

---

# Fondamenti della Programmazione: Metodi Evoluti

**Prof. Enrico Nardelli**

Lezione 8: Riferimenti

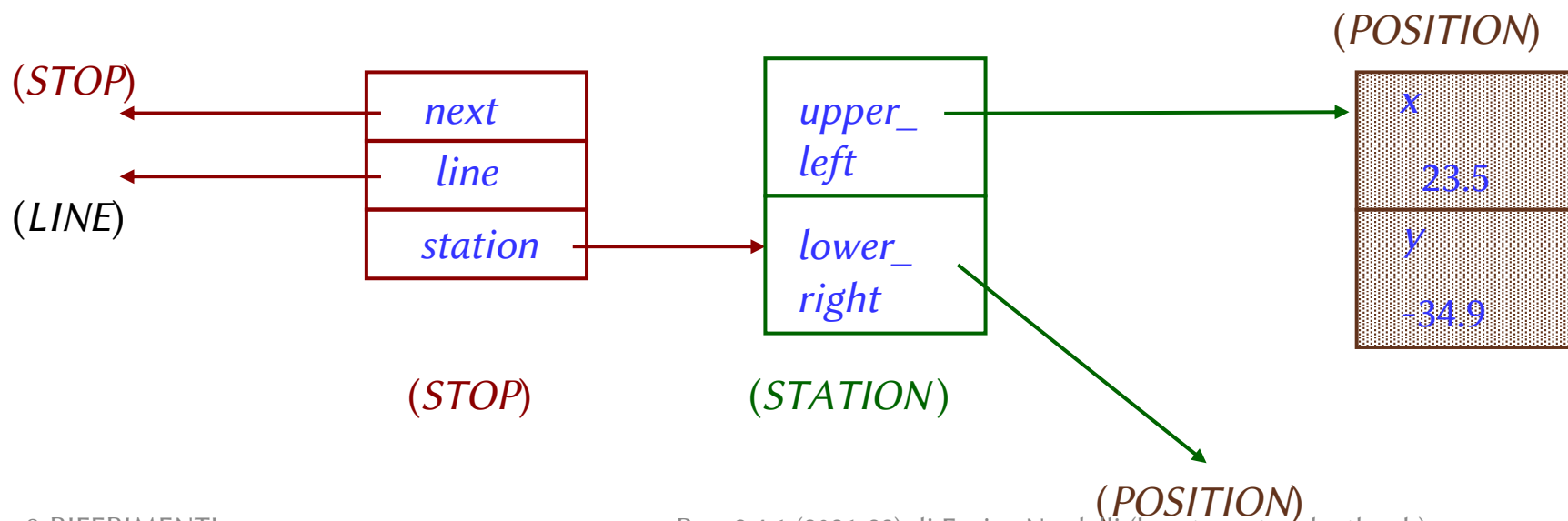
# Object structure (static property)

An object is made of **fields**

Each field is a **value**, which is either:

A **reference** to another object

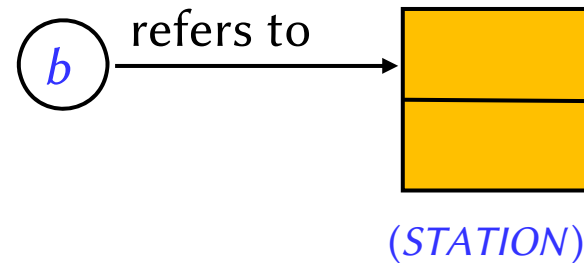
A **basic value**: integer, character, “real” number...  
(known as an **expanded** value)



# There are two types of values

Reference types: the value is a reference.

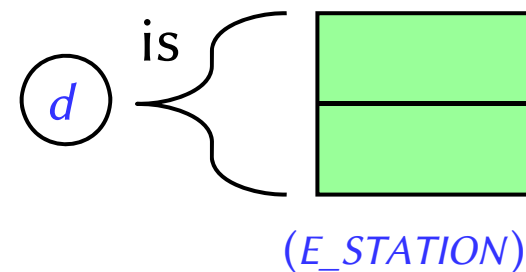
$b: STATION$



A car *has* a brand

Expanded types: the value is an object.

$d: E\_STATION$



A car *has* an engine

## Expanded classes

---

A class may be declared as

**expanded class** *E\_STATION*

... The rest as in *STATION* ...

Then any entity declared

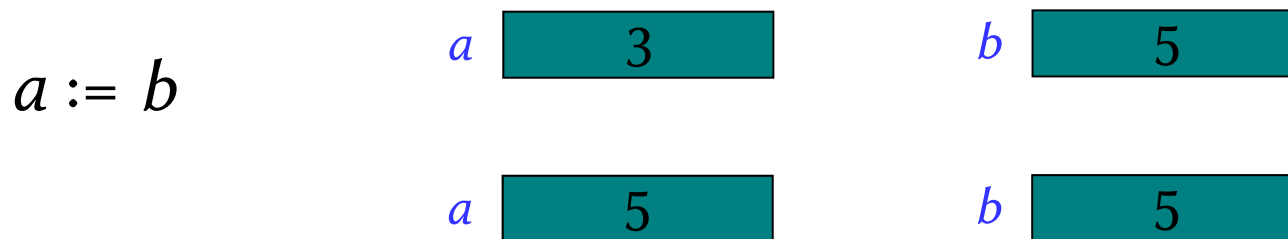
*d: E\_STATION*

has the expanded semantics just described:

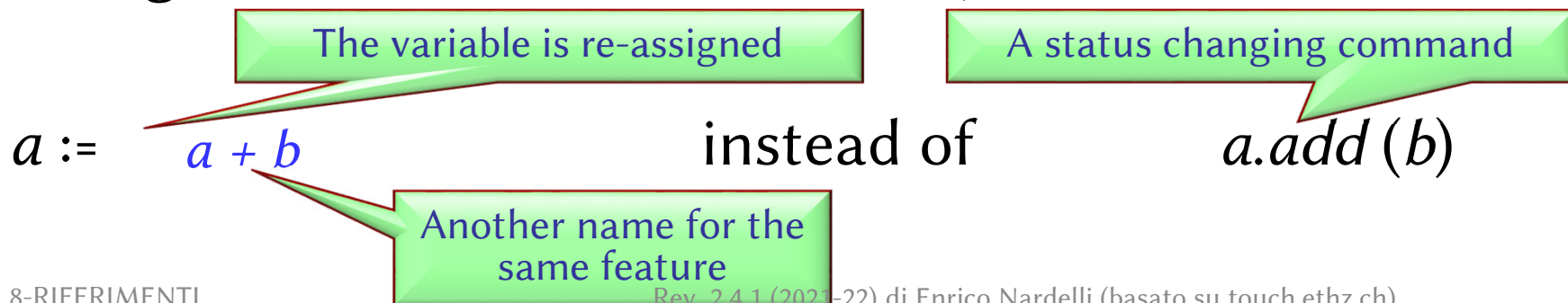
its value is an object.

# Basic types

- So called basic types (*BOOLEAN*, *INTEGER*, *NATURAL*, *REAL*, *CHARACTER*, *STRING*) in Eiffel are classes – just like all other types
- **Most** of them are expanded...



- ... and immutable (they do not contain commands to change the state of their instances)...



## Basic types as expanded classes

---

expanded class *INTEGER* ...

(internally: *INTEGER\_32*, *INTEGER\_64* etc.)

expanded class *BOOLEAN* ...

expanded class *CHARACTER* ...

expanded class *REAL* ...

(internally: *REAL\_32*, *REAL\_64* etc.)

