Ramanujan in his first paper states that the Bernoulli numbers satisfy the following equations

(actually, since Ramanujan Bernoulli numbers are the moduli of ours, his equations are a bit different).

Call

$$R \begin{bmatrix} B_2(0) \\ B_4(0) \\ B_6(0) \\ \vdots \end{bmatrix} = \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ \vdots \end{bmatrix}$$
 (RamanSyst)

such semi-infinite Ramanujan linear system. We observe that such system is equivalent to the following linear systems:

$$R\begin{bmatrix} \frac{2!}{x} & \frac{4!}{x^2} \\ & \frac{4!}{x^2} \end{bmatrix} \begin{bmatrix} \frac{x}{2!} & \frac{x^2}{4!} \\ & \ddots \end{bmatrix} \begin{bmatrix} B_2(0) \\ B_4(0) \\ B_6(0) \\ \vdots \end{bmatrix} = \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ \vdots \end{bmatrix}$$

$$\begin{bmatrix} \frac{2!}{x} & \frac{4!}{x^2} \\ & \frac{4!}{x^2} \end{bmatrix} \tilde{R}\begin{bmatrix} \frac{x}{2!} & \frac{x^2}{4!} \\ & \frac{x^2}{4!} \end{bmatrix} \begin{bmatrix} B_2(0) \\ B_4(0) \\ B_6(0) \\ \vdots \end{bmatrix} = \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ \vdots \end{bmatrix}$$

$$\tilde{R}\begin{bmatrix} \frac{x}{2!} & \frac{x^2}{4!} \\ & \ddots \end{bmatrix} \begin{bmatrix} B_2(0) \\ B_4(0) \\ B_6(0) \\ \vdots \end{bmatrix} = \begin{bmatrix} \frac{x^2}{4!} & \frac{x^2}{4!} \\ & \ddots \end{bmatrix} \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ \vdots \end{bmatrix}, \quad \tilde{R} = \sum_{j=0}^{+\infty} \alpha_j Z^j = \sum_{s=0}^{+\infty} \frac{2x^{3s}}{(6s+2)!(2s+1)} Z^{3s},$$

$$\alpha_0 = 1, \quad \alpha_1 = \alpha_2 = 0, \quad \alpha_2 = \frac{2}{2} x^3, \quad \alpha_4 = \alpha_5 = 0, \quad \alpha_6 = \frac{2}{2} x^6, \quad \alpha_7 = \alpha_9 = 0$$

 $\alpha_0=1,\ \alpha_1=\alpha_2=0,\ \alpha_3=\frac{2}{8!3}x^3,\ \alpha_4=\alpha_5=0,\ \alpha_6=\frac{2}{14!5}x^6,\ \alpha_7=\alpha_8=0,\ \alpha_9=\frac{2}{20!7}x^9,\ldots$  The symbol Z denotes the semi-infinite lower shift matrix:

$$Z = \left[ \begin{array}{ccc} 0 & & & \\ 1 & 0 & & \\ & 1 & 0 & \\ & & \cdot & \cdot \end{array} \right]$$

So, the Ramanujan system is equivalent to a lower triangular Toeplitz (sparse) linear system. Can the latter system be obtained from the lower triangular Toeplitz (dense) even and odd linear systems introduced in toe\_1n? First let us prepare things.

Theorem.

Let  $Z_{n-1}$  and  $Z_n$  be the upper-left  $n-1\times n-1$  and  $n\times n$  submatrices of Z. Then

$$\sum_{j=0}^{n-1} \alpha_j Z_n^j \begin{bmatrix} B_0(0) \\ \frac{x}{2!} B_2(0) \\ \frac{x^2}{4!} B_4(0) \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} B_{2(n-1)}(0) \end{bmatrix} = \begin{bmatrix} \Box \\ \frac{x}{2!} f_1 \\ \frac{x^2}{4!} f_2 \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} f_{n-1} \end{bmatrix} + B_0(0) \begin{bmatrix} ARB \\ \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_{n-1} \end{bmatrix} \left( = \begin{bmatrix} w_0 \\ \frac{x}{2!} w_1 \\ \frac{x^2}{4!} w_2 \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} w_{n-1} \end{bmatrix} \right)$$

if and only if

$$\sum_{j=0}^{n-2}\alpha_{j}Z_{n-1}^{j}\left[\begin{array}{c}\frac{\frac{x}{2!}}{2!}B_{2}(0)\\\frac{x^{2}}{4!}B_{4}(0)\\\vdots\\\frac{x^{n-1}}{(2(n-1))!}B_{2(n-1)}(0)\end{array}\right]=\left[\begin{array}{c}\frac{\frac{x}{2!}}{2!}f_{1}\\\frac{\frac{x^{2}}{2!}}{4!}f_{2}\\\vdots\\\frac{x^{n-1}}{(2(n-1))!}f_{n-1}\end{array}\right]\left(\begin{array}{c}=\left[\begin{array}{c}\frac{x}{2!}w_{1}\\\frac{x^{2}}{2!}w_{2}\\\vdots\\\frac{x^{n-1}}{(2(n-1))!}w_{n-1}\end{array}\right]-B_{0}(0)\left[\begin{array}{c}\alpha_{1}\\\alpha_{2}\\\vdots\\\alpha_{n-1}\end{array}\right]\right)$$

The implication " $\Leftarrow$ " holds if  $\alpha_0 B_0(0) = \Box + B_0(0) \cdot ARB \ (\alpha_0 B_0(0) = w_0)$ .

Application of the Theorem:

We have shown that the Ramanujan system is equivalent to the system

Then, by the Theorem,

$$\sum_{j=0}^{n-1} \alpha_j Z_n^j \begin{bmatrix} B_0(0) \\ \frac{x}{2!} B_2(0) \\ \frac{x^2}{2!} B_4(0) \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} B_{2(n-1)}(0) \end{bmatrix} = \begin{bmatrix} \Box \\ \frac{x}{2!} f_1 \\ \frac{x^2}{2!} f_2 \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} f_{n-1} \end{bmatrix} + B_0(0) \begin{bmatrix} ARB \\ \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_{n-1} \end{bmatrix}, \ \Box + B_0(0) \cdot ARB = \alpha_0 B_0(0),$$

$$(I + \frac{2}{8!3}x^3Z^3 + \frac{2}{14!5}x^6Z^6 + \frac{2}{20!7}x^9Z^9 + \dots) \begin{bmatrix} B_0(0) \\ \frac{x}{2!}B_2(0) \\ \frac{x}{4!}B_4(0) \\ \frac{x^6}{8!}B_8(0) \\ \frac{x^6}{12!}B_{12}(0) \\ \vdots \end{bmatrix} = \begin{bmatrix} 1 \\ \frac{x^2}{2!}\frac{1}{6} \\ \frac{x^2}{4!}(-\frac{1}{30}) \\ \frac{x^3}{6!}\frac{1}{4!} + \frac{2x^3}{8!3} \\ \frac{x^4}{6!}\frac{1}{4!5} \\ \frac{x^5}{8!}\frac{1}{4!5} \\ \frac{x^5}{10!}(-\frac{1}{132}) \\ \frac{x^5}{12!}\frac{1}{4!5} + \frac{2x^6}{12!} \\ \frac{x^5}{4!}\frac{1}{120} \\ \frac{x^5}{12!}\frac{1}{4!5} + \frac{2x^6}{14!5} \\ \frac{x^5}{12!}(-\frac{1}{306}) \\ \frac{x^5}{12!}\frac{3}{6!}(-\frac{1}{306}) \\ \frac{x^5}{12!}(-\frac{1}{10}) \\ \frac{x^5}{12!}(-\frac{1}{10})$$

$$\sum_{j=0}^{+\infty} \alpha_j Z^j D_x \mathbf{b} = D_x \mathbf{q}^R,$$

$$\alpha_j = \delta_{j=0 \, \text{mod} \, 3} \frac{2x^j}{(2j+2)!(\frac{2}{3}j+1)}, \ \mathbf{q}^R = (\frac{1}{(2i+1)(i+1)}(1-\delta_{i=2 \, \text{mod} \, 3} \frac{3}{2})), \ i=0,1,2,3,\dots.$$

