ b_i \in (0,1).$$

Page-rank computation [Berkhin, et al]

Consider an oriented graph with set of verteces $V = \{1, 2, ..., n\}$, and set of edges $E, i, j \in E$ if there is a link from i to j.

Associate to such graph the adiacency matrix:

$$L = \left\{ \begin{array}{ll} 1 & ij \in E \\ 0 & ij \notin E \end{array} \right.$$

Call deg(i) the number of edges starting from i. Of course, deg(i) = 0 (no edge starts from i) iff $L_{ij} = 0, \forall j$. Moreover, deg(i) = $\sum_j L_{ij}$.

Associate to the graph the transition matrix:

$$P = \begin{cases} \frac{1}{\deg(i)} & ij \in E\\ 0 & ij \notin E \end{cases}$$

Note that $P_{ij} = L_{ij}/\deg(i)$ if *i* is such that $\deg(i) > 0$, and $P_{ij} = L_{ij} = 0$ otherwise. Moreover, $\sum_j P_{ij} = 1$ if $\deg(i) > 0$ and $\sum_j P_{ij} = 0$ otherwise. So, the matrix *P* is a non negative matrix *quasi*-stochastic by rows.

Remark. Row i of P is null iff no edge starts from i; column j of P is null iff no edge points to j.

Let $\mathbf{p} \in \mathbb{R}^n$ be the vector whose entry j, p_j , is the importance (authority) of the vertex j. Then

$$p_{j} = \sum_{i: i \to j} \frac{p_{i}}{\deg(i)} = \sum_{i=1}^{n} P_{ij} p_{i} = \sum_{i=1}^{n} P_{ji}^{T} p_{i} = (P^{T} \mathbf{p})_{j}, \quad \mathbf{p} = P^{T} \mathbf{p}$$

(note that such fixed point \mathbf{p} may not exist, or, if exists, may be not unique or with zero entries; see below).

Let $p_j^{(k+1)}$ be the probability that at step k+1 of my visit of the graph (navigation on the web) I am on the vertex (page) j. Then

$$p_{j}^{(k+1)} = \sum_{i:i \to j} \frac{p_{i}^{(k)}}{\deg(i)} = \sum_{i=1}^{n} P_{ij} p_{i}^{(k)}$$
$$= \sum_{i=1}^{n} P_{ji}^{T} p_{i}^{(k)} = (P^{T} \mathbf{p}^{(k)})_{j}, \quad \mathbf{p}^{(k+1)} = P^{T} \mathbf{p}^{(k)}$$

(note that such sequence of vector probabilities exists and is uniquely defined, once $\mathbf{p}^{(0)}$ is given, but the $\mathbf{p}^{(k)}$ may loss the possible ddp property of $\mathbf{p}^{(0)}$; see below). [A vector \mathbf{w} is said ddp (discrete distribution of probability) if \mathbf{w} is positive and $\|\mathbf{w}\|_1 = 1$].

Note that it is natural to require that: $p_i^{(k)} > 0$ (at step k there is a probability that I am on vertex i), $\sum_j p_j^{(k)} = 1$ (at step k I am on some vertex); $p_i > 0$ (any vertex has a portion of importance ...), $\sum_i p_i = 1$ (... of the total 1).

So, the following facts must be true:

a) **p** such that $\mathbf{p} = P^T \mathbf{p}, \mathbf{p} > \mathbf{0}, \|\mathbf{p}\|_1 = 1$, exists and is uniquely defined, b) the method $\mathbf{p}^{(0)} = ddp, \mathbf{p}^{(k+1)} = P^T \mathbf{p}^{(k)}$ converges to **p**.

We now show that in order to have a) and b), the matrix P^T must be both stochastic by columns and irreducible.

Theorem(stoc). If the above facts a) and b) are true, then P^T must be stochastic by columns.

proof: We know that $\exists ! \mathbf{p}$ such that $\mathbf{p} = P^T \mathbf{p}, \mathbf{p} > \mathbf{0}, \|\mathbf{p}\|_1 = 1$. Assume that P^T is quasi-stochastic but not stochastic by columns. Then

$$\begin{aligned} \|P^T \mathbf{p}\|_1 &= \sum_i (P^T \mathbf{p})_i = \sum_i \sum_j (P^T)_{ij} p_j = \sum_j \sum_i P_{ji} p_j \\ &= \sum_{j: \deg(j) > 0} 1 \cdot p_j + \sum_{j: \deg(j) = 0} 0 \cdot p_j < \sum_j p_j = \|\mathbf{p}\|_1, \end{aligned}$$

analogously,

$$\|\mathbf{p}^{(k+1)}\|_1 = \|P^T \mathbf{p}^{(k)}\|_1 \le \|\mathbf{p}^{(k)}\|_1 \le \|\mathbf{p}^{(1)}\|_1 < \|\mathbf{p}^{(0)}\|_1 = 1.$$

Thus **p** cannot be equal to P^T **p**, and **p**^(k) cannot converge to a ddp. \Box

Theorem(irred). If the above facts a) and b) are true, then ${\cal P}^T$ must be irreducible.