*n*: *INTEGER*

*c*: *CHARACTER*

*s*: *BOOLEAN*

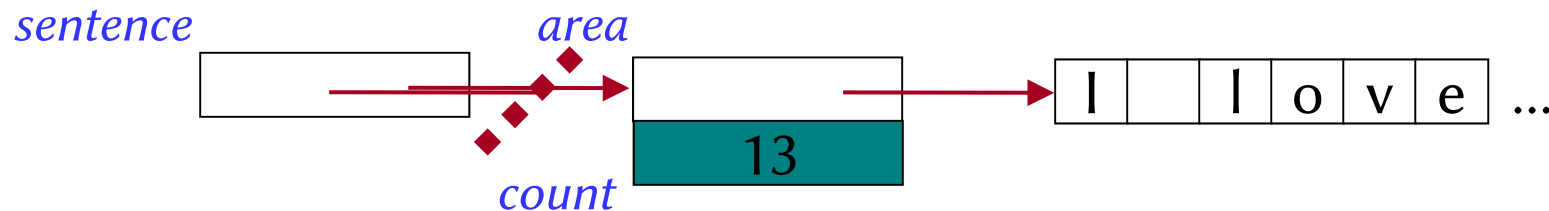
What about strings?

# Strings are a bit different

Strings in Eiffel are **not** expanded:

*sentence* : *STRING*

is a reference to an object



... and **not** immutable

*sentence* := “I love Eiffel”

*sentence.append* (“ very much!”) [compare with  $a := a+b$ ]

## Basic types

---

... their only privilege is to use **manifest constants** to construct their instances:

*b* : *BOOLEAN*

*x* : *INTEGER*

*c* : *CHARACTER*

*s* : *STRING*

...

*b* := **True**

*x* := 5      -- instead of **create** *x.make\_five*

*c* := 'c'

*s* := "I love Eiffel"



It's not an expanded class!



# Initialization

---

Default value of any **reference** type is **Void**

Default values of **basic expanded** types are:

- **False** for *BOOLEAN*
- **0** (zero) for numeric types (*INTEGER*, *NATURAL*, *REAL*)
- “null” character (its *code* = 0) for *CHARACTER*

Default value of an **expanded** type is an object, whose fields have default values of their types

These rules apply to:

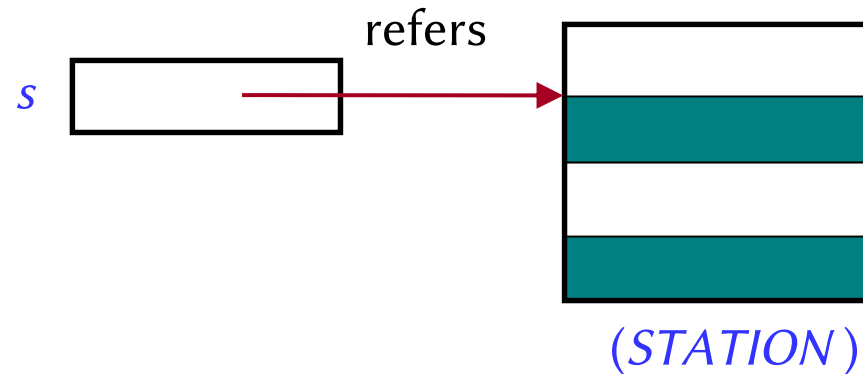
- Fields (from class attributes), on **object creation**
- Local variables, on **start of routine execution** (includes **Result**)

# Two kinds of types

**Reference** types: value of any entity is a reference.

Example:

*s: STATION*



**Expanded** types: value of an entity is an object.

Example:

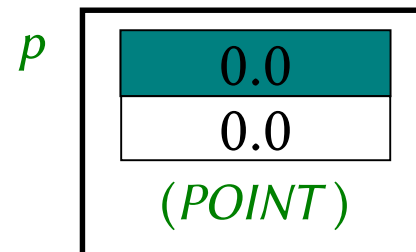
*p: POINT*

... ..

**expanded class** *POINT*

**feature**

*x, y: REAL*



# Objects of reference or expanded types

---

Objects of **reference** types: they don't exist when we declare them (they are initially *Void*).

*s: STATION*

We need to explicitly create them with a create instruction.

*create s*

Objects of **expanded** types: they exist by just declaring them (they are never *Void*)

*p: POINT*

No need to use a **create** instruction

Feature *default\_create* from *ANY* is implicitly invoked on them when creating the containing instance

# How to declare an expanded type

To create an expanded type, declare the class with keyword **expanded**:

**expanded class** *COUPLE*

**feature** -- *Access*

*man, woman : HUMAN*

*years\_together : INTEGER*

**end**

An expanded class

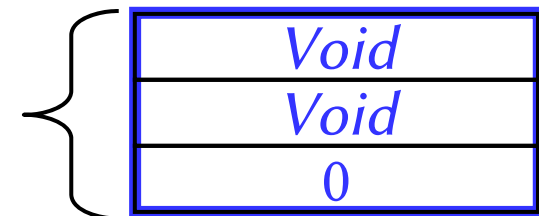
Reference

Expanded

Any entity of type *COUPLE* is expanded:

*adam\_and\_eve : COUPLE*

*adam\_and\_eve*



# Object creation

A standard reference class

**class** *FAMILY*

**feature**

*parents* : *COUPLE*

*home* : *HOUSE*

*phone* : *STRING*

...

**class** *CLAN*

**feature**

*main\_family* : *FAMILY*

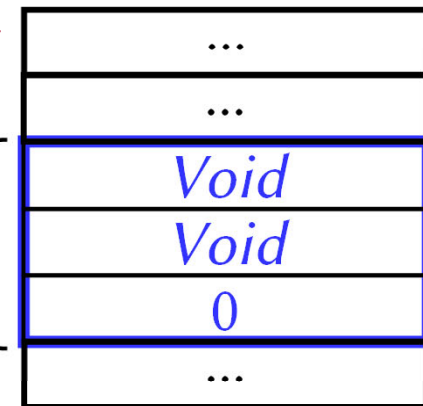
...

**create** *main\_family*

An expanded class

The instance referred to by *main\_family* is created only when this instruction is executed...

*parents*



(*FAMILY*)

# Object creation

```
class RECTANGLE  
feature
```

```
top_left : POINT  
bottom_right : POINT  
cost : INTEGER  
color : STRING  
...
```

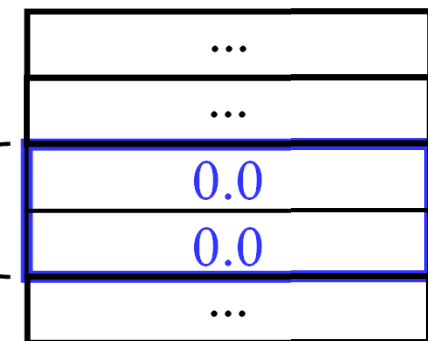
```
class DRAWING  
feature
```

```
rect : RECTANGLE  
...  
create rect
```

A reference class

When this is executed...