Note that from the explicit expression of  $\mathbf{q}^R$  it follows an explicit expression for the  $f_i$  of the Ramanujan system:

$$f_i = \frac{1}{(2i+1)(i+1)} (1 - \delta_{i=2 \mod 3} \frac{3}{2} - \delta_{i=0 \mod 3} \frac{1}{\frac{2}{3}i+1}), \quad i = 1, 2, 3, \dots$$

Note also that (\*) can be rewritten as

$$L(\underline{\alpha}_{n-1})I_n^2 D_x \mathbf{b} = \operatorname{diag}(z_i, i = 1, 2, \dots, n-1)I_n^2 D_x \mathbf{q}^R$$

for suitable  $z_i$ . Let us look for such  $z_i$ :

$$\begin{split} &(1-\delta_{i=2\,\mathrm{mod}\,3}\tfrac{3}{2}-\delta_{i=0\,\mathrm{mod}\,3}\tfrac{1}{\tfrac{2}{3}i+1})=z_i(1-\delta_{i=2\,\mathrm{mod}\,3}\tfrac{3}{2}),\\ &z_i=1-\tfrac{\delta_{i=0\,\mathrm{mod}\,3}\tfrac{1}{\tfrac{3}{2}i+1}}{1-\delta_{i=2\,\mathrm{mod}\,3}\tfrac{3}{2}}\\ &=1-\delta_{i=0\,\mathrm{mod}\,3\,\tfrac{2}{3}i+1}. \end{split}$$

Now let us consider the two even and odd systems introduced in toe\_1n whose solution is again the vector of Bernoulli numbers:

$$\sum_{j=0}^{+\infty} \frac{x^{j}}{(2j+2)!} Z^{j} D_{x} \mathbf{b} = D_{x} \mathbf{q}^{e}, \ \mathbf{q}^{e} = \begin{bmatrix} \frac{1}{2} \\ \frac{1}{6} \\ \frac{1}{10} \end{bmatrix}, \sum_{j=0}^{+\infty} \frac{x^{j}}{(2j+1)!} Z^{j} D_{x} \mathbf{b} = D_{x} \mathbf{q}^{o}, \ \mathbf{q}^{o} = \begin{bmatrix} \frac{1}{2} \\ \frac{1}{2} \\ \frac{1}{2} \end{bmatrix};$$

$$\sum_{j=0}^{n-1} \frac{x^{j}}{(2j+2)!} Z_{n}^{j} D_{x} \mathbf{b} = D_{x} \mathbf{q}^{e}, \quad \sum_{j=0}^{n-1} \frac{x^{j}}{(2j+1)!} Z_{n}^{j} D_{x} \mathbf{b} = D_{x} \mathbf{q}^{o}.$$

And apply to them the Theorem:

$$\sum_{j=0}^{n-2} \frac{x^{j}}{(2j+2)!} Z_{n-1}^{j} \begin{bmatrix} \frac{x_{1}}{2!} B_{2}(0) \\ \frac{x_{2}}{4!} B_{4}(0) \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} B_{2(n-1)}(0) \end{bmatrix} = \begin{bmatrix} \frac{x_{1}}{2!} \frac{1}{6} \\ \frac{x^{2}}{4!} \frac{1}{10} \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} \frac{1}{4n-2} \end{bmatrix} -B_{0}(0) \begin{bmatrix} \frac{x_{1}}{4!} \\ \frac{x^{2}}{2!} \frac{1}{6!} \\ \frac{x^{2}}{6!} \\ \vdots \\ \frac{x^{n-1}}{(2n)!} \end{bmatrix} = \begin{bmatrix} \frac{x_{1}}{4!} \frac{1}{10} \\ \frac{x^{2}}{8!} \frac{1}{8!} \\ \vdots \\ \frac{x^{n-1}}{(2n)!} \end{bmatrix}$$

$$\sum_{j=0}^{n-2} \frac{x^{j}}{(2j+1)!} Z_{n-1}^{j} \begin{bmatrix} \frac{x_{1}}{2!} B_{2}(0) \\ \frac{x_{1}}{2!} B_{4}(0) \\ \vdots \\ \frac{x^{n-1}}{(2n-1)!} B_{2(n-1)}(0) \end{bmatrix} = \begin{bmatrix} \frac{x_{1}}{2!} \frac{1}{2} \\ \frac{x_{1}^{2}}{4!} \frac{1}{2} \\ \vdots \\ \frac{x^{n-1}}{(2n-1)!} \frac{1}{2!} \end{bmatrix} -B_{0}(0) \begin{bmatrix} \frac{x_{1}}{3!} \\ \frac{x^{2}}{5!} \\ \frac{x^{2}}{5!} \\ \frac{x^{n-1}}{(2n-1)!} \end{bmatrix} = \begin{bmatrix} \frac{x_{1}}{3!} \frac{1}{2} \\ \frac{x^{2}}{5!} \frac{1}{2} \\ \frac{x^{n-1}}{(2n-1)!} \frac{1}{2!} \end{bmatrix}$$

from which it follows

$$\sum_{j=0}^{n-2} \frac{x^{j}}{(2j+2)!} Z_{n-1}^{j} \begin{bmatrix} \frac{x}{2!} B_{2}(0) \\ \frac{x^{2}}{4!} B_{4}(0) \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} B_{2(n-1)}(0) \end{bmatrix} = \begin{bmatrix} \frac{x}{2!} \frac{1}{4 \cdot 3} \\ \frac{x^{2}}{6!} \frac{2}{6 \cdot 5} \\ \frac{x^{3}}{6!} \frac{3}{8 \cdot 7} \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} \frac{n-1}{2n(2n-1)} \end{bmatrix} = \operatorname{diag}(\frac{i}{i+1}, i = 1 \dots n-1) I_{n-1}^{1}(Z_{n}^{T} D_{x} \mathbf{q}^{e})$$

$$\sum_{j=0}^{n-2} \frac{x^{j}}{(2j+1)!} Z_{n-1}^{j} \begin{bmatrix} \frac{x}{2!} B_{2}(0) \\ \frac{x^{2}}{4!} B_{4}(0) \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} B_{2(n-1)}(0) \end{bmatrix} = \begin{bmatrix} \frac{x}{2!} \frac{1}{2 \cdot 3} \\ \frac{x^{2}}{4!} \frac{3}{2 \cdot 5} \\ \frac{x^{3}}{6!} \frac{5}{2 \cdot 7} \\ \vdots \\ \frac{x^{n-1}}{(2(n-1))!} \frac{2n-3}{2(2n-1)} \end{bmatrix} = \operatorname{diag}(\frac{i}{i+1}, i = 1 \dots n-1) I_{n-1}^{1}(Z_{n}^{T} D_{x} \mathbf{q}^{o})$$