proof: assume P^T reducible. Then there exists a permutation matrix Q such that

$$Q^T P^T Q = \begin{bmatrix} A_{11} & A_{12} \\ 0 & A_{22} \end{bmatrix}, \qquad (*)$$

with A_{11} and A_{22} at least 1×1 square matrices. Note that $Q^T P^T Q$ is stochastic by columns, like P^T . We know that $\exists ! \mathbf{p}$ such that $\mathbf{p} = P^T \mathbf{p}, \mathbf{p} > \mathbf{0}, \|\mathbf{p}\|_1 = 1$, but this is equivalent to say that $\exists ! Q^T \mathbf{p}$ such that

$$Q^{T}\mathbf{p} = \begin{bmatrix} A_{11} & A_{12} \\ 0 & A_{22} \end{bmatrix} Q^{T}\mathbf{p}, \quad Q^{T}\mathbf{p} > \mathbf{0}, \quad ||Q^{T}\mathbf{p}||_{1} = 1.$$
(**)

Case 1: assume $A_{12} = 0$. Then A_{11} and A_{22} are non negative stochastic by columns matrices. We can assume they are also irreducible (why?). Then, by the Perron-Frobenius theory,

$$\exists ! \mathbf{y}_1, \mathbf{y}_2 > \mathbf{0}, \ \|\mathbf{y}_1\|_1 = \|\mathbf{y}_2\|_1 = 1, \ \mathbf{y}_1 = A_{11}\mathbf{y}_1, \ \mathbf{y}_2 = A_{22}\mathbf{y}_2.$$

and, as a consequence, for all $\alpha \in (0, 1)$ we have

$$\begin{bmatrix} \alpha \mathbf{y}_1 \\ (1-\alpha)\mathbf{y}_2 \\ \alpha \mathbf{y}_1 \\ (1-\alpha)\mathbf{y}_2 \end{bmatrix} = \begin{bmatrix} A_{11} & 0 \\ 0 & A_{22} \end{bmatrix} \begin{bmatrix} \alpha \mathbf{y}_1 \\ (1-\alpha)\mathbf{y}_2 \end{bmatrix},$$
$$\mathbf{0}, \| \begin{bmatrix} \alpha \mathbf{y}_1 \\ (1-\alpha)\mathbf{y}_2 \end{bmatrix} \|_1 = 1.$$

So, all vectors \mathbf{p} such that $Q^T \mathbf{p} = \begin{bmatrix} \alpha \mathbf{y}_1 \\ (1-\alpha)\mathbf{y}_2 \end{bmatrix}$ satisfy the properties $\mathbf{p} > \mathbf{0}$, $\|\mathbf{p}\|_1 = 1$, $\mathbf{p} = P^T \mathbf{p}$, which is against the hypothesis of unicity.

Case 2: assume $A_{12} \neq 0$. Then A_{22} in (*) is not stochastic by columns, thus A_{22} may have no eigenvalue equal to 1, i.e. the equations in (**) involving A_{22} may be verified only if part of **p** is null, against the hypothesis of positiveness of **p**. \Box

Viceversa, we know that if the non negative matrix P^T is irreducible and stochastic by columns, then $1 = \rho(P^T)$ is a simple eigenvalue of P^T and there exists a unique vector \mathbf{p} such that $\mathbf{p} = P^T \mathbf{p}, \mathbf{p} > \mathbf{0}, \|\mathbf{p}\|_1 = 1$. We also know that such hypotheses are not sufficient to assure the convergence (to \mathbf{p}) of the sequence $\mathbf{p}^{(k+1)} = P^T \mathbf{p}^{(k)}, \mathbf{p}^{(0)} > \mathbf{0}, \|\mathbf{p}_0\|_1 = 1$, or, equivalently, to assure that the remaining eigenvalues of P^T have absolute value smaller than 1. We can only say that the sequence $\{\mathbf{p}^{(k)}\}$ is a well defined sequence of ddp.

We have to modify P (the graph) so to make well posed $(\exists !)$ the mathematical problem and to make convergent the algorithm for solving it.

Make P^T stochastic:

$$P' = P + \mathbf{d}\mathbf{v}^{T}, \quad \mathbf{d} = \begin{bmatrix} \delta_{\deg(1),0} \\ \vdots \\ \delta_{\deg(n),0} \end{bmatrix}, \quad \mathbf{v} = \begin{bmatrix} \frac{1}{n} \\ \vdots \\ \frac{1}{n} \end{bmatrix}$$

i.e, where P has null rows P' has the row vector \mathbf{v}^T . The vertex *i* with deg (*i*) = 0 now links to all verteces of the graph. [We discuss the uniform case, but what follows can be repeated for the more general case $\mathbf{v} = ddp$].

Observe that $(P')^T$ is stochastic by columns, $(P')^T \ge 0$, thus 1 is eigenvalue of $(P')^T$ and the other eigenvalues of $(P')^T$, $\lambda'_2, \ldots, \lambda'_n$, are such that $|\lambda'_i| \le 1$.