*top\_left*



# Initialization



```
class SOME_CLASS
feature p: POINT v: detachable VECTOR
       s: detachable STRING
end
```

When creating an instance of *SOME\_CLASS* which is the default initial value given to its attributes?

```
expanded class POINT
feature x, y: REAL end
class VECTOR
feature x, y: REAL end
STRING
```

<i>x</i>	0.0
<i>y</i>	0.0

(*POINT*)

**Void**

**Void**

# Is this correct?

```
expanded class POINT
create make
feature x, y: REAL
feature make
  do
    x := 5.0
    y := 5.0
  end
  ...
end
```

Incorrect

## ➤ REMEMBER:

- Instances of expanded classes are automatically created when the object containing them is created

For *x*: *POINT*

There is no need of a **create x**



# Custom initialization for expanded types?

---

- Expanded classes can be created **only** in the default way
  - i.e. using *default\_create*, possibly redefining it, if needed for a proper initialization

```
expanded class POINT
inherit ANY
  redefine default_create
feature
  default_create
  do
    x := 5.0; y := 5.0
  end
end
```

## Do you remember this expanded class?

---

expanded class *COUPLE*

feature -- *Access*

*man, woman : HUMAN*

*years\_together : INTEGER*

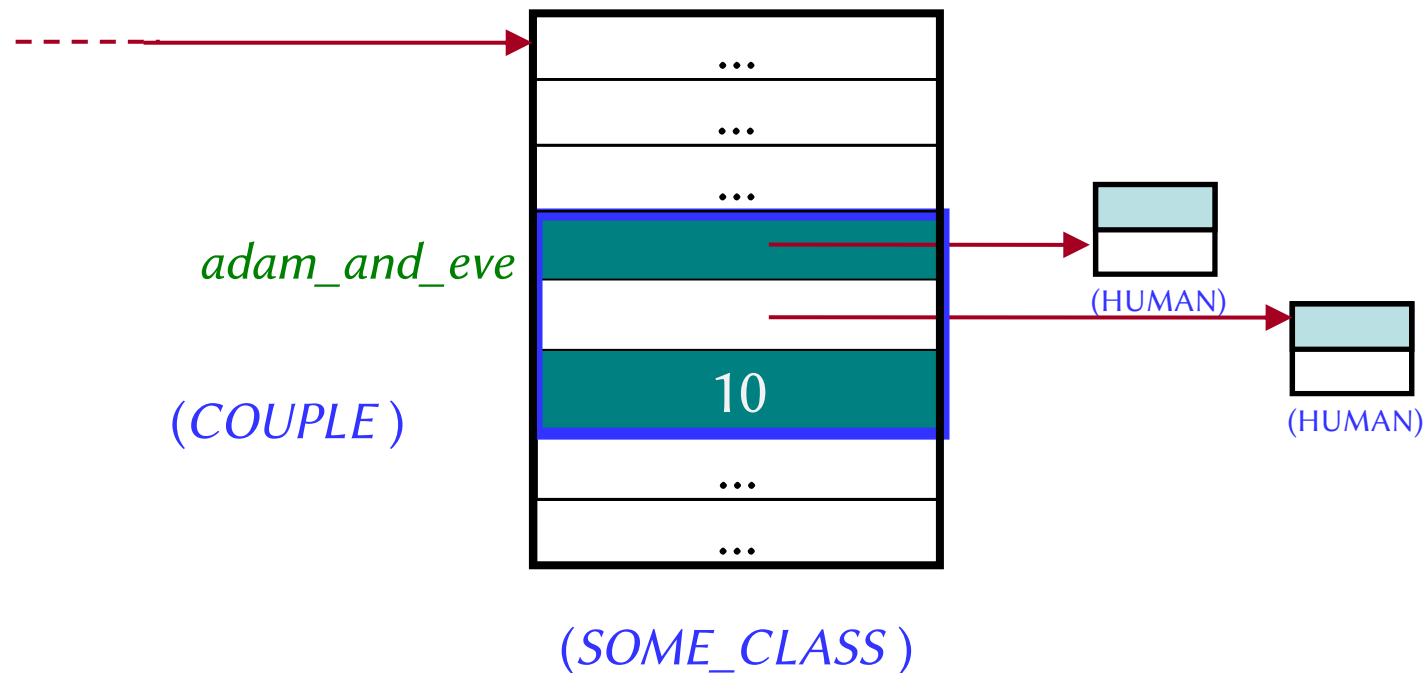
**end**

Assume there is *SOME\_CLASS* definition with this declaration:

*adam\_and\_eve : COUPLE*

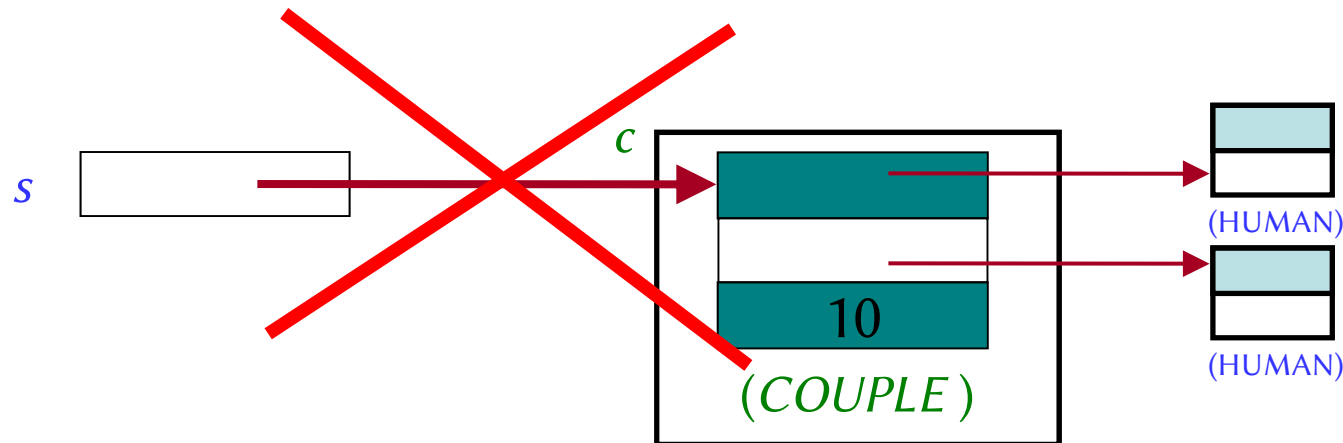
# Can expanded types contain reference types?

An instance of reference type (*SOME\_CLASS*) contains an instance of expanded types (*COUPLE*) which contain instances of a reference type (*HUMAN*)



# Who can reference what?

Objects of expanded types can contain references to other objects...



... but they cannot be referenced by other objects!

# Changing variable values: assignment

*target := source*

*source* is an **expression** and may be:

- Call to a query:
  - *position*
  - *upper\_left.position*
- Arithmetic or boolean expression:
  - $a + (b * c)$
  - $(a < b)$  **and**  $(c = d)$

*target* is a variable entity and may be:

- An attribute
- **Result** in a function
- A “local variable” of a routine

## Semantics

- after the assignment *source* equals *target*
- the value of *source* is not changed by the assignment

# Assignment to attributes

- GENERAL RULE: Direct assignment to an attribute is only allowed if an attribute is called in an unqualified way (i.e., by the object itself). Are the following allowed?

$y := 5$

OK

$x.y := 5$

Error

**Current.y** := 5

Error

- There are two main reasons for the general rule:
1. An other client may not be aware of the restrictions put on the attribute value and interdependencies with other attributes => class invariant violation
  2. The *uniform access principle* (client access independent from the implementation (memory/computation))