Set

$$\mathbf{b} = \begin{bmatrix} B_0(0) \\ B_2(0) \\ B_4(0) \\ \vdots \end{bmatrix}, \quad D_x = \operatorname{diag}(\frac{x^i}{(2i)!}, i = 0, 1, 2, \ldots).$$

Then

$$L(\alpha)D_x\mathbf{b} = D_x\mathbf{q}, \quad L(\alpha)Z^TD_x\mathbf{b} = d(\mathbf{z})Z^TD_x\mathbf{q},$$

where the vectors  $\alpha = (\alpha_i)_{i=0}^{+\infty}$ ,  $\mathbf{q} = (q_i)_{i=0}^{+\infty}$ , and  $\mathbf{z} = (z_i)_{i=1}^{+\infty}$ , can assume respectively the values:

$$\begin{array}{c} \alpha_i^R = \delta_{i=0 \bmod 3} \frac{2x^i}{(2i+2)!(\frac{2}{3}i+1)}, \quad q_i^R = \frac{1}{(2i+1)(i+1)} (1 - \delta_{i=2 \bmod 3} \frac{3}{2}), \ i = 0, 1, 2, 3, \dots \\ z_i^R = 1 - \delta_{i=0 \bmod 3} \frac{1}{\frac{2}{3}i+1}, \ i = 1, 2, 3, \dots, \end{array}$$

$$\alpha_i^e = \frac{x^i}{(2i+2)!}, \quad q_i^e = \frac{1}{2(2i+1)}, \quad i = 0, 1, 2, 3, \dots$$
$$z_i^e = \frac{i}{i+1}, \quad i = 1, 2, 3, \dots,$$

$$\alpha_i^o = \frac{x^i}{(2i+1)!}, \ i = 0, 1, 2, 3, \dots, \ q_0^o = 1, \ q_i^o = \frac{1}{2}, \ i = 1, 2, 3, \dots$$
$$z_i^o = \frac{2i-1}{2i+1}, \ i = 1, 2, 3, \dots$$

( Z is the semi-infinite lower shift matrix

$$Z = \left[ \begin{array}{ccc} 0 & & \\ 1 & 0 & \\ & 1 & 0 \\ & & \cdot & \cdot \end{array} \right],$$

 $L(\alpha)$  is the semi-infinite lower triangular Toeplitz matrix with first column  $\alpha$ , i.e.

$$L(\alpha) = \sum_{i=0}^{+\infty} \alpha_i Z^i = \begin{bmatrix} \alpha_0 \\ \alpha_1 & \alpha_0 \\ \alpha_2 & \alpha_1 & \alpha_0 \\ \alpha_3 & \alpha_2 & \alpha_1 & \alpha_0 \\ \vdots & \vdots & \ddots & \vdots \end{bmatrix}.$$

Exercise.

Try to prove that the system  $L(\alpha^R)D_x\mathbf{b} = D_x\mathbf{q}^R$  (the lower triangular Toeplitz sparse system equivalent to the Ramanujan lower triangular sparse system) can be obtained as a consequence of the system  $L(\alpha^o)D_x\mathbf{b} = D_x\mathbf{q}^o$  (or  $L(\alpha^e)D_x\mathbf{b} = D_x\mathbf{q}^e$ ).

For example, try to find  $\beta^e$ ,  $\beta^o$  such that  $L(\beta^e)L(\alpha^e) = L(\alpha^R) = L(\beta^o)L(\alpha^o)$  [??  $L(\beta^e)D_x\mathbf{q}^e = D_x\mathbf{q}^R = L(\beta^o)D_x\mathbf{q}^o$  ??], i.e.  $L(\alpha^e)\beta^e = \alpha^R = L(\alpha^o)\beta^o$ .

Find  $\gamma$  such that  $L(\gamma)L(\alpha^e)=L(\alpha)$  (or  $L(\gamma)L(\alpha^o)=L(\alpha)$ ) with  $\alpha$  more sparse than  $\alpha^R$ .

Exercise.

Prove that  $d(\mathbf{z})Z^TL(\alpha)(D_x\mathbf{b}) = L(\alpha)Z^T(D_x\mathbf{b}).$ 

Note that

$$L(\mathbf{a}) = \begin{bmatrix} 1 & & & & & \\ a_1 & 1 & & & & \\ a_2 & a_1 & 1 & & & \\ a_3 & a_2 & a_1 & 1 & & \\ a_4 & a_3 & a_2 & a_1 & 1 & \\ a_5 & a_4 & a_3 & a_2 & a_1 & 1 \\ & & & & & & & & & \\ \end{bmatrix} \begin{bmatrix} 1 \\ -a_1 \\ a_2 \\ -a_3 \\ a_4 \\ -a_5 \\ & & \end{bmatrix} = \begin{bmatrix} 1 \\ 2a_2 - a_1^2 \\ 0 \\ 2a_4 - 2a_1a_3 + a_2^2 \\ 0 \\ 2a_6 - 2a_1a_5 + 2a_2a_4 - a_3^2 \\ 0 \\ 2a_8 - 2a_1a_7 + 2a_2a_6 - 2a_3a_5 + a_4^2 \\ 0 \\ & & & \\ \end{bmatrix},$$

$$L(\mathbf{a})\left(\mathbf{e}_{1} + \sum_{i=1}^{+\infty} (-1)^{i} a_{i} \mathbf{e}_{i+1}\right) = \mathbf{e}_{1} + \sum_{i=1}^{+\infty} \delta_{i=0 \bmod 2} \left(2a_{i} + \sum_{i=1}^{i-1} (-1)^{i} a_{j} a_{i-j}\right) \mathbf{e}_{i+1}.$$

If we introduce the following two semi-infinite matrices,

$$I_{\pm} = \operatorname{diag}((-1)^{i}, i = 0, 1, 2, 3, \ldots), \quad E = \begin{bmatrix} 1 & 0 & & \\ 0 & 0 & & \\ 0 & 1 & 0 & \\ 0 & 0 & 0 & \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \end{bmatrix},$$

then the previous equality can be rewritten as

$$L(\mathbf{a})I_{\pm}\mathbf{a} = E\mathbf{a}^{(1)}, \ a_0^{(1)} = 1, \ a_s^{(1)} = 2a_{2s} + \sum_{j=1}^{2s-1} (-1)^j a_j a_{2s-j}, \ s = 1, 2, 3, \dots$$

Of course,  $L(\mathbf{a})L(I_{\pm}\mathbf{a}) = L(I_{\pm}\mathbf{a})L(\mathbf{a}) = L(E\mathbf{a}^{(1)})$ , so, for example, a lower triangular Toeplitz (l.t.T.) system of the type  $L(\mathbf{a})\mathbf{z} = \mathbf{c}$  is equivalent to the l.t.T. system  $L(E\mathbf{a}^{(1)})\mathbf{z} = L(I_{\pm}\mathbf{a})\mathbf{c}$  whose coefficient matrix has null diagonals alternating with non-null ones. We shall introduce, as a consequence of this result, an algorithm of complexity  $O(n(\log_2 n)^2)$  for solving a lower triangular Toeplitz system (of order  $n = 2^k$ ).