Make $(P')^T$ irreducible:

$$P'' = cP' + (1-c)\mathbf{e}\mathbf{v}^T, \quad \mathbf{e} = \begin{bmatrix} 1\\ \vdots\\ 1 \end{bmatrix}, \ c \in (0,1).$$

Since

$$\mathbf{e}_{i}^{T}P^{\prime\prime} = \begin{cases} c\mathbf{v}^{T} + (1-c)\mathbf{v}^{T} = \mathbf{v}^{T} & \deg\left(i\right) = 0\\ c[\cdots 0 \ \frac{1}{\deg\left(i\right)} \ 0 \cdots] + (1-c)\mathbf{v}^{T} & \deg\left(i\right) > 0 \end{cases}$$

we are assuming that a visitor of the graph can go from the vertex i to one of its neighborhoods with probability $c/\deg(i) + (1-c)/n$, and with probability (1-c)/n to an arbitrary vertex of the graph. Of course the parameter c must be chosen near to 1, in order to maintain our model faithful to the way the graph (web) is visited.

Observe that $(P'')^T$ is stochastic by columns and positive, therefore, in particular, it is non negative and irreducible. So, we have all we need.

Theorem(page-rank). $1 = \rho((P'')^T)$ is a simple eigenvalue of $(P'')^T$, there exists a unique vector \mathbf{p} such that $\mathbf{p} = (P'')^T \mathbf{p}, \mathbf{p} > \mathbf{0}, \|\mathbf{p}\|_1 = 1$ (i.e. we have fact (a)), and the other eigenvalues of $(P'')^T, \lambda''_2, \ldots, \lambda''_n$, are such that $|\lambda''_j| < 1$ (by Theorem(positive)). Thus, $\mathbf{p}^{(k+1)} = (P'')^T \mathbf{p}^{(k)}, \mathbf{p}^{(0)} > \mathbf{0}, \|\mathbf{p}^{(0)}\|_1 = 1$, is a sequence of ddp convergent to \mathbf{p} and, in case P'' is diagonalizable,

$$\|\mathbf{p}^{(k)} - \mathbf{p}\| = O((\max_{j=2,\dots,n} |\lambda_j''|)^k)$$

[for the general case see the Remark in Theorem(power) of Section 1] (i.e. we have fact (b)). Moreover, for the particular choice of $(P'')^T$, the cost of each step of the power method is O(n) and is dominated by the cost of the matrix-vector multiplication $P^T \mathbf{z}$, and if A is diagonalizable, then the rate of convergence is

$$\|\mathbf{p}^{(k)} - \mathbf{p}\| = O(c^k)$$

[for the general case see the Remark in Theorem(power) of Section 1] (Google-search engine sets c = 0.85 [Berkhin]).

proof: We have to prove only the final assertions.

O(n) arithmetic operations are sufficient to perform each step of the power method. We have

$$(P'')^T = c(P^T + \mathbf{vd}^T) + (1 - c)\mathbf{ve}^T,$$

and, if $\mathbf{x} \ge \mathbf{0}$, then

$$(P'')^T \mathbf{x} = cP^T \mathbf{x} + \gamma \mathbf{v},$$

$$\gamma = c\mathbf{d}^T \mathbf{x} + (1-c)\mathbf{e}^T \mathbf{x} = \mathbf{e}^T \mathbf{x} - c[\mathbf{e}^T \mathbf{x} - \mathbf{d}^T \mathbf{x}] = \|\mathbf{x}\|_1 - c\|P^T \mathbf{x}\|_1$$

(why the latter equality holds?). Thus, in order to compute $\mathbf{p}^{(k+1)}$ from $\mathbf{p}^{(k)}$ one can use the following function

$$\mathbf{y} = cP^T \mathbf{x}, \gamma = \|\mathbf{x}\|_1 - \|\mathbf{y}\|_1, (P'')^T \mathbf{x} = \mathbf{y} + \gamma \mathbf{v}$$

where the dominant operation is the matrix-vector multiplication $P^T \mathbf{x}$. Note that each row j of $P^{\hat{T}}$ has (on the average) a very small number of nonzero entries, i.e. exactly the number of vertices pointing to j, so $P^T \mathbf{x}$ can be computed with O(n) arithmetic operations. Note also that in order to implement the above function one needs only 2n memory allocations.

Rate of convergence $\|\mathbf{p}^{(k)} - \mathbf{p}\| = O(c^k)$. We first prove that if $\mathbf{e}^T \mathbf{v} = 1$ then

$$p_{P'}(\lambda) = (\lambda - 1)p_{n-1}(\lambda) \quad \Rightarrow \quad p_{P' + \frac{1-c}{c} \mathbf{e} \mathbf{v}^T}(\lambda) = (\lambda - \frac{1}{c})p_{n-1}(\lambda). \quad (* * *)$$

In fact,

$$S = \begin{bmatrix} \mathbf{e} & \mathbf{y}_{2} & \cdots & \mathbf{y}_{n} \end{bmatrix}, \ \det(S) \neq 0,$$
$$S^{-1}P'S = \begin{bmatrix} 1 & \mathbf{u}^{T} \\ \mathbf{0} & M \end{bmatrix}, \ p_{P'}(\lambda) = (\lambda - 1)p_{M}(\lambda),$$
$$S^{-1}(P' + \frac{1-c}{c}\mathbf{e}\mathbf{v}^{T})S = \begin{bmatrix} 1 & \mathbf{u}^{T} \\ \mathbf{0} & M \end{bmatrix} + \frac{1-c}{c}S^{-1}\mathbf{e}\mathbf{v}^{T}S$$
$$= \begin{bmatrix} 1 & \mathbf{u}^{T} \\ \mathbf{0} & M \end{bmatrix} + \frac{1-c}{c}\begin{bmatrix} 1 & \cdots & * \\ 0 & \cdots & 0 \end{bmatrix} = \begin{bmatrix} \frac{1}{c} & \tilde{\mathbf{u}}^{T} \\ \mathbf{0} & M \end{bmatrix}$$