## Effect of an assignment (1)

---

**Reference** type: value of any entity is a reference.

$b, c: STATION$

$c := b$

Assignment: copy the reference

**Expanded** type: value of an entity is an object

$d, e: E\_STATION$

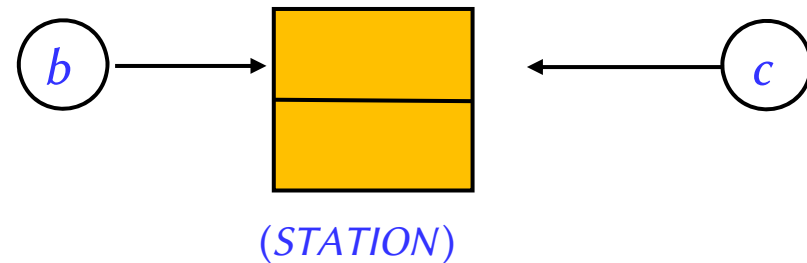
$e := d$

Assignment: copy the object

# Effect of an assignment (2)

**Reference** type: value of any entity is a reference.

$b, c: STATION$

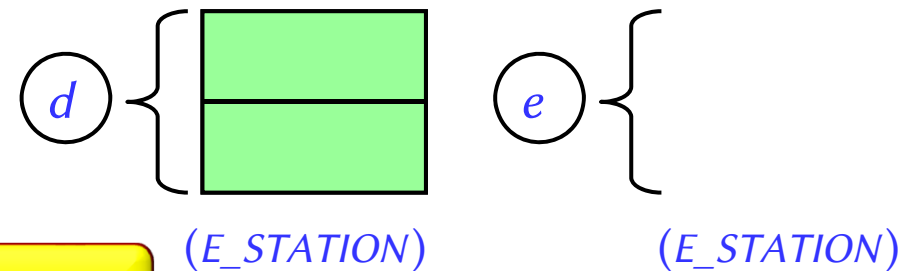


Assignment: copy the reference

$c := b$

**Expanded** type: value of an entity is an object

$d, e: E\_STATION$

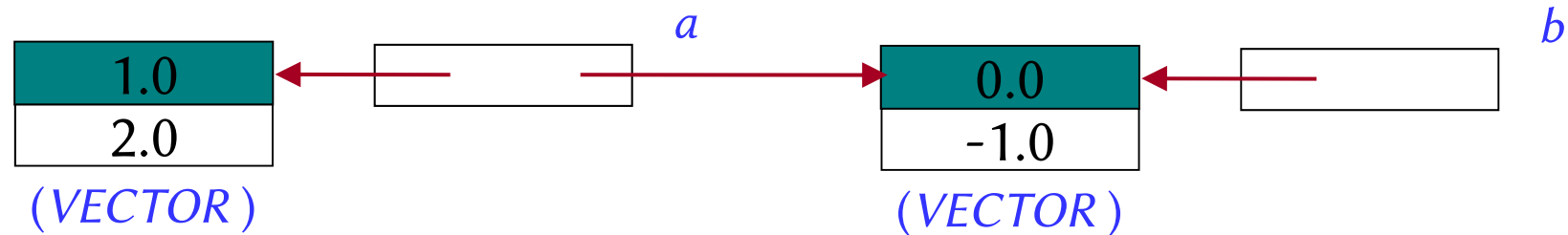


Assignment: copy the object

$e := d$



# Reference assignment



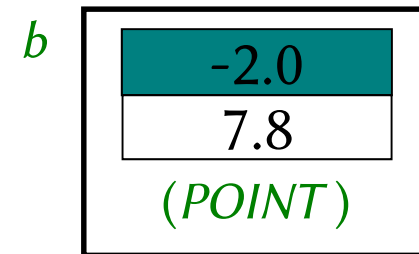
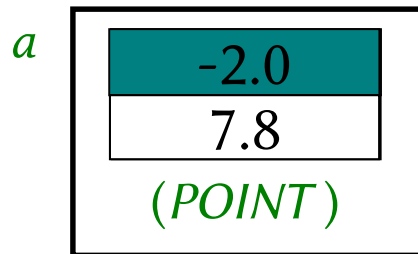
$a := b$

Which is its effect?

now *a* references the same object as *b*:

$a = b$

# Expanded assignment



$a := b$

Which is its effect?

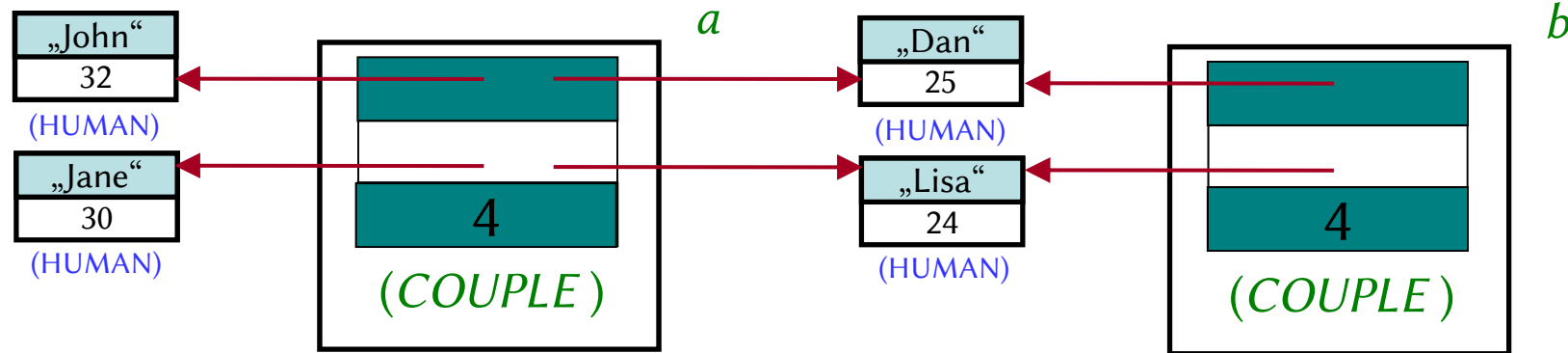
The value of *b* is copied into *a* and, again:

$a = b$

# Assignment

**Hands-On**


Explain graphically the effect of an assignment:



$a := b$

Here **COUPLE** is an expanded class, **HUMAN** is a reference class

# Do not confuse assignment with equality

x  y

Instruction  
(prescriptive and destructive)

if x  y then...

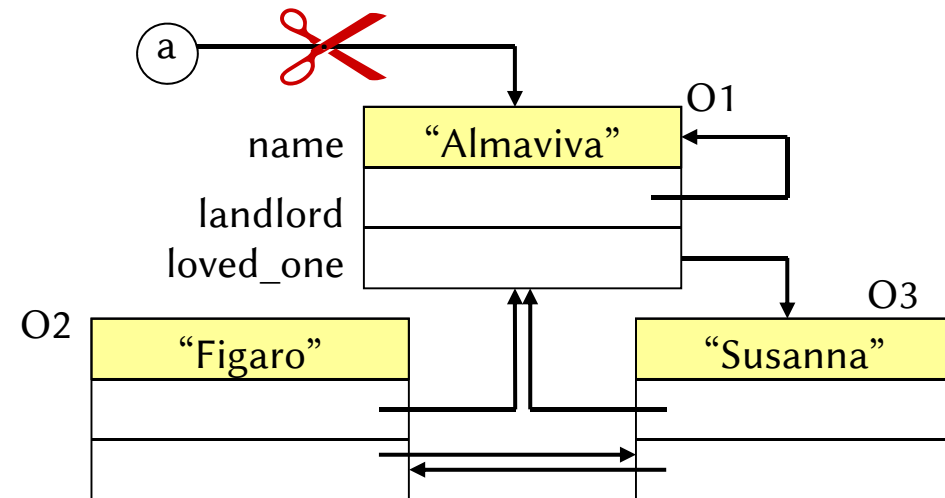
Expression (descriptive)

if x  Current then...

Expression (descriptive)

# What to do with unreachable objects

Reference assignments may make some objects useless



Two possible approaches:

- Manual “free” (C++, Pascal)
- Automatic **garbage collection** (Eiffel, Oberon, Java, .NET)

# Arguments for automatic collection

---

Manual reclamation is dangerous for reliability.

- Wrong “frees” are among the most difficult bugs to detect and correct.

Manual reclamation is tedious.

Modern garbage collectors have acceptable performance overhead.

GC is tunable: disabling, activation, parameterization....