Exercise. Apply the above result to the odd (even) system (the one solved by Bernoulli numbers) in order to obtain a system where null diagonals alternate with the non-null ones. Is it possible to write an explicit formula for the entries of the non-null diagonals?

Exercise. Use twice the above result in order to show that a l.t.T. system of the type  $L(\mathbf{a})\mathbf{z} = \mathbf{c}$  is equivalent to a l.t.T. system where the coefficient matrix has three-null diagonals alternating with non-null ones.

Exercise. Answer to the quotation marks in the following equality:

$$L(\mathbf{a}) = \begin{bmatrix} 1 & & & & & \\ a_1 & 1 & & & & \\ a_2 & a_1 & 1 & & & \\ a_3 & a_2 & a_1 & 1 & & \\ a_4 & a_3 & a_2 & a_1 & 1 & \\ a_5 & a_4 & a_3 & a_2 & a_1 & 1 \\ & & & & & & & & & & & & & & \end{bmatrix} \begin{bmatrix} 1 \\ ? \\ ? \\ ? \\ ? \\ ? \\ ? \\ . \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ ? \\ ? \\ 0 \\ 0 \\ ? \\ . \end{bmatrix}.$$

As a consequence of this result,

- (i) introduce an algorithm of complexity  $O(n(\log_3 n)^2)$  for solving a lower triangular Toeplitz system (for example, in case the system has order n=27 such algorithm is more efficient rather than embed the system in one of order n=32 and use the algorithm ok for  $n=2^k$ );
- (ii) obtain the Ramanujan system solved by Bernoulli numbers as a consequence of the odd (even) system.

A lower triangular Toeplitz system solver of complexity  $O(n(\log_2 n)^2)$  Set

$$n = 2^{k}, \ \mathbf{a} = \mathbf{a}^{(0)} = \begin{bmatrix} 1 \\ a_{1} \\ \vdots \\ a_{n-1} \end{bmatrix} = \begin{bmatrix} 1 \\ a_{1}^{(0)} \\ \vdots \\ a_{n-1}^{(0)} \end{bmatrix}, \ L(\mathbf{a}) = \begin{bmatrix} 1 \\ a_{1} & 1 \\ a_{2} & a_{1} & 1 \\ \vdots & \vdots & \ddots & \vdots \\ a_{n-1} & \cdot & a_{2} & a_{1} & 1 \end{bmatrix},$$

and let  $I_{\pm}$  and E denote the  $n \times n$  upper-left submatrices of the previously defined semi-infinite matrices  $I_{\pm}$  and E. In the following we multiply  $L(\mathbf{a})$  on the left by suitable l.t.T matrices so to reduce the number of its non-null diagonals; we do this in  $k = \log_2 n$  steps:

step 1

$$\begin{array}{l} L(\mathbf{a})I_{\pm}\mathbf{a} =: E\mathbf{a}^{(1)}, \\ a_0^{(1)} = 1, \ a_s^{(1)} = 2a_{2s}^{(0)} + \sum_{j=1}^{2s-1} (-1)^j a_j^{(0)} a_{2s-j}^{(0)}, \ s = 1, \dots, \frac{n}{2} - 1, \ a_s^{(1)} = 0, \ s = \frac{n}{2}, \dots, n-1, \\ L(\mathbf{a})L(I_{\pm}\mathbf{a}) = L(I_{\pm}\mathbf{a})L(\mathbf{a}) = L(E\mathbf{a}^{(1)}) \end{array}$$

Note: One here needs to compute  $\mathbf{a}^{(1)}$ , i.e.  $(\frac{n}{2} - 1 \text{ odd entries of})$  the vector  $L(\mathbf{a}) \cdot I_{\pm}\mathbf{a}$ . Denote by  $\varphi_n$  the cost (the number of multiplications) of such operation.

step 2

$$\begin{array}{l} L(E\mathbf{a}^{(1)})EI_{\pm}\mathbf{a}^{(1)} =: E^2\mathbf{a}^{(2)}, \\ a_0^{(2)} = 1, \ a_s^{(2)} = 2a_{2s}^{(1)} + \sum_{j=1}^{2s-1} (-1)^j a_j^{(1)} a_{2s-j}^{(1)}, \ s = 1, \dots, \frac{n}{4} - 1, \ a_s^{(2)} = 0, \ s = \frac{n}{4}, \dots, n-1, \\ L(E\mathbf{a}^{(1)})L(EI_{\pm}\mathbf{a}^{(1)}) = L(EI_{\pm}\mathbf{a}^{(1)})L(E\mathbf{a}^{(1)}) = L(E^2\mathbf{a}^{(2)}) \end{array}$$

Note: One here needs to compute  $\mathbf{a}^{(2)}$ , i.e.  $(\frac{n}{4}-1 \text{ odd entries of})$  the order  $\frac{n}{2}$  non null first part of the vector  $L(\mathbf{a}^{(1)}) \cdot I_{\pm} \mathbf{a}^{(1)}$ . Denote by  $\varphi_{\frac{n}{2}}$  the cost of such operation.

step 3

$$L(E^{2}\mathbf{a}^{(2)})E^{2}I_{\pm}\mathbf{a}^{(2)} =: E^{3}\mathbf{a}^{(3)},$$

$$a_{0}^{(3)} = 1, \ a_{s}^{(3)} = 2a_{2s}^{(2)} + \sum_{j=1}^{2s-1} (-1)^{j} a_{j}^{(2)} a_{2s-j}^{(2)}, \ s = 1, \dots, \frac{n}{8} - 1, \ a_{s}^{(3)} = 0, \ s = \frac{n}{8}, \dots, n-1,$$

$$L(E^{2}\mathbf{a}^{(2)})L(E^{2}I_{\pm}\mathbf{a}^{(2)}) = L(E^{2}I_{\pm}\mathbf{a}^{(2)})L(E^{2}\mathbf{a}^{(2)}) = L(E^{3}\mathbf{a}^{(3)})$$

Note: One here needs to compute  $\mathbf{a}^{(3)}$ , i.e.  $(\frac{n}{8}-1 \text{ odd entries of})$  the order  $\frac{n}{4}$  non null first part of the vector  $L(\mathbf{a}^{(2)}) \cdot I_{\pm} \mathbf{a}^{(2)}$ . Denote by  $\varphi_{\frac{n}{4}}$  the cost of such operation.