As a consequence of (***), if $1, \lambda'_2, \ldots, \lambda'_n$ are the eigenvalues of P' (recall that $|\lambda'_j| \leq 1, j = 2, \ldots, n$), then the eigenvalues of $P' + \frac{1-c}{c} \mathbf{e} \mathbf{v}^T$ are $\frac{1}{c}, \lambda'_2, \dots, \lambda'_n, \text{ and thus the eigenvalues of } cP' + (1-c)\mathbf{e}\mathbf{v}^T \text{ are } 1, c\lambda'_2, \dots, c\lambda'_n.$ It follows that if A is diagonalizable, then $\|\mathbf{p}^{(k)} - \mathbf{p}\| = O((\max_{j=2,\dots,n} |c\lambda'_j|)^k) =$

 $O(c^k)$. \Box

Why to compute **p**? [Berkhin].

QUERY: Berkhin survey.

Go in the inverted terms document file, which is a table containing a row for each term of a collection's dictionary. In such file, for each term there is a list of all documents that contain such term

$$\begin{array}{l} \vdots \\ \text{term} \quad \rightarrow \quad \text{LISTA}_{term} = \{i_1, i_2, \dots, i_k\} \subset \{1, 2, \dots, n\} \\ \vdots \\ \text{Berkhin} \quad \rightarrow \quad \text{LISTA}_{Berkhin} = \{1, 4, 6\} \\ \vdots \\ \text{survey} \quad \rightarrow \quad \text{LISTA}_{survey} = \{1, 3\} \\ \vdots \end{array}$$

Define the set of relevance of the query

$$\cup_{term \in \text{QUERY}} \text{LISTA}_{term} = \{1, 3, 4, 6\}$$

Reading \mathbf{p} (\mathbf{p} is updated once each month) considers and orders the corresponding set of authorities $\{p_1, p_3, p_4, p_6\}$, for example $p_4 \ge p_6 \ge p_3 \ge p_1$

Finally, show the titles of the documents 1, 3, 4, 6 in the order 4, 6, 3, 1, from the one with greatest authority to the one with smallest authority.

Main criticism: This procedure, being independent from the query allows a fast answer, but does not make distinction between pages with authority from pages with authority on a specific subject.

Exercise. Draw the graph whose transition matrix is

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix}.$$

Note that P is non negative, reducible and quasi-stochastic, but not stochastic, by rows. Prove that there is no positive vector \mathbf{p} such that $\mathbf{p} = P^T \mathbf{p}$. Starting from P and proceeding as indicated in the theory, introduce a non negative matrix P' stochastic by rows. Note that P' is reducible, like P. Prove that there is no positive vector \mathbf{p} such that $\mathbf{p} = (P')^T \mathbf{p}$. Starting from P' and proceeding as indicated in the theory, introduce a non negative matrix P'' irreducible and stochastic by rows. Prove that there exists a unique positive vector \mathbf{p} such that $\mathbf{p} = (P'')^T \mathbf{p}$, $\|\mathbf{p}\|_1 = 1$, and describe an algorithm for the computation of \mathbf{p} . Here below is an approximation of such vector \mathbf{p} :

 $\mathbf{p}^{T} = [0.03721 \ 0.05396 \ 0.04151 \ 0.3751 \ 0.206 \ 0.2862].$

Assume, for example, that the set of relevance of a query is $\{1, 3, 4, 6\}$. Then the documents 1, 3, 4, 6 are listed in the order 4, 6, 3, 1, being $p_4 \ge p_6 \ge p_3 \ge p_1$.

APPENDIX

Proof (Theorem(inverse power)):

Assume A diagonalizable, $A\mathbf{x}_j = \lambda_j \mathbf{x}_j$, $\{\mathbf{x}_j\}_{j=1}^n$ linearly independent. Then

$$A\mathbf{x}_j - \lambda_i^* \mathbf{x}_j = (\lambda_j - \lambda_i^*) \mathbf{x}_j \quad \Rightarrow \quad (A - \lambda_i^* I)^{-m} \mathbf{x}_j = \frac{1}{(\lambda_j - \lambda_i^*)^m} \mathbf{x}_j.$$

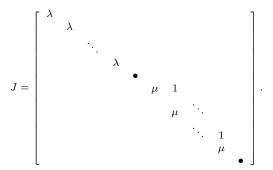
Call α_j the numbers for which $\mathbf{v}_0 = \sum_{j: \lambda_j = \lambda_i} \alpha_j \mathbf{x}_j + \sum_{j: \lambda_j \neq \lambda_i} \alpha_j \mathbf{x}_j$. Note that the first sum is a non null vector (by the assumption on \mathbf{v}_0). Then the following equality holds

$$(\lambda_i - \lambda_i^*)^m (A - \lambda_i^* I)^{-m} \mathbf{v}_0 = \sum_{j: \lambda_j = \lambda_i} \alpha_j \mathbf{x}_j + \sum_{j: \lambda_j \neq \lambda_i} \alpha_j \Big(\frac{\lambda_i - \lambda_i^*}{\lambda_j - \lambda_i^*}\Big)^m \mathbf{x}_j$$

which, for $m \to +\infty$, yields the thesis.