# Properties of a garbage collector (GC)

---

**Consistency** (never reclaim a reachable object).

**Completeness** (reclaim every unreachable object – eventually).

Consistency (also called **safety**) is an absolute requirement. Better no GC than an unsafe GC.

But: safe automatic garbage collection is hard in C-based languages.

# The trouble with reference assignment

---

A comfortable mode of reasoning:

-- Here *SOME\_PROPERTY* holds of *a*

“Apply *SOME\_OPERATION* to *b*”

-- Here *SOME\_PROPERTY* still holds of *a*

It holds for “expanded” values, e.g. integers ( $a=2$ ,  $b=3$ )

-- Here  $P(a)$  holds, e.g. *even* (*a*)

*OPER* (*b*) , e.g. *increment* (*b*)

-- Here  $P(a)$  still holds of *a* , e.g. *even* (*a*)

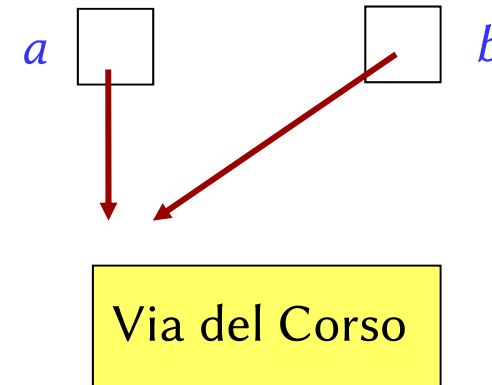


# Dynamic aliasing

*a, b: STRING*

*a := "Via del Corso"*

*b := a*



-- Here *a.item* has value "*Via del Corso*"

*b := "Piazza Venezia"*



-- Here *a.item* has value ??????

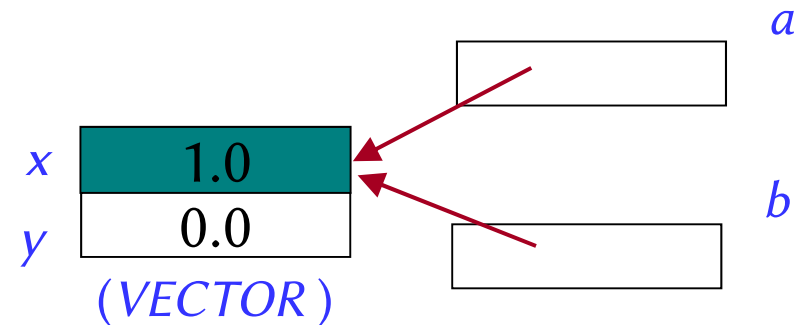
# Dynamic aliasing

$a, b$ : VECTOR

...

**create**  $b.make(1.0, 0.0)$

$a := b$



- now  $a$  and  $b$  reference the same object (are two names or aliases of the same object)
- any change to the object attached to  $a$  will be reflected, when accessing it using  $b$
- any change to the object attached to  $b$  will be reflected, when accessing it using  $a$

# Dynamic aliasing (1)



What are the values of  $a.x$ ,  $a.y$ ,  $b.x$  and  $b.y$  after executing each instruction 1 to 3?

$a, b$ : VECTOR

...

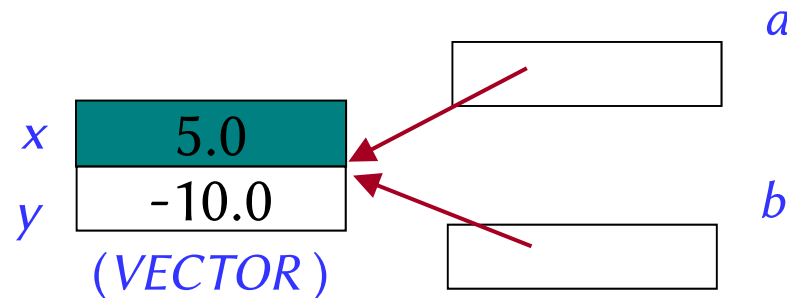
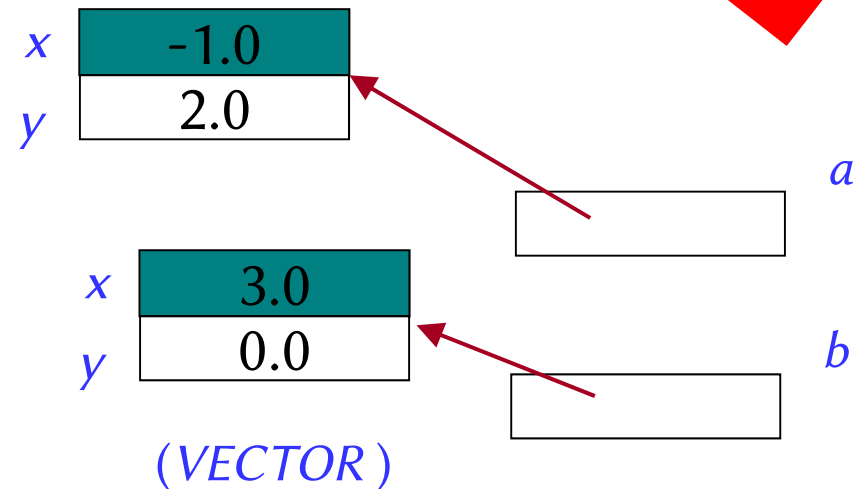
**create**  $a.make$  (-1.0, 2.0)

**create**  $b.make$  (3.0, 0.0)

1  $a := b$

2  $b.set\_x$  (5.0)

3  $a.set\_y$  (-10.0)



# Dynamic aliasing (2)



What are the values of  $a.x$ ,  $a.y$ ,  $b.x$  and  $b.y$  after executing each instruction 1-3?

Now assume `VECTOR` is an *expanded* class

$a, b$ : `VECTOR`

...

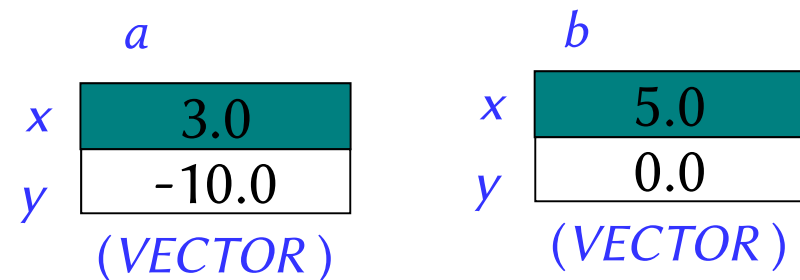
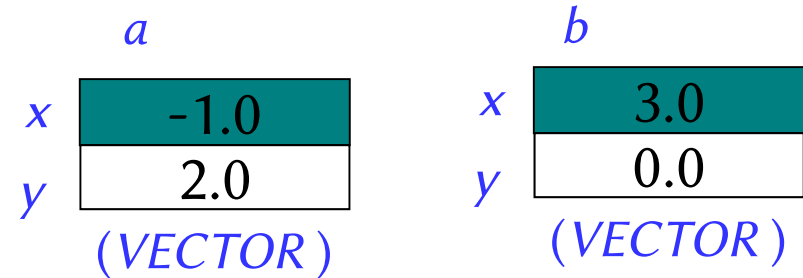
$a.set(-1.0, 2.0)$