. . .

$$\begin{split} & \text{step } k-1 = \log_2 n - 1 \\ & L(E^{k-2}\mathbf{a}^{(k-2)})E^{k-2}I_{\pm}\mathbf{a}^{(k-2)} =: E^{k-1}\mathbf{a}^{(k-1)}, \\ & a_0^{(k-1)} = 1, \\ & a_s^{(k-1)} = 2a_{2s}^{(k-2)} + \sum_{j=1}^{2s-1} (-1)^j a_j^{(k-2)} a_{2s-j}^{(k-2)}, \ s = 1, \dots, \frac{n}{2^{k-1}} - 1, \\ & a_s^{(k-1)} = 0, \ s = \frac{n}{2^{k-1}}, \dots, n-1, \\ & L(E^{k-2}\mathbf{a}^{(k-2)})L(E^{k-2}I_{\pm}\mathbf{a}^{(k-2)}) = L(E^{k-2}I_{\pm}\mathbf{a}^{(k-2)})L(E^{k-2}\mathbf{a}^{(k-2)}) = L(E^{k-1}\mathbf{a}^{(k-1)}) \end{split}$$

Note: One here needs to compute  $\mathbf{a}^{(k-1)}$ , i.e.  $(\frac{n}{2^{k-1}}-1)$  odd entries of) the order  $\frac{n}{2^{k-2}}$  non null first part of the vector  $L(\mathbf{a}^{(k-2)}) \cdot I_{\pm} \mathbf{a}^{(k-2)}$ . Denote by  $\varphi_{\frac{n}{2^{k-2}}}$  the cost of such operation.

step  $k = \log_2 n$ 

$$\begin{split} &L(E^{k-1}\mathbf{a}^{(k-1)})E^{k-1}I_{\pm}\mathbf{a}^{(k-1)} =: E^{k}\mathbf{a}^{(k)}, \\ &a_{0}^{(k)} = 1, \\ &a_{s}^{(k)} = 0, \ s = \frac{n}{2^{k}}, \dots, n-1, \\ &L(E^{k-1}\mathbf{a}^{(k-1)})L(E^{k-1}I_{\pm}\mathbf{a}^{(k-1)}) = L(E^{k-1}I_{\pm}\mathbf{a}^{(k-1)})L(E^{k-1}\mathbf{a}^{(k-1)}) = L(E^{k}\mathbf{a}^{(k)}) \end{split}$$

$$(E^k = \mathbf{e}_1 \mathbf{e}_1^T, \, \mathbf{a}^{(k)} = \mathbf{e}_1, \, L(E^k \mathbf{a}^{(k)}) = I)$$

Note: One here needs to compute  $\mathbf{a}^{(k)}$ , i.e.  $(\frac{n}{2^k} - 1 \text{ odd entries of})$  the order  $\frac{n}{2^{k-1}}$  non null first part of the vector  $L(\mathbf{a}^{(k-1)}) \cdot I_{\pm} \mathbf{a}^{(k-1)}$ . Denote by  $\varphi_{\frac{n}{2^{k-1}}}$  the cost of such operation  $(\varphi_{\frac{n}{2^{k-1}}} = 0)$ .

Finally note that the following equality holds

$$L(E^{k-1}I_{\pm}\mathbf{a}^{(k-1)})L(E^{k-2}I_{\pm}\mathbf{a}^{(k-2)})\cdots L(E^{2}I_{\pm}\mathbf{a}^{(2)})L(EI_{\pm}\mathbf{a}^{(1)})L(I_{\pm}\mathbf{a})\Big[L(\mathbf{a})\Big]=I,$$

in other words we have transformed  $L(\mathbf{a})$  into the identity matrix.

Example: n = 8

$$n=2^3, \ \mathbf{a}=\mathbf{a}^{(0)}=\begin{bmatrix} 1\\ a_1\\ a_2\\ a_3\\ a_4\\ a_5\\ a_6\\ a_7 \end{bmatrix}=\begin{bmatrix} 1\\ a_1^{(0)}\\ a_2^{(0)}\\ a_3^{(0)}\\ a_3^{(0)}\\ a_3^{(0)}\\ a_6^{(0)}\\ a_6^{(0)}\\ a_7^{(0)} \end{bmatrix}, \ L(\mathbf{a})=\begin{bmatrix} 1\\ a_1 & 1\\ a_2 & a_1 & 1\\ a_3 & a_2 & a_1 & 1\\ a_4 & a_3 & a_2 & a_1 & 1\\ a_5 & a_4 & a_3 & a_2 & a_1 & 1\\ a_6 & a_5 & a_4 & a_3 & a_2 & a_1 & 1\\ a_7 & a_6 & a_5 & a_4 & a_3 & a_2 & a_1 & 1 \end{bmatrix}$$

and let  $I_{\pm}$  and E denote the  $8 \times 8$  upper-left submatrices of the previously defined semi-infinite matrices  $I_{\pm}$  and E.

step 1

$$L(\mathbf{a})I_{\pm}\mathbf{a} = \begin{bmatrix} 1 \\ a_1 & 1 \\ a_2 & a_1 & 1 \\ a_3 & a_2 & a_1 & 1 \\ a_4 & a_3 & a_2 & a_1 & 1 \\ a_5 & a_4 & a_3 & a_2 & a_1 & 1 \\ a_7 & a_6 & a_5 & a_4 & a_3 & a_2 & a_1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ -a_1 \\ a_2 \\ -a_3 \\ a_4 \\ -a_5 \\ a_6 \\ -a_7 \end{bmatrix} =: E\mathbf{a}^{(1)} = \begin{bmatrix} 1 \\ 0 \\ a_1^{(1)} \\ 0 \\ a_2^{(1)} \\ 0 \\ 0 \end{bmatrix},$$
 
$$a_0^{(1)} = 1, \ a_s^{(1)} = 2a_{2s}^{(0)} + \sum_{j=1}^{2s-1} (-1)^j a_j^{(0)} a_{2s-j}^{(0)}, \ s = 1, \dots, \frac{8}{2} - 1, \ a_s^{(1)} = 0, \ s = \frac{8}{2}, \dots, 8 - 1$$
 
$$a_1^{(1)} = 2a_2 - a_1^2, \ a_2^{(1)} = 2a_4 - 2a_1a_3 + a_2^2, \ a_3^{(1)} = 2a_6 - 2a_1a_5 + 2a_2a_4 - a_3^2,$$
 
$$L(\mathbf{a})L(I_{\pm}\mathbf{a}) = L(I_{\pm}\mathbf{a})L(\mathbf{a}) = L(E\mathbf{a}^{(1)})$$

Note: One here needs to compute  $\mathbf{a}^{(1)}$ , i.e.  $(\frac{8}{2}-1 \text{ odd entries of})$  the vector  $L(\mathbf{a}) \cdot I_{\pm}\mathbf{a}$ . Denote by  $\varphi_8$  the cost of such operation  $(\varphi_8 \leq 6 \text{ if multiplications})$  by 2 are not counted).

step 2

$$L(E\mathbf{a}^{(1)})EI_{\pm}\mathbf{a}^{(1)} = \begin{bmatrix} 1 & & & & \\ & 1 & & & & \\ & a_1^{(1)} & 1 & & & \\ & a_1^{(1)} & 1 & & & \\ & a_1^{(1)} & 1 & & & \\ & a_2^{(1)} & a_1^{(1)} & 1 & & \\ & a_2^{(1)} & a_1^{(1)} & 1 & \\ & a_3^{(1)} & a_2^{(1)} & a_1^{(1)} & 1 \\ & a_3^{(1)} & a_2^{(1)} & a_1^{(1)} & 1 \end{bmatrix} \begin{bmatrix} 1 & \\ 0 & \\ -a_1^{(1)} \\ 0 & \\ -a_3^{(1)} & 0 \\ 0 & \end{bmatrix} = : E^2\mathbf{a}^{(2)} = \begin{bmatrix} 1 & \\ 0 & \\ 0 & \\ 0 & \\ 0 & \\ 0 & \end{bmatrix},$$
 
$$\mathbf{a}_0^{(2)} = 1, \ a_s^{(2)} = 2a_{2s}^{(1)} + \sum_{j=1}^{2s-1} (-1)^j a_j^{(1)} a_{2s-j}^{(1)}, \ s = 1, \dots, \frac{8}{4} - 1, \ a_s^{(2)} = 0, \ s = \frac{8}{4}, \dots, 8 - 1,$$
 
$$a_1^{(2)} = 2a_2^{(1)} - (a_1^{(1)})^2,$$
 
$$L(E\mathbf{a}^{(1)})L(EI_{\pm}\mathbf{a}^{(1)}) = L(EI_{\pm}\mathbf{a}^{(1)})L(E\mathbf{a}^{(1)}) = L(E^2\mathbf{a}^{(2)})$$