Assume that A is not diagonalizable. In this case call \mathbf{x}_j , $j = 1, \ldots, n$, the columns of the non singular matrix X for which $X^{-1}AX = J$ where J is the block diagonal Jordan form of A. Assume that the upper-left diagonal block of J is diagonal and its diagonal contains all λ_j such that $\lambda_j = \lambda := \lambda_i$. Assume that there are t of such eigenvalues. Then consider any other diagonal block of J; of course it corresponds to an eigenvalue λ_j different from λ , call such

eigenvalue μ . The block under consideration has an order, say $s = s_{\mu}$.



Restrict the matrix equation $X^{-1}AX = J$ to this block, thus, for some $r \ge t$ we have

$$A[\mathbf{x}_{r+1} \ \mathbf{x}_{r+2} \ \mathbf{x}_{r+3} \ \cdots \ \mathbf{x}_{r+s}] = [\mathbf{x}_{r+1} \ \mathbf{x}_{r+2} \ \mathbf{x}_{r+3} \ \cdots \ \mathbf{x}_{r+s}] \begin{bmatrix} \mu & 1 & & \\ & \mu & \ddots & \\ & & \ddots & 1 \\ & & & \mu \end{bmatrix}.$$

It follows that

$$\begin{split} &A\mathbf{x}_{r+1} = \mu \mathbf{x}_{r+1}, \ (A - \lambda^* I)\mathbf{x}_{r+1} = (\mu - \lambda^*)\mathbf{x}_{r+1}, \\ &(A - \lambda^* I)^{-1}\mathbf{x}_{r+1} = \frac{1}{\mu - \lambda^*}\mathbf{x}_{r+1}, \ (A - \lambda^* I)^{-m}\mathbf{x}_{r+1} = \frac{1}{(\mu - \lambda^*)^m}\mathbf{x}_{r+1}. \\ &A\mathbf{x}_{r+2} = \mu \mathbf{x}_{r+2} + \mathbf{x}_{r+1}, \ (A - \lambda^* I)\mathbf{x}_{r+2} = (\mu - \lambda^*)\mathbf{x}_{r+2} + \mathbf{x}_{r+1}, \\ &(A - \lambda^* I)^{-1}\mathbf{x}_{r+2} = \frac{1}{\mu - \lambda^*}\mathbf{x}_{r+2} - \frac{1}{(\mu - \lambda^*)^2}\mathbf{x}_{r+1}, \\ &(A - \lambda^* I)^{-m}\mathbf{x}_{r+2} = \frac{1}{(\mu - \lambda^*)^m}\mathbf{x}_{r+2} - \frac{m}{(\mu - \lambda^*)^m}\mathbf{x}_{r+1}. \\ \\ &A\mathbf{x}_{r+3} = \mu \mathbf{x}_{r+3} + \mathbf{x}_{r+2}, \ (A - \lambda^* I)\mathbf{x}_{r+3} = (\mu - \lambda^*)\mathbf{x}_{r+3} + \mathbf{x}_{r+2}, \\ &(A - \lambda^* I)^{-m}\mathbf{x}_{r+3} = \frac{1}{\mu - \lambda^*}\mathbf{x}_{r+3} - \frac{1}{(\mu - \lambda^*)^2}\mathbf{x}_{r+2} + \frac{1}{(\mu - \lambda^*)^3}\mathbf{x}_{r+1}, \\ &(A - \lambda^* I)^{-m}\mathbf{x}_{r+3} = \frac{1}{(\mu - \lambda^*)^m}\mathbf{x}_{r+3} - \frac{m}{(\mu - \lambda^*)^{m+1}}\mathbf{x}_{r+2} + \frac{1}{(\mu - \lambda^*)^{m+2}}\mathbf{x}_{r+1}. \\ \\ &A\mathbf{x}_{r+4} = \mu \mathbf{x}_{r+4} + \mathbf{x}_{r+3}, \ (A - \lambda^* I)\mathbf{x}_{r+4} = (\mu - \lambda^*)\mathbf{x}_{r+4} + \mathbf{x}_{r+3}, \\ &(A - \lambda^* I)^{-m}\mathbf{x}_{r+4} = \frac{1}{\mu - \lambda^*}\mathbf{x}_{r+4} - \frac{1}{(\mu - \lambda^*)^2}\mathbf{x}_{r+3} \\ &+ \frac{1}{(\mu - \lambda^*)^m}\mathbf{x}_{r+2} - \frac{1}{(\mu - \lambda^*)^{m}}\mathbf{x}_{r+1}, \\ &(A - \lambda^* I)^{-m}\mathbf{x}_{r+4} = \frac{1}{(\mu - \lambda^*)^m}\mathbf{x}_{r+4} - \frac{m}{(\mu - \lambda^*)^{m+1}}\mathbf{x}_{r+3} \\ &+ \frac{1}{(\mu - \lambda^*)^{m+2}}\mathbf{x}_{r+2} - \frac{1}{m}\frac{m^3 + \frac{1}{2}m^2 + \frac{1}{3}m}{(\mu - \lambda^*)^{m+3}}\mathbf{x}_{r+1}, \\ &(A - \lambda^* I)^{-m}\mathbf{x}_{r+8} = \frac{1}{(\mu - \lambda^*)^m}\mathbf{x}_{r+8} - \dots + (-1)^{s-1}\frac{q_{s-1}(m)}{(\mu - \lambda^*)^{m+s-1}}\mathbf{x}_{r+1}, \\ &(A - \lambda^* I)^{-m}[\mathbf{x}_{r+1}\mathbf{x}_{r+2}\mathbf{x}_{r+3}\cdots] \\ &= [\mathbf{x}_{r+1}\mathbf{x}_{r+2}\mathbf{x}_{r+3}\cdots] \begin{bmatrix} \frac{1}{(\mu - \lambda^*)^m} - \frac{m}{(\mu - \lambda^*)^{m+1}} & \frac{1}{(\mu - \lambda^*)^m} - \frac{1}{(\mu - \lambda^*)^m} \\ &0 & 0 & \frac{1}{(\mu - \lambda^*)^m} \end{bmatrix} \\ &\vdots & \ddots \end{bmatrix}$$