$b.set(3.0, 0.0)$

1  $a := b$

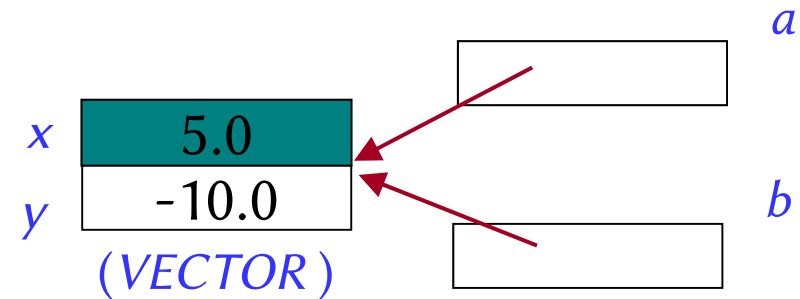
2  $b.set_x(5.0)$

3  $a.set_y(-10.0)$

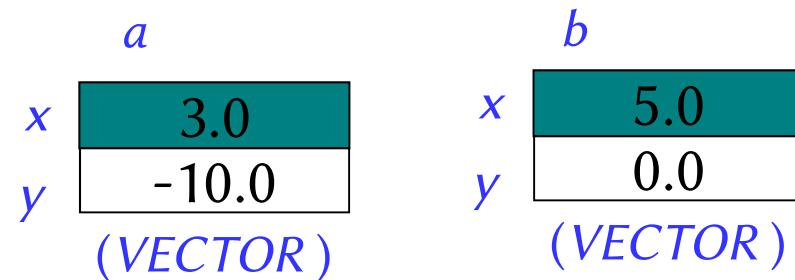


# Compare the final result

VECTOR defined as **reference** type



VECTOR defined as **expanded** type



## Practical advice

---

Reference assignment is useful

It's also potentially tricky

As much as possible, leave it to specialized libraries of general data structures

## Variants of assignment and copy

---

Reference assignment (*a* and *b* of the same reference types):

*b := a*

Object duplication (**shallow** – creates a **new** object whose fields receive copy of values):

*c := a.twin*

Object duplication (**deep** – creates a **new** object whose reference fields, if any, are deep duplicated):

*d := a.deep\_twin*

Also: shallow field-by-field copy (an **existing** object receives copy of field values of another object):

*e.copy(a)*

NOTE: *c* and *d* are “*created on the fly*” by the duplication operation (or reassigned). Instead *e* must already exist

# Shallow and deep cloning

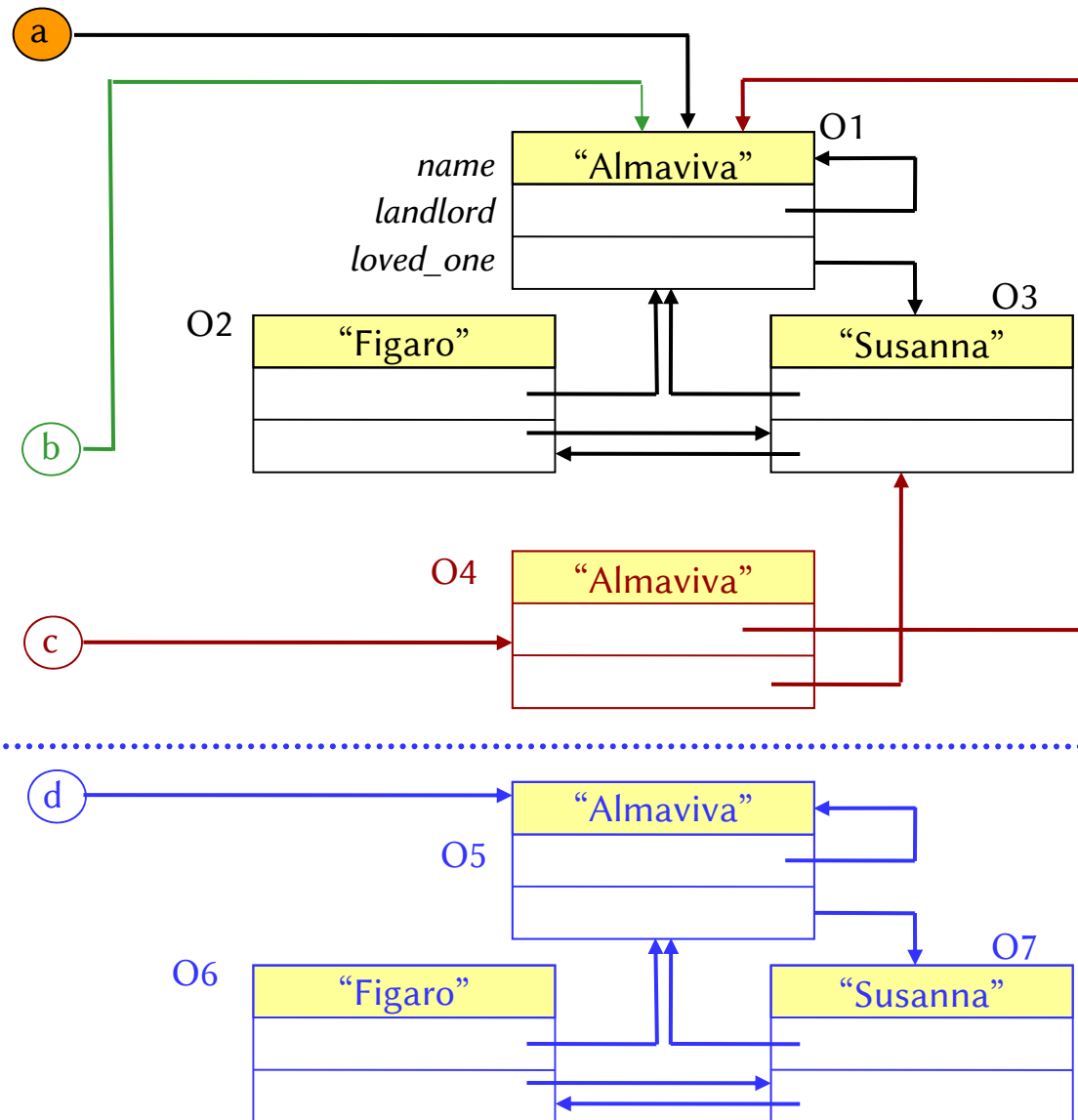
Initial situation:

Result of:

$b := a$

$c := a.twin$

$d := a.deep\_twin$





## Equality testing (1)

---

Let  $a$  and  $b$  be variable of a **reference** type

To test equality of their values (which are references) use the instruction for *reference equality*

$$a = b$$

To test equality of their referenced objects, i.e. of items referenced by their values, use the instruction for ***object equality***

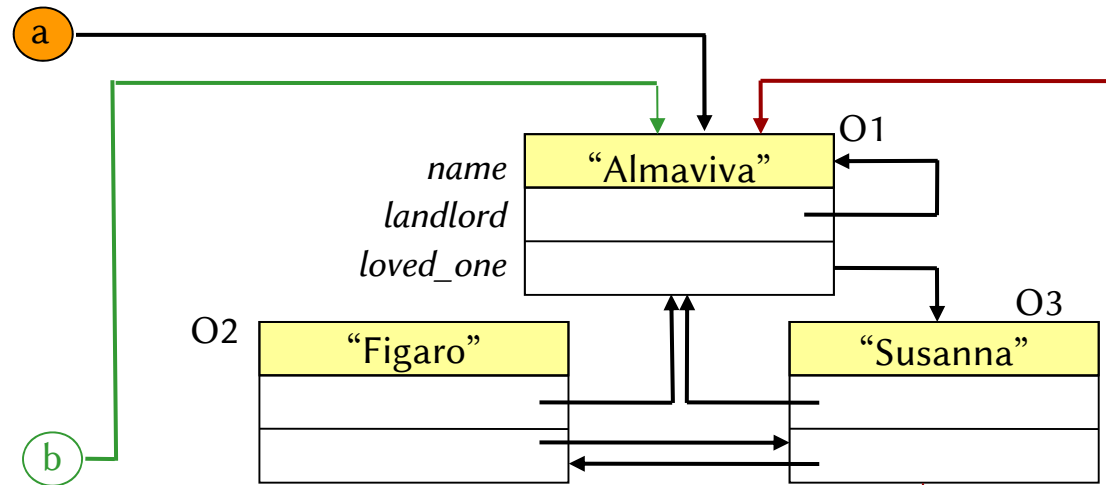
$$a \sim b$$

which is implemented by testing that all attributes are “*object equal*”