Note: One here needs to compute  $\mathbf{a}^{(2)}$ , i.e.  $(\frac{8}{4}-1 \text{ odd entries of})$  the order  $\frac{8}{2}$  non null first part of the vector  $L(\mathbf{a}^{(1)}) \cdot I_{\pm} \mathbf{a}^{(1)}$ . Denote by  $\varphi_{\frac{8}{2}}$  the cost of such operation  $(\varphi_{\frac{8}{2}} \leq 1 \text{ if multiplications by 2 are not counted}).$ 

step  $3 = \log_2 8$ 

$$L(E^2\mathbf{a}^{(2)})E^2I_{\pm}\mathbf{a}^{(2)} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ a_1^{(2)} \\ a_1^{(2)} \\ a_1^{(2)} \\ a_1^{(2)} \\ a_1^{(2)} \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \\ -a_1^{(2)} \\ 0 \\ 0 \\ 0 \end{bmatrix} =: E^3\mathbf{a}^{(3)} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix},$$
 
$$a_0^{(3)} = 1, \ a_s^{(3)} = 2a_{2s}^{(2)} + \sum_{j=1}^{2s-1} (-1)^j a_j^{(2)} a_{2s-j}^{(2)}, \ s = 1, \dots, \frac{8}{8} - 1, \ a_s^{(3)} = 0, \ s = \frac{8}{8}, \dots, 8 - 1,$$
 
$$L(E^2\mathbf{a}^{(2)})L(E^2I_{\pm}\mathbf{a}^{(2)}) = L(E^2I_{\pm}\mathbf{a}^{(2)})L(E^2\mathbf{a}^{(2)}) = L(E^3\mathbf{a}^{(3)})$$

 $(E^3=\mathbf{e}_1\mathbf{e}_1^T,\,\mathbf{a}^{(3)}=\mathbf{e}_1,\,L(E^3\mathbf{a}^{(3)})=I)$  Note: One here needs to compute  $\mathbf{a}^{(3)}$ , i.e.  $(\frac{8}{8}-1 \text{ odd entries of})$  the order  $\frac{8}{4}$  non null first part of the vector  $L(\mathbf{a}^{(2)})\cdot I_\pm\mathbf{a}^{(2)}$ . Denote by  $\varphi_{\frac{8}{4}}$  the cost of such operation  $(\varphi_{\frac{8}{4}}=0)$ .

By using the obtained identities, one realizes that

$$L(E^2I_{\pm}\mathbf{a}^{(2)})L(EI_{\pm}\mathbf{a}^{(1)})L(I_{\pm}\mathbf{a})[L(\mathbf{a})] = I,$$

i.e. we have performed a kind of Gaussian elimination . . .

Upper bounds for the cost  $\sum_{j=k}^{1} \varphi_{2^{j}}$  of the computation of  $\mathbf{a}^{(1)}, \mathbf{a}^{(2)}, \dots, \mathbf{a}^{(k-1)}, \mathbf{a}^{(k)}$ 

Since the cost of a matrix-vector product involving a generic vector and a  $2^j \times 2^j$  lower triangular Toeplitz matrix with ones on the diagonal is obviously bounded by  $1+2+3+\ldots+(2^j-1)=\frac{1}{2}2^j(2^j-1)$ , we can say that

$$\sum_{i=1}^{k} \varphi_{2i} \le \frac{1}{2} \sum_{i=1}^{k} (2^{2j} - 2^{j}) = \frac{2}{3} (2^{k})^{2} - 2^{k} + \frac{1}{3}.$$

However, by exploiting the particularity of our matrix-vector products, we observe that  $\varphi_2 = 0$ ,  $\varphi_4 \le 1$ ,  $\varphi_8 \le 3 + 1 + 2$ ,  $\varphi_{16} \le 7 + 1 + 2 + 3 + 4 + 5 + 6$ ,  $\varphi_{2^j} \le \frac{1}{2} 2^{j-1} (2^{j-1} - 1)$ . So,

$$\sum_{i=1}^{k} \varphi_{2^{j}} \le \frac{1}{2} \sum_{i=2}^{k} (2^{2j-2} - 2^{j-1}) = \frac{1}{6} (2^{k})^{2} - \frac{1}{2} 2^{k} + \frac{1}{3}.$$

Finally, since a matrix-vector product involving a generic vector and a  $2^j \times 2^j$  lower triangular Toeplitz matrix can be computed by means of the three FFT of order  $2^{j+1}$ , i.e. with a cost  $O((j+1)2^{j+1})$ , we have

$$\sum_{j=1}^{k} \varphi_{2^{j}} = O(k^{2} 2^{k}) = O(n(\log_{2} n)^{2}) \text{ prove this!}$$

Computing the first column of  $L(\mathbf{a})^{-1}$ 

Now observe that

$$\begin{split} L(E^{k-1}I_{\pm}\mathbf{a}^{(k-1)})L(E^{k-2}I_{\pm}\mathbf{a}^{(k-2)})\cdots L(E^{2}I_{\pm}\mathbf{a}^{(2)})L(EI_{\pm}\mathbf{a}^{(1)})L(I_{\pm}\mathbf{a}) & \left[L(\mathbf{a})\right] = I, \\ L(\mathbf{a})\mathbf{z} & = \mathbf{c}, \\ \mathbf{z} & = L(E^{k-1}I_{\pm}\mathbf{a}^{(k-1)})L(E^{k-2}I_{\pm}\mathbf{a}^{(k-2)})\cdots L(E^{2}I_{\pm}\mathbf{a}^{(2)})L(EI_{\pm}\mathbf{a}^{(1)})L(I_{\pm}\mathbf{a})\mathbf{c} \\ & = L(I_{\pm}\mathbf{a})L(EI_{\pm}\mathbf{a}^{(1)})L(E^{2}I_{\pm}\mathbf{a}^{(2)})\cdots L(E^{k-2}I_{\pm}\mathbf{a}^{(k-2)})L(E^{k-1}I_{\pm}\mathbf{a}^{(k-1)})\mathbf{c}, \end{split}$$

$$\begin{split} L(\mathbf{a})\mathbf{z} &= u\mathbf{e}_1 + v\mathbf{e}_{\frac{n}{2}+1}, \\ \mathbf{z} &= L(I_{\pm}\mathbf{a})L(EI_{\pm}\mathbf{a}^{(1)})L(E^2I_{\pm}\mathbf{a}^{(2)})\cdots L(E^{k-2}I_{\pm}\mathbf{a}^{(k-2)})L(E^{k-1}I_{\pm}\mathbf{a}^{(k-1)})\left(u\mathbf{e}_1 + v\mathbf{e}_{\frac{n}{2}+1}\right) \\ &= L(I_{\pm}\mathbf{a})EL(I_{\pm}\mathbf{a}^{(1)})EL(I_{\pm}\mathbf{a}^{(2)})\cdots EL(I_{\pm}\mathbf{a}^{(k-2)})EL(I_{\pm}\mathbf{a}^{(k-1)})\left(u\mathbf{e}_1 + v\mathbf{e}_2\right). \end{split}$$