Set also $\lambda^* := \lambda_i^* (\lambda_i^* \text{ is the given approximation of } \lambda = \lambda_i)$. Choose \mathbf{v}_0 , and call α_j the numbers for which $\mathbf{v}_0 = \sum_{j:\lambda_j=\lambda} \alpha_j \mathbf{x}_j + \{ \dots + \sum_{j=r+1}^{r+s} \alpha_j \mathbf{x}_j + \dots \}$. Note that $\sum_{j:\lambda_j=\lambda} \alpha_j \mathbf{x}_j = \sum_{j=1}^t \alpha_j \mathbf{x}_j$. Then

$$\begin{aligned} &(A - \lambda^* I)^{-m} \mathbf{v}_0 \\ &= \sum_{j: \, \lambda_j = \lambda} \alpha_j (A - \lambda^* I)^{-m} \mathbf{x}_j + \{ \cdots + \sum_{j=r+1}^{r+s} \alpha_j (A - \lambda^* I)^{-m} \mathbf{x}_j + \cdots \} \\ &= \sum_{j: \, \lambda_j = \lambda} \alpha_j \frac{1}{(\lambda - \lambda^*)^m} \mathbf{x}_j + \{ \cdots + [\alpha_{r+1} (\frac{1}{(\mu - \lambda^*)^m} \mathbf{x}_{r+1}) \\ &+ \alpha_{r+2} (\frac{1}{(\mu - \lambda^*)^m} \mathbf{x}_{r+2} - \frac{m}{(\mu - \lambda^*)^{m+1}} \mathbf{x}_{r+1}) \\ &+ \alpha_{r+3} (\frac{1}{(\mu - \lambda^*)^m} \mathbf{x}_{r+3} - \frac{m}{(\mu - \lambda^*)^{m+1}} \mathbf{x}_{r+2} + \frac{\frac{1}{2} (m^2 + m)}{(\mu - \lambda^*)^{m+2}} \mathbf{x}_{r+1}) \\ &+ \alpha_{r+4} (\frac{1}{(\mu - \lambda^*)^m} \mathbf{x}_{r+4} - \frac{m}{(\mu - \lambda^*)^{m+1}} \mathbf{x}_{r+3} \\ &+ \frac{\frac{1}{2} (m^2 + m)}{(\mu - \lambda^*)^{m+2}} \mathbf{x}_{r+2} - \frac{\frac{1}{6} m^3 + \frac{1}{2} m^2 + \frac{1}{3} m}{(\mu - \lambda^*)^{m+3}} \mathbf{x}_{r+1}) \\ &+ \ldots + \alpha_{r+s} (\frac{1}{(\mu - \lambda^*)^m} \mathbf{x}_{r+s} - \ldots + (-1)^{s-1} \frac{q_{s-1}(m)}{(\mu - \lambda^*)^{m+s-1}} \mathbf{x}_{r+1})] + \cdots \}. \end{aligned}$$