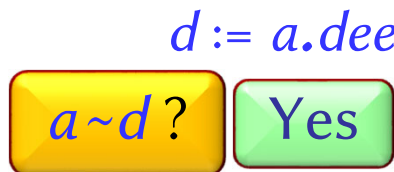
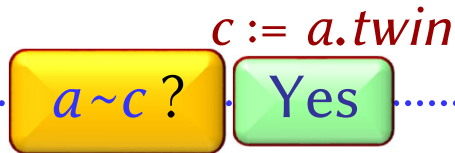
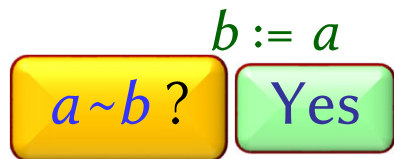
N.B.: a **recursive** definition

# Shallow and deep cloning: equality testing

Initial situation:



After:



(b)

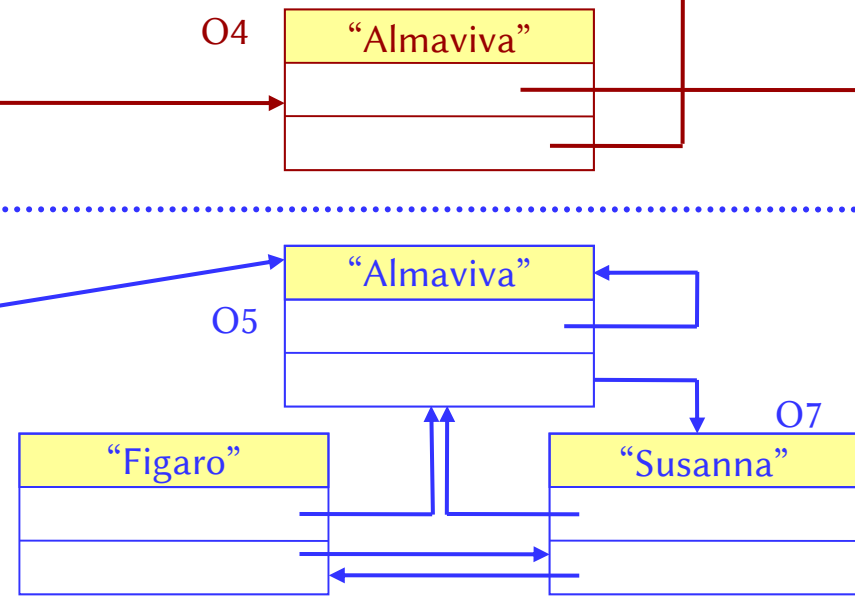
(c)

(d)

O4

O5

O7

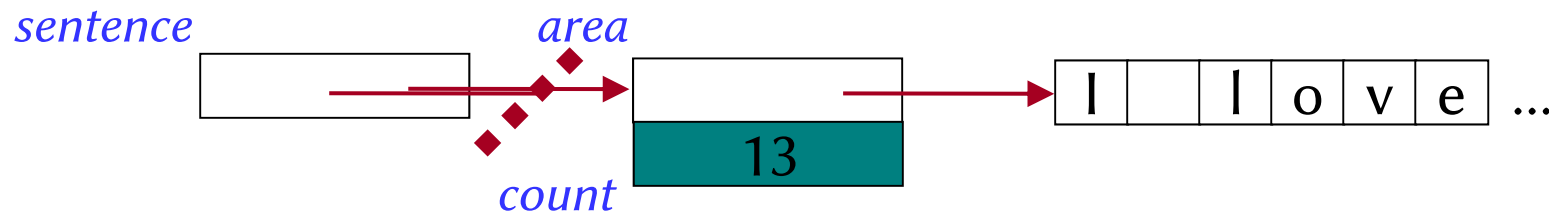


# Remember strings

Strings in Eiffel are **not** expanded:

*sentence* : *STRING*

is a reference to an object



... and **not** immutable

*sentence* := “I love Eiffel”

*sentence.append* (“ very much!”) [compare with *a := a+b*]

# Strings assignment and comparison

*my\_name, your\_name: STRING*

Three ways of testing equality

`is_equal` (a *STRING* feature)

`=` (equality of entities' values)

`~` (equality of referenced objects)

object equality

reference equality

object equality

*my\_name := "mario"*  
*your\_name := "mario"*

`is_equal ?`

`= ?`

`~ ?`

True

False

True

*my\_name := "mario"*  
*your\_name := my\_name*

True

True

True

*my\_name := "mario"*  
*your\_name := my\_name.twin*

True

False

True

# Object comparison: = versus ~

*s1: STRING* = “Teddy”

*s2: STRING* = “Teddy”

...

*s1 = s2*

False

reference comparison on different objects

*s1 ~ s2*

True

object comparison

...

Now you know what to do if interested in comparing the content of two strings

# Reading and assigning strings

---

`lo` is a predefined object in Eiffel referring to the system input (keyboard) and output (screen)

`last_string` is a query providing the (reference to the) last string that was read by `read_line` feature (but for the new line character – which is consumed but is not part of `last_string`)

Subsequent invocations of `last_string` do **not** provide a new references. Instead, the string referred to by `last_string` is rewritten every time a new string is read from the file

Then after

```
lo.read_line; s1 := lo.last_string
```

```
lo.read_line; s2 := lo.last_string
```

it will always be  $s1 = s2 =$  the last input line read

Use instead

```
lo.read_line; s1 := lo.last_string.twin
```

```
lo.read_line; s2 := lo.last_string.twin
```

to read in `s1` and `s2` the two input lines

## Equality testing (2)

---

Given Eiffel is a strongly typed language:

$a = b$  implies  $a \sim b$

while

when  $a \sim b$  it can be  $a \neq b$

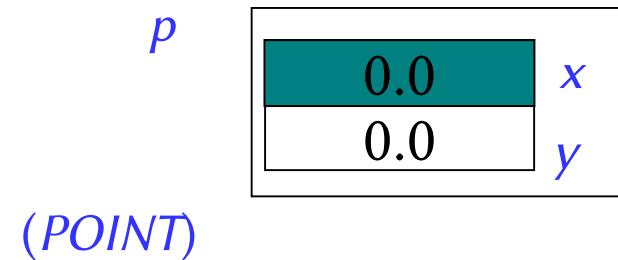
For variables of "pure" expanded type both tests always provide the same result!

An expanded type is "pure" if it is "basic" or all subtypes composing it are "pure"

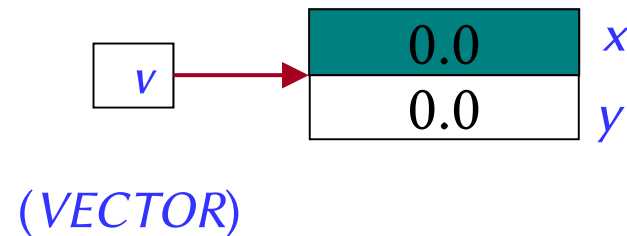
But for variables of "mixed" expanded types the two tests can have different results

# Remember

expanded class *POINT*  
 feature  $x, y: REAL$   
 end



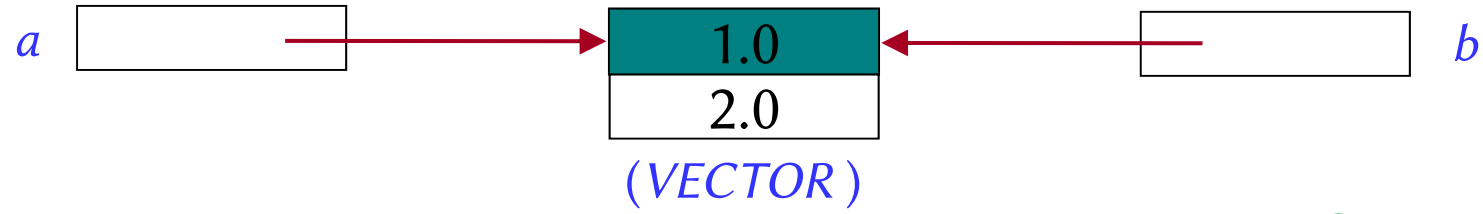
class *VECTOR*  
 feature  $x, y: REAL$   
 end





# Object equality (reference types)

Hands-On

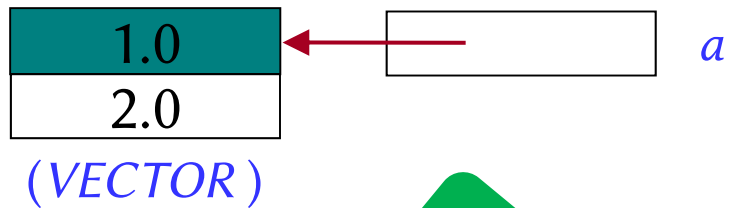


*a* ~ *b*?

True

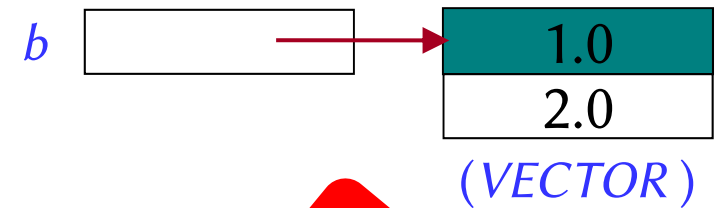
*a* = *b*?

True



*a* ~ *b*?

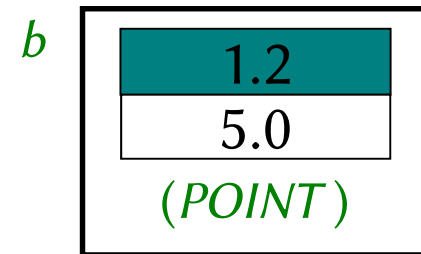
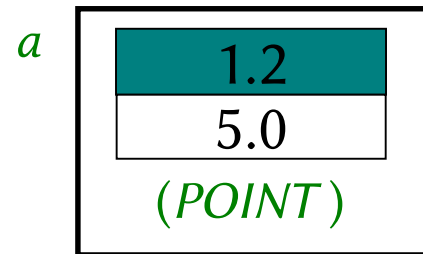
True



*a* = *b*?

False

# Object equality (expanded types)



Hands-On

$a \sim b?$

$a = b?$

True

True

Entities of expanded types are compared by value!  
*POINT* is a "pure" expanded type: for such a type, reference equality and object equality always provide the same result!

# Remember

---

expanded class *COUPLE*

feature -- Access

*man, woman : HUMAN*

*years\_together : INTEGER*

end



Reference type

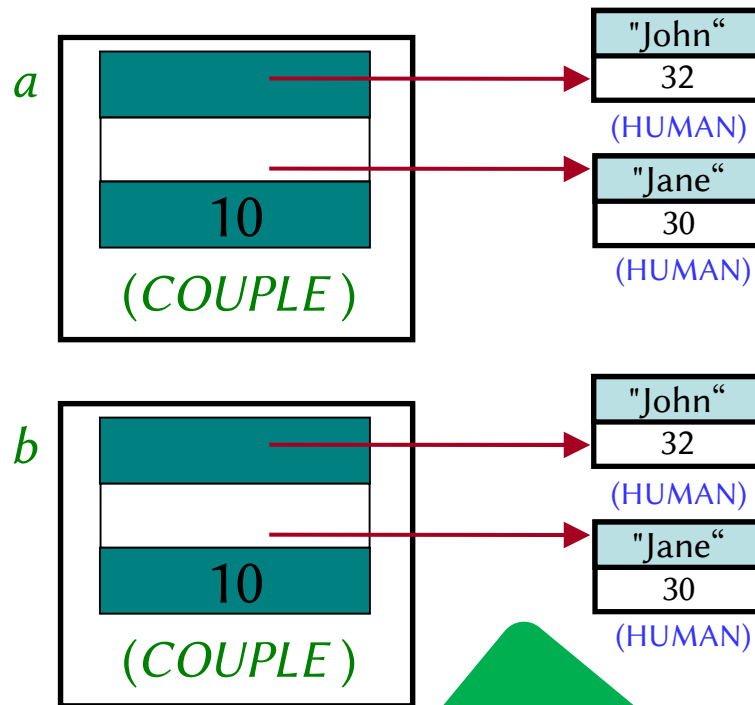
Any entities of type *COUPLE* is automatically an expanded entity:

*adam\_and\_eve : COUPLE*



Mixed expanded type

# Expanded entities equality



$a \sim b?$

True

$a = b?$

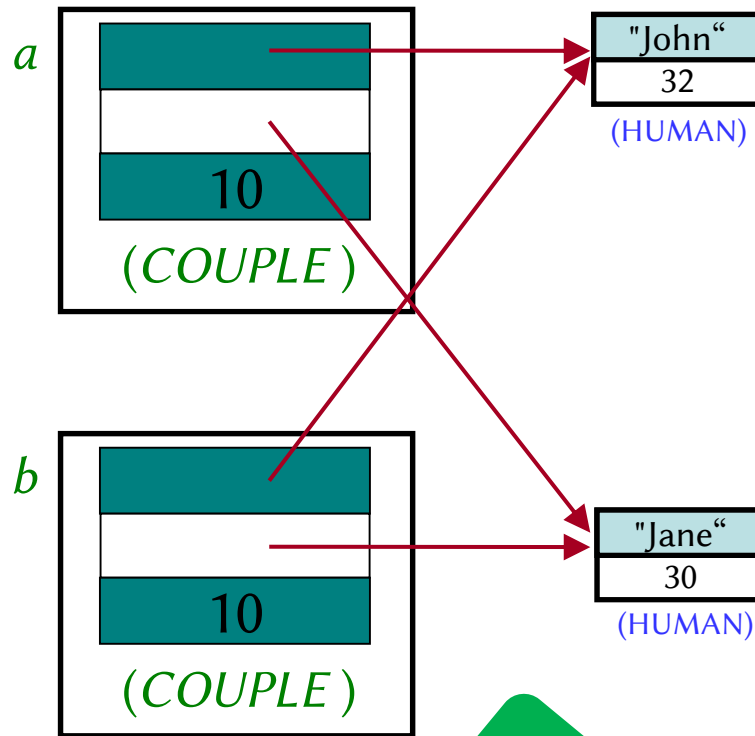
False

Hands-On

*COUPLE* is NOT a "pure" expanded type: for such a type reference equality and object equality do not always provide the same result!

# Expanded entities equality

Hands-On



$a \sim b?$

True

$a = b?$

True

*COUPLE* is NOT a "pure" expanded type, but in this case both equality tests give the same result!

# Attachment

---

➤ Is a more general term than assignment

➤ Includes:

- Assignment

$a := b$

- Passing arguments to a routine

$f(a: \text{SOME\_TYPE})$

**do ... end**

... ..

$f(b)$

➤ Same semantics