The latter identity, whose proof is left to the reader, implies that the solution of the system  $L(\mathbf{a})\mathbf{z} = u\mathbf{e}_1 + v\mathbf{e}_{\frac{n}{2}+1}$  (the first column of  $L(\mathbf{a})^{-1}$ , if u=1, v=0) can be computed by performing  $k = \log_2 n$  l.t.T.matrix-vector products where the matrices are, respectively,  $2 \times 2$  (no operation, if u=1, v=0),  $4 \times 4$  (one multiplication, if u=1, v=0),  $8 \times 8, \ldots, 2^k \times 2^k$ , i.e.  $2^j \times 2^j, j=1,2,\ldots,k$ . If we perform such products by means of FFT (as soon as j becomes so large that using FFT is cheaper than the economized direct product), then the total amount of operations is  $O(n(\log_2 n)^2)$  (prove it!).

Example:  $n = 8 = 2^3, k = 3$ 

is such that

$$\begin{bmatrix} 1 & & & & & & & & \\ a_1 & 1 & & & & & & \\ a_2 & a_1 & 1 & & & & & \\ a_3 & a_2 & a_1 & 1 & & & & \\ a_4 & a_3 & a_2 & a_1 & 1 & & & \\ a_5 & a_4 & a_3 & a_2 & a_1 & 1 & & \\ a_6 & a_5 & a_4 & a_3 & a_2 & a_1 & 1 & & \\ a_7 & a_6 & a_5 & a_4 & a_3 & a_2 & a_1 & 1 \end{bmatrix} \begin{bmatrix} z_0 \\ z_1 \\ z_2 \\ z_3 \\ z_4 \\ z_5 \\ z_6 \\ z_7 \end{bmatrix} = \begin{bmatrix} u \\ 0 \\ 0 \\ v \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

How to determine **z** such that  $[L(\mathbf{a})\mathbf{z}]_i = 0, i = 1, 2, 4, 5, 7, 8, 10, 11, \dots$ ?

$$\begin{bmatrix} 1 & & & & & & \\ a_1 & 1 & & & & & \\ a_2 & a_1 & 1 & & & & \\ a_3 & a_2 & a_1 & 1 & & & \\ a_4 & a_3 & a_2 & a_1 & 1 & & \\ a_5 & a_4 & a_3 & a_2 & a_1 & 1 \\ & & & & & & & \\ \end{bmatrix} \begin{bmatrix} 1 \\ -a_1 \\ a_1^2 - a_2 \\ x \\ y \\ z \\ \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ a_3 - 2a_1a_2 + a_1^3 + x \\ a_4 - a_1a_3 + a_2a_1^2 - a_2^2 + a_1x + y \\ a_5 - a_1a_4 + a_3a_1^2 - a_2a_3 + a_2x + a_1y + z \end{bmatrix}$$

We have to choose x, y, z such that

$$a_4 - a_1 a_3 + a_2 a_1^2 - a_2^2 + a_1 x + y = 0,$$
  

$$a_5 - a_1 a_4 + a_3 a_1^2 - a_2 a_3 + a_2 x + a_1 y + z = 0.$$

Note that it must be  $y = -a_4 + a_2^2 - a_1t$ ,  $x = a_3 - a_1a_2 + t$ , and then z is determined by

$$z = -a_5 + a_1 a_4 - a_3 a_1^2 + a_2 a_3 - a_2 a_3 + a_1 a_2^2 - a_2 t + a_1 a_4 - a_1 a_2^2 + a_1^2 t = -a_5 + 2a_1 a_4 - a_3 a_1^2 - a_2 t + a_1^2 t.$$

A possible choice of t is  $t = a_3$ , so that in z there are not terms of type  $a_i^r a_j$  with r greater than 1:

$$z = -a_5 + 2a_1a_4 - a_2a_3.$$

Exercise. What choice of t allows one to obtain the Ramanujan system from the odd (even) system? ...

Proceed analogously in order to obtain  $[L(\mathbf{a})\mathbf{z}]_i = 0, i = 7, 8$ :

$$a_7 + a_1X + Y + a_1x + y = 0$$
,  $a_8 \dots + a_2x + a_1y + z = 0$ .

By setting x=-X+t,  $y=-a_7-Y-a_1t$  we satisfy the first equation, and the second equation becomes  $\ldots a_2t+a_6a_1^2-a_1^2t\ldots$ . Thus choose  $t=a_6$  so that the only "strange" term of the second equation,  $a_6a_1^2$ , is neglected, and x,y,z are determined. Idem in order to obtain  $[L(\mathbf{a})\mathbf{z}]_i=0$ , i=10,11. In such case the second equation becomes  $\ldots a_2t+a_9a_1^2-a_1^2t\ldots$  Thus choose  $t=a_9\ldots$ 

In this way we prove that

$$a_1^{(1)} = 3a_3 - 3a_1a_2 + a_1^3, ?a_2^{(1)} = 3a_6 - 3a_1a_5 - 3a_2a_4 + 3a_3^2 - 3a_1a_2a_3 + 3a_1^2a_4 + a_2^3?, a_3^{(1)} = ?.$$

Set  $a_0 = 1$ . Conjecture on the next entry of the solution:

$$2a_0a_{12} - a_1a_{11} - a_2a_{10} + 2a_3a_9 - a_4a_8 - a_5a_7 + a_6^2$$

Why? ...

Conjecture on the solution:

$$-\sum_{r=0}^{\lfloor \frac{i-1}{2} \rfloor} a_r a_{i-r} + \delta_{i=0 \bmod 2} a_{\frac{i}{2}}^2 + 3 \left\{ \begin{array}{ll} \sum_{s \geq \frac{3-i}{6}}^0 a_{\frac{i-3}{2} + 3s} a_{\frac{i+3}{2} - 3s} & i \text{ odd} \\ \sum_{s \geq \frac{6-i}{6}}^0 a_{\frac{i-6}{2} + 3s} a_{\frac{i+6}{2} - 3s} & i \text{ even} \end{array} \right., \ i = 0, 1, 2, 3, 4, 5, \ldots$$

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			- 0	$\begin{array}{c c} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$	1 a1	1																	$     \begin{array}{c}       a_3 \\       a_4 \\       a_5 \\       a_6 \\       a_7 \\       a_8 \\       a_9 \\       a_{10}     \end{array} $
			- 0	$\begin{array}{c c} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$ \begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \end{array} $	$\begin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$	$\frac{1}{a_1}$	1																$egin{array}{c} a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ \end{array}$
			- 0	$\begin{array}{c c} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$	_		$\frac{1}{a_1}$	1															$egin{array}{c} a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ a_{12} \\ \end{array}$
_			- 0	$\begin{array}{c c} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$		$\frac{1}{a_1}$	1														$egin{array}{c} a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ \end{array}$
_			- 0	$\begin{array}{c c} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$	$a_1 \\ a_2 \\ a_3 \\ a_4$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$		1 a1	1													$egin{array}{c} a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ a_{12} \\ a_{13} \\ \end{array}$
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ \end{array}$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$	-	1 a1	1												$egin{array}{c} a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ a_{12} \\ a_{13} \\ a_{14} \\ \end{array}$
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$	-	1 a1	1											$\begin{array}{c c} a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ a_{12} \\ a_{13} \\ a_{14} \\ a_{15} \end{array}$
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$	$a_1$ $a_2$ $a_3$ $a_4$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$		1 a1	1										a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$	$a_1$ $a_2$ $a_3$ $a_4$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$	$a_1$		1									a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16       a17
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$	$a_1$ $a_2$ $a_3$ $a_4$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$	$a_1$ $a_2$	$a_1$		_								a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16       a17       a18
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$ $a_7$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$	$a_1 \\ a_2 \\ a_3 \\ a_4$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$ $a_3$	$a_1$ $a_2$ $a_2$									) ( ( ( ( ( ( ( ( ( ( ( ( ( ( ( ( ( ( (	a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16       a17       a18       a19
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 & & & & \\ a_2 & & & & \\ a_3 & & & & \\ a_4 & & & & \\ a_5 & & & & \\ a_6 & & & & \\ a_7 & & & & \\ a_8 & & & & \\ \end{array}$	$egin{array}{c c} a_1 & & \\ a_2 & & \\ a_3 & & \\ a_4 & & \\ a_5 & & \\ a_6 & & \\ a_7 & & \\ \hline \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 & & & & \\ a_2 & & & & \\ a_3 & & & & \\ a_4 & & & & \\ a_5 & & & & \\ \end{array}$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$a_1$ $a_2$ $a_3$ $a_4$	$\frac{1}{2}$ $\frac{1}{a}$	1							a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16       a17       a18       a19       a20       a21       a22
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	$egin{array}{c} a_1 & & & & \\ a_2 & & & & \\ a_3 & & & & \\ a_4 & & & & \\ a_5 & & & & \\ a_6 & & & & \\ a_7 & & & & \\ a_8 & & & & \\ a_9 & & & & \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$ $a_7$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ \end{array}$	$egin{array}{c cccc} a_1 & & & & & \\ \hline a_2 & a_2 & & & & \\ a_3 & a_4 & & & \\ \hline a_4 & a_5 & & & \\ \hline \end{array}$	$a_1$ $a_2$ $a_3$ $a_4$ $a_4$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$\begin{array}{c c} 1 & 1 \\ 2 & a_1 \end{array}$							a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16       a17       a18       a19       a20       a21       a22       a23
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	$egin{array}{c} a_1 & & & \\ a_2 & & & \\ a_3 & & & \\ a_4 & & & \\ a_5 & & & \\ a_6 & & & \\ a_7 & & & \\ a_8 & & & \\ a_9 & & & \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$egin{array}{c c} a_1 & & \\ a_2 & & \\ a_3 & & \\ a_4 & & \\ a_5 & & \\ a_6 & & \\ a_7 & & \\ \end{array}$	$egin{array}{c cccc} a_1 & & & & \\ a_2 & a_3 & a_4 & a_5 & a_6 & a_6 & a_6 & a_6 \\ \hline \end{array}$	$egin{array}{cccccccccccccccccccccccccccccccccccc$	$\begin{bmatrix} 1 & 1 \\ 2 & a \\ 3 & a_2 \\ 4 & a_3 \end{bmatrix}$	$\begin{array}{c cc} & 1 & 1 \\ 2 & a_1 & a_2 & a_2 \end{array}$	$a_1$						a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16       a17       a18       a19       a20       a21       a22       a23       a24
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	a1       a2       a3       a4       a5       a6       a7       a8       a9	a1       a2       a3       a4       a5       a6       a7       a8	$egin{array}{cccccccccccccccccccccccccccccccccccc$	$     \begin{array}{ccccccccccccccccccccccccccccccccc$	$\begin{bmatrix} 1 & 1 \\ 2 & a \end{bmatrix}$ $\begin{bmatrix} 1 & 3 & a_2 \\ 4 & a_3 \end{bmatrix}$ $\begin{bmatrix} 1 & 3 & a_2 \\ 4 & a_3 \end{bmatrix}$	$\begin{bmatrix} 1 & 1 \\ 2 & a_1 \\ 3 & a_2 \\ 4 & a_3 \end{bmatrix}$	$a_1$ $a_2$	$a_1$	11	_			a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16       a17       a18       a19       a20       a21       a22       a23       a24       a25
_			- 0	$\begin{array}{c cc} a_1 & 1 \\ a_2 & a \\ \hline & a_1 \end{array}$	$\begin{array}{c cccc} 1 & 1 & 1 & & & & & & & & & & & & & & $	$\begin{array}{c c} 1 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \end{array}$	$ \begin{array}{c c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \end{array} $	$a_1$ $a_2$ $a_3$ $a_4$ $a_5$ $a_6$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ a_{11} \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ a_9 \\ a_{10} \\ \end{array}$	$egin{array}{c} a_1 & & & \\ a_2 & & & \\ a_3 & & & \\ a_4 & & & \\ a_5 & & & \\ a_6 & & & \\ a_7 & & & \\ a_8 & & & \\ a_9 & & & \\ \end{array}$	$egin{array}{c} a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \\ a_8 \\ \end{array}$	a1       a2       a3       a4       a5       a6       a7       a8	$egin{array}{c cccc} a_1 & & & & & & \\ \hline a_2 & a_2 & & & & \\ \hline a_3 & a_4 & & & & \\ \hline a_4 & a_5 & & & & \\ \hline a_6 & a_7 & & & \\ \hline a_8 & & & & \\ \hline \end{array}$	$egin{array}{cccccccccccccccccccccccccccccccccccc$	$\begin{bmatrix} 1 & 1 \\ 2 & a \\ 3 & a_2 \\ 4 & a_3 \end{bmatrix}$	$\begin{bmatrix} 1 & 1 \\ 2 & a_1 \\ 3 & a_2 \\ 4 & a_3 \\ 5 & a_4 \end{bmatrix}$	$a_1$ $a_2$	$a_1$	+	. 1			a3       a4       a5       a6       a7       a8       a9       a10       a11       a12       a13       a14       a15       a16       a17       a18       a19       a20       a21       a22       a23       a24

QUESTION Can the above  $3^j \times 3^j$  matrix by vector product be computed in at most  $O(3^j j)$  arithmetic operations ?

If yes, then we would have a method which solves  $3^k \times 3^k$  lower triangular Toeplitz linear systems in at most  $O(3^k k)$  arithmetic operations.

If no, then look for another solution  $\hat{\mathbf{a}}$  of the system  $L(\mathbf{a})\hat{\mathbf{a}} = [1\ 0\ 0\bullet 0\ 0\cdot]^T$  such that  $\{\hat{\mathbf{a}}\}_{3^j}$  is computable from  $\{\mathbf{a}\}_{3^j}$  in at most  $O(3^jj)$  arithmetic operations.