But this implies

$$\begin{aligned} &(\lambda-\lambda^*)^m (A-\lambda^*I)^{-m} \mathbf{v}_0 \\ &= \sum_{j:\,\lambda_j=\lambda} \alpha_j \mathbf{x}_j + \{ \dots + \left(\frac{\lambda-\lambda^*}{\mu-\lambda^*}\right)^m \\ &\left[\left(\alpha_{r+1} - \alpha_{r+2} \frac{m}{\mu-\lambda^*} + \alpha_{r+3} \frac{\frac{1}{2}(m^2+m)}{(\mu-\lambda^*)^2} \dots + (-1)^{s-1} \alpha_{r+s} \frac{q_{s-1}(m)}{(\mu-\lambda^*)^{s-1}} \right) \mathbf{x}_{r+1} \\ &+ \left(\alpha_{r+2} - \alpha_{r+3} \frac{m}{\mu-\lambda^*} + \alpha_{r+4} \frac{\frac{1}{2}(m^2+m)}{(\mu-\lambda^*)^2} \dots + (-1)^{s-2} \alpha_{r+s} \frac{q_{s-2}(m)}{(\mu-\lambda^*)^{s-2}} \right) \mathbf{x}_{r+2} \\ &+ \dots + \left(\alpha_{r+s}\right) \mathbf{x}_{r+s} \right] + \cdots \}, \end{aligned}$$

from which it is clear that the sequence $(\lambda - \lambda^*)^m (A - \lambda^* I)^{-m} \mathbf{v}_0$ converges to an eigenvector of A associated with λ . The assertions about the rate of convergence follow by setting

$$p_{s-1}(\lambda) = \alpha_{r+1} - \alpha_{r+2} \frac{m}{\mu - \lambda^*} + \alpha_{r+3} \frac{\frac{1}{2}(m^2 + m)}{(\mu - \lambda^*)^2} - \alpha_{r+4} \frac{\frac{1}{6}m^3 + \frac{1}{2}m^2 + \frac{1}{3}m}{(\mu - \lambda^*)^3} \dots + (-1)^{s-1} \alpha_{r+s} \frac{q_{s-1}(m)}{(\mu - \lambda^*)^{s-1}}.$$

Proof (Theorem(power)):

Assume that A is diagonalizable, so $A\mathbf{x}_j = \lambda_j \mathbf{x}_j$, $\{\mathbf{x}_j\}_{j=1}^n$ linearly independent. Choose \mathbf{v}_0 , and call α_j so that $\mathbf{v}_0 = \sum_{j: \lambda_j = \lambda_1} \alpha_j \mathbf{x}_j + \sum_{j: \lambda_j \neq \lambda_1} \alpha_j \mathbf{x}_j$. Note that the first sum is non null by assumption. Then we have the equalities:

$$A^{m}\mathbf{v}_{0} = \lambda_{1}^{m}\sum_{j:\,\lambda_{j}=\lambda_{1}}\alpha_{j}\mathbf{x}_{j} + \sum_{j:\,\lambda_{j}\neq\lambda_{1}}\alpha_{j}\lambda_{j}^{m}\mathbf{x}_{j},$$

$$\frac{1}{\lambda_{1}^{m}}A^{m}\mathbf{v}_{0} = \sum_{j:\,\lambda_{j}=\lambda_{1}}\alpha_{j}\mathbf{x}_{j} + \sum_{j:\,\lambda_{j}\neq\lambda_{1}}\alpha_{j}\left(\frac{\lambda_{j}}{\lambda_{1}}\right)^{m}\mathbf{x}_{j}.$$

The thesis follows letting m go to infinite.

Assume that A is not diagonalizable. Set $\mu := \lambda_j$, where $\lambda_j \neq \lambda_1$, and restrict, as above, the Jordan matrix equation $X^{-1}AX = J$ to a diagonal Jordan block of order s associated with μ . Then

$$A\mathbf{x}_{r+1} = \mu \mathbf{x}_{r+1}, \ A^m \mathbf{x}_{r+1} = \mu^m \mathbf{x}_{r+1}, A\mathbf{x}_{r+2} = \mu \mathbf{x}_{r+2} + \mathbf{x}_{r+1}, \ A^m \mathbf{x}_{r+2} = \mu^m \mathbf{x}_{r+2} + m\mu^{m-1} \mathbf{x}_{r+1}, A\mathbf{x}_{r+3} = \mu \mathbf{x}_{r+3} + \mathbf{x}_{r+2}, A^m \mathbf{x}_{r+3} = \mu^m \mathbf{x}_{r+3} + m\mu^{m-1} \mathbf{x}_{r+2} + \frac{1}{2}(m^2 - m)\mu^{m-2} \mathbf{x}_{r+1},$$

$$A\mathbf{x}_{r+4} = \mu \mathbf{x}_{r+4} + \mathbf{x}_{r+3}, A^m \mathbf{x}_{r+4} = \mu^m \mathbf{x}_{r+4} + m\mu^{m-1} \mathbf{x}_{r+3} + \frac{1}{2}(m^2 - m)\mu^{m-2} \mathbf{x}_{r+2} + (\frac{1}{6}m^3 - \frac{1}{2}m^2 + \frac{1}{3}m)\mu^{m-3} \mathbf{x}_{r+1},$$

and so on. Set $\lambda := \lambda_1$ and $t = m_a(\lambda_1) = m_g(\lambda_1)$, so we can assume that the upper-left $t \times t$ submatrix of J is diagonal with diagonal entries all equal to λ . Then

$$\begin{aligned}
A^{m}\mathbf{v}_{0} &= A^{m}\left(\sum_{j=1}^{t} \alpha_{j}\mathbf{x}_{j} + \{\dots + \sum_{j=r+1}^{r+s} \alpha_{j}\mathbf{x}_{j} + \dots\}\right) \\
&= \sum_{j=1}^{t} \alpha_{j}\lambda^{m}\mathbf{x}_{j} + \{\dots + \sum_{j=r+1}^{r+s} \alpha_{j}A^{m}\mathbf{x}_{j} + \dots\} \\
&= \lambda^{m}\sum_{j=1}^{t} \alpha_{j}\mathbf{x}_{j} + \{\dots + [\alpha_{r+1}A^{m}\mathbf{x}_{r+1} + \alpha_{r+2}A^{m}\mathbf{x}_{r+2} \\
&+ \alpha_{r+3}A^{m}\mathbf{x}_{r+3} + \alpha_{r+4}A^{m}\mathbf{x}_{r+4} + \dots + \alpha_{r+s}A^{m}\mathbf{x}_{r+s}] + \dots\} \\
&= \lambda^{m}\sum_{j=1}^{t} \alpha_{j}\mathbf{x}_{j} \\
&+ \{\dots + [\alpha_{r+1}(\mu^{m}\mathbf{x}_{r+1}) + \alpha_{r+2}(\mu^{m}\mathbf{x}_{r+2} + m\mu^{m-1}\mathbf{x}_{r+1}) \\
&+ \alpha_{r+3}(\mu^{m}\mathbf{x}_{r+3} + m\mu^{m-1}\mathbf{x}_{r+2} + \frac{1}{2}(m^{2} - m)\mu^{m-2}\mathbf{x}_{r+1}) \\
&+ \alpha_{r+4}(\mu^{m}\mathbf{x}_{r+4} + m\mu^{m-1}\mathbf{x}_{r+3} + \frac{1}{2}(m^{2} - m)\mu^{m-2}\mathbf{x}_{r+2} \\
&+ (\frac{1}{6}m^{3} - \frac{1}{2}m^{2} + \frac{1}{3}m)\mu^{m-3}\mathbf{x}_{r+1}) \\
&+ \dots + \alpha_{r+s}(\mu^{m}\mathbf{x}_{r+s} + \dots + q_{s-1}(m)\mu^{m-s+1}\mathbf{x}_{r+1})] + \dots \},
\end{aligned}$$

$$\begin{aligned}
A^{m}\mathbf{v}_{0} &= \lambda^{m}\sum_{j=1}^{t}\alpha_{j}\mathbf{x}_{j} \\
&+\{\dots + [(\alpha_{r+1}\mu^{m} + \alpha_{r+2}m\mu^{m-1} + \alpha_{r+3}\frac{1}{2}(m^{2} - m)\mu^{m-2} \\
&+\alpha_{r+4}(\frac{1}{6}m^{3} - \frac{1}{2}m^{2} + \frac{1}{3}m)\mu^{m-3} \\
&+\dots + \alpha_{r+s}q_{s-1}(m)\mu^{m-s+1})\mathbf{x}_{r+1} \\
&+(\alpha_{r+2}\mu^{m} + \alpha_{r+3}m\mu^{m-1} + \alpha_{r+4}\frac{1}{2}(m^{2} - m)\mu^{m-2} + \dots \\
&+\alpha_{r+s}q_{s-2}(m)\mu^{m-s+2})\mathbf{x}_{r+2} + \dots + (\alpha_{r+s}\mu^{m})\mathbf{x}_{r+s}] + \dots \}, \\
\frac{1}{\lambda^{m}}A^{m}\mathbf{v}_{0} &= \sum_{j=1}^{t}\alpha_{j}\mathbf{x}_{j} + \{\dots + \left(\frac{\mu}{\lambda}\right)^{m}[(\alpha_{r+1} + \alpha_{r+2}\frac{m}{\mu} + \alpha_{r+3}\frac{\frac{1}{2}(m^{2} - m)}{\mu^{2}} \\
&+\alpha_{r+4}\frac{\frac{1}{6}m^{3} - \frac{1}{2}m^{2} + \frac{1}{3}m}{\mu^{3}} + \dots + \alpha_{r+s}\frac{q_{s-1}(m)}{\mu^{s-1}})\mathbf{x}_{r+1} \\
&+(\alpha_{r+2} + \alpha_{r+3}\frac{m}{\mu} + \alpha_{r+4}\frac{\frac{1}{2}(m^{2} - m)}{\mu^{2}} + \dots + \alpha_{r+s}\frac{q_{s-2}(m)}{\mu^{s-2}})\mathbf{x}_{r+2} \\
&+\dots + (\alpha_{r+s})\mathbf{x}_{r+s}] + \dots \}.
\end{aligned}$$

It is clear that, as m goes to infinite, the sequence $\frac{1}{\lambda^m}A^m\mathbf{v}_0$ converges to an eigenvector of A associated with λ , the dominant eigenvalue of A. Finally, the assertions on the rate of convergence follow by setting

$$p_{s-1}(m) = \alpha_{r+1} + \alpha_{r+2} \frac{m}{\mu} + \alpha_{r+3} \frac{\frac{1}{2}(m^2 - m)}{\mu^2} + \alpha_{r+4} \frac{\frac{1}{6}m^3 - \frac{1}{2}m^2 + \frac{1}{3}m}{\mu^3} + \dots + \alpha_{r+s} \frac{q_{s-1}(m)}{\mu^{s-1}}